# **Numerical Methods for Engineers**



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### **Preface**

[View the promotional video on YouTube](https://youtu.be/qFJGMBDfFMY)

These are the lecture notes for my upcoming Coursera course , *[Numerical Methods for](https://www.coursera.org/learn/numerical-methods-engineers/) [Engineers](https://www.coursera.org/learn/numerical-methods-engineers/)* (for release in January 2021). Before students take this course, they should have some basic knowledge of single-variable calculus, vector calculus, differential equations and matrix algebra. Students should also be familiar with at least one programming language. In this course, however, I will exclusively use MATLAB. I teach the basics of Matlab in the first week, but this may be too short an introduction for students of limited programming ability and they may need to supplement their programming lessons elsewhere.

I have divided these notes into chapters called Lectures, with each Lecture corresponding to a video on Coursera. I have also uploaded all my Coursera videos to YouTube, and links are placed at the top of each Lecture.

There are problems at the end of each lecture, some that require analytical solutions and others that require Matlab programs. Solutions to the analytical questions and Learner Templates for the MATLAB programs can be found in the Appendix.

On the Coursera platform, at the end of each week there is also both an assessed multiple-choice quiz and a Matlab project. Details of the Matlab projects and their Learner Templates can also be found in these lecture notes.

> Jeffrey R. Chasnov Hong Kong Nov 2020

# **Contents**









### **Week I**

## <span id="page-7-0"></span>**Scientific Computing**

In this week's lectures, we learn how to program using MATLAB. We learn how real numbers are *represented in double precision and how to do basic arithmetic with* Matlab*. We learn how to use scripts and functions, how to represent vectors and matrices, how to draw line plots, how to use logical variables, conditional statements,* for *loops and* while *loops.*

*Your programming project will be to write a* Matlab *code to compute the bifurcation diagram for the logistic map.*

### <span id="page-8-0"></span>**Lecture 1 | Binary numbers**

[View this lecture on YouTube](https://youtu.be/2zd-TgHJgVU)

We do our arithmetic using decimals, which is a base-ten positional number system. For example, the meaning of the usual decimal notation is illustrated by

$$
524.503=5\times 10^2 + 2\times 10^1 + 4\times 10^0 + 5\times 10^{-1} + 0\times 10^{-2} + 3\times 10^{-3}.
$$

Each position in a decimal number corresponds to a power of 10. A digit in this number system is defined as any of the numerals from 0 to 9, while digit is also the English word meaning a finger or a toe. Decimals probably arose from counting using ten fingers.

Computers count using two states, and this has led to the use of binary numbers, which is a base-two positional number system. Here, instead of digits, we use bits, defined as either a 0 or a 1. The meaning of a binary number is illustrated by

$$
101.011 = 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 + 0 \times 2^{-1} + 1 \times 2^{-2} + 1 \times 2^{-3}.
$$

In both decimal and binary, certain fractions can be represented only by infinitely repeated numerals. In decimal, only fully reduced fractions whose denominators are a product of powers of 2 and 5 do not repeat. For example, we have

$$
\frac{1}{2} = 0.5
$$
,  $\frac{1}{4} = 0.25$ ,  $\frac{1}{5} = 0.2$ ,  $\frac{1}{8} = 0.125$ ;

but we also have

$$
\frac{1}{3} = 0.\overline{3}, \quad \frac{1}{6} = 0.1\overline{6}, \quad \frac{1}{7} = 0.\overline{142857}, \quad \frac{1}{9} = 0.\overline{1},
$$

where we use the bar notation to indicate repeating numerals.

In binary, only fully reduced fractions whose denominators are a product of powers of 2 do not repeat. For example, using our more familiar digits to represent the fraction, the corresponding binary numbers are given by

$$
\frac{1}{2} = 0.1, \quad \frac{1}{4} = 0.01, \quad \frac{1}{8} = 0.001,
$$

and

$$
\frac{1}{3} = 0.\overline{01}, \quad \frac{1}{5} = 0.\overline{0011}, \quad \frac{1}{6} = 0.00\overline{10}, \quad \frac{1}{7} = 0.\overline{001}, \quad \frac{1}{9} = 0.\overline{000111}.
$$

Binary numbers with infinite repeats can not be represented exactly using a finite number of bits and are a potential source of round-off errors.

**1.** Using binary, round the fractions 1/3, 1/5, 1/6, 1/7 and 1/9 to six places after the binary point. Which numbers round down and which numbers round up?

### <span id="page-10-0"></span>**Lecture 2 | Double precision**

[View this lecture on YouTube](https://youtu.be/3RTbnzHtf7Q)

Eight bits make a byte. Most numerical computation is done in double precision, with numbers stored using eight bytes. The format of a double precision number is



where s is the sign bit, e is the biased exponent, and 1.f is the significand. Here, *e* is a whole number written in binary, 1023 is in decimal, and *f* is the fractional part of the binary number following a binary point.

The 64 bits of a double precision number are distributed so that the sign uses one bit, the exponent uses 11 bits (in decimal,  $0 \le e \le 2047$ ), and the significand uses 52 bits. The distribution of bits reconciles two conflicting needs: that the numbers should range from the very large to the very small, and that numbers should be closely spaced.

Both the largest and smallest exponents are reserved. When  $e$  is all ones ( $e = 2047$ ) in decimal),  $f = 0$  is used to represent infinity (written in MATLAB as Inf) and  $f \neq 0$  is used to represent 'not a number' (written in MATLAB as NaN). NaN typically results from a 0/0, ∞/∞ or ∞ − ∞ operation. When *e* is all zeros, the double precision representation changes from 1.*f* to 0.*f*, allowing these denormal numbers to gracefully underflow towards zero. The largest positive double precision number is  $\text{realmax} = 1.7977e+308$ , and the smallest positive normal number is  $r = 2.2251e-308$ .

Another important number is called machine epsilon (written in Matlab as eps) and is defined as the distance between one and the next largest machine number. If  $0 \le \delta$ eps/2, then  $1 + \delta = 1$  in computer arithmetic. And since

$$
x + y = x(1 + y/x),
$$

when  $y/x < \text{eps}/2$ , then  $x + y = x$ . In double precision, machine epsilon is equal to eps = 2.2204e-16. Note that the spacing between numbers is uniform between powers of 2, but changes by a factor of two with each additional power of two. For example, the spacing of numbers between one and two is  $\epsilon_{\rm PS}$ , between two and four is  $2*\epsilon_{\rm PS}$ , between four and eight is  $4 \star eps$ , and so on.

We have already learned that not all rational numbers can be represented exactly in double precision. For example, the numbers 1/3 and 1/5 are inexact. In most computations this doesn't matter. But one needs to be aware that a calculation using inexact numbers that should result in an integer (such as zero) may not because of roundoff errors.

**1.** Determine the double precision formats of the numbers 1, 1/2 and 1/3.

**2.** Using the format of a double precision number, determine the largest machine number realmax.

**3.** Using the format of a double precision number, determine the smallest positive normal machine number realmin.

**4.** Using the format of a double precision number, determine the distance between one and the next largest machine number.

### <span id="page-12-0"></span>**Lecture 3 |** Matlab **as a calculator**

[View this lecture on YouTube](https://youtu.be/Z2b5sAq-inQ)

The basic arithmetic operators are the usual  $+$ ,  $-$ ,  $*$ ,  $/$ ,  $\hat{ }$ . The single quote  $\hat{ }$  transposes a matrix. The order of operations follow the standard mathematical rules, and parenthesis can be used. Common mathematical functions are sin, cos, tan, exp, log. The base-10 logarithm is log10. The square-root is sqrt and the absolute value is abs. There is a constant  $pi$ i, though one will need to define  $e = exp(1)$ . Complex numbers use the imaginary unit i or j. Some examples are

```
>> log(exp(1))
ans =
    1
\gg (sqrt(5)+1)/2
ans =
     1.6180
\gt
```
MATLAB also has two special numbers called  $Inf$  (for infinity) and NaN (for not-a-number). These special numbers commonly occur when your program either has a bug or some other run-time error. Simple occurrences of Inf and NaN are

```
>> 1/0
ans =
     Inf
>> 0/0
ans =NaN
\rightarrow
```
A semicolon at the end of a command suppresses output. Two commands can also be placed on one line, using either a comma or a semicolon.:

```
>> x=0; y=1;>> x, yx =0
y =1
>>
```
1. Use MATLAB to compute the following expressions.

a) 
$$
\frac{\sqrt{5}-1}{2}
$$
  
\nb)  $\frac{\left(\frac{\sqrt{5}+1}{2}\right)^{10}-\left(\frac{\sqrt{5}-1}{2}\right)^{10}}{\sqrt{5}}$   
\nc)  $\frac{2^5}{2^5-1}$   
\nd)  $\left(1-\frac{1}{2^5}\right)^{-1}$   
\ne)  $e^3$   
\nf)  $\ln(e^3)$   
\ng)  $\sin(\pi/6)$   
\nh)  $\cos(\pi/6)$   
\ni)  $\sin^2(\pi/6) + \cos^2(\pi/6)$ 

### <span id="page-14-0"></span>**Lecture 4 | Scripts and functions**

[View this lecture on YouTube](https://youtu.be/LY0IsQvEqhE)

For a quick solution to a problem, one can type directly into the Command Window. But when you want to type a longer sequence of commands, or write a program that will be run several times, you are better off writing either a script or a function. These saved programs are called m-files because of their .m extension.

A script is essentially a sequence of commands that can be typed directly into the Command Window, but are saved into a file. When working with a file, editing and debugging become easier, and commands can be quickly added or changed.

You can easily start a new script in the HOME tab, and save it with a name and in a directory of your choice. When developing code, many Matlab programmers like to start a script with the commands

clear all; close all; clc;

These three commands clear the work space, close all figures, and clear the command window. It is a quick way to reset Matlab to prevent previously work from interfering with your new script.

It is often useful for a script to print data during execution. To print to the command window or a file, you will need to learn how to use the MATLAB function  $f$ printf.m.

A function is different than a script because the variables defined in a function are removed from the Workspace after the function ends. You can also create a new function in the HOME tab. When creating a new function, my Editor helpfully shows me the following:

```
function [outputArg1,outputArg2] = untitled(inputArg1,inputArg2)
%UNTITLED Summary of this function goes here
   Detailed explanation goes here
outputArg1 = inputArg1;
outputArg2 = inputArg2;
end
```
A function can optionally have both input arguments and output arguments. The title of the function should match the name of the file it is saved under. Functions may also be written inside of scripts. MATLAB calls these local functions.

In this course, we will mainly write scripts. But we will also make use of many Matlab provided functions. After this course, you may sometimes want to download functions from the Matlab File Exchange. Maybe you will even write a function yourself for others to use.

**1.** The first few Fibonacci numbers  $F_n$  are given by  $1, 1, 2, 3, 5, 8, 13, 21, 34, \ldots$ , where starting from 2, each number is the sum of the preceding two numbers. Write a function  $F = Fibonacci(n)$ 

that returns the *n*th Fibonacci number. You should use Binet's formula, given by

$$
F_n = \frac{\Phi^n - (-\phi)^n}{\sqrt{5}},
$$

where

$$
\Phi = \frac{\sqrt{5} + 1}{2}, \qquad \phi = \frac{\sqrt{5} - 1}{2}.
$$

Because of round-off errors, your function should return a number rounded to the nearest integer using the built-in MATLAB function round.m.

# <span id="page-16-0"></span>**Lecture 5 | Vectors**

#### [View this lecture on YouTube](https://youtu.be/1sFzTZ0mpi0)

Matlab can work directly with both vectors and matrices. We discuss vectors first. A vector in Matlab can either be a row vector or a column vector. An example of a row vector is  $x=[1 \ 2 \ 3]$ , or equivalently,  $x=[1,2,3]$ . On the command line,

```
>> x=[1,2,3]x =1 2 3
\rightarrow
```
An example of a column vector is  $x=[1;2;3]$ . On the command line,

```
>> x=[1;2;3]x =1
     2
     3
\rightarrow
```
Spaces or commas stay in the same row; semicolons move to the next row. Extra spaces are ignored. Row vectors can be changed to column vectors, or vice-a-versa, by the transpose operator ', e.g.,

```
>> x=[1;2;3]'
x =1 2 3
\rightarrow
```
There are some useful functions for constructing vectors. The zero row vector of size n is constructed from  $x = zeros(1, n)$  (and the column vector from  $x = zeros(n, 1)$  or by using the transpose operator). The row vector of size n consisting of all ones is constructed from x=ones(1,n). For example,

```
>> x = zeros(1, 3)x =\begin{matrix} 0 & 0 & 0 \end{matrix}>> x = ones(1,3)x =1 1 1
>>
```
Another useful function constructs a vector of evenly spaced points. To construct a row vector of six evenly space points between zero and one, we can write  $x =$ linspace(0,1,6). If the number of points is not specified, the default is 100. The colon operator also works,

where the spacing is specified rather than the number of points:  $x=0:0.2:1$ . On the command line,

```
>> x =linspace(0,1,6)
x =0 0.2000 0.4000 0.6000 0.8000 1.0000
>> x=0:0.2:1x =0 0.2000 0.4000 0.6000 0.8000 1.0000
>>
```
If  $x$  is a vector, then  $x \times x$  results in an error. If you want to perform element-wise operations, then you need to add a period (or dot) before the operator. There is a dot-multiply operator  $(.*)$ , a dot-divide operator  $(.')$  and a dot-exponentiation operator  $(.')$ . On the command line,

```
>> x=[1 2 3];
>> x.*x
ans =
     1 4 9
>> x./x
ans =
    1 1 1
>> x.^{\wedge}xans =
     1 4 27
\rightarrow
```
All the standard built-in mathematical functions such as cos, sin, exp, log, etc. also work element-wise on vectors (and matrices, in general). For example,

```
>> x=0:pi/2:2*pi;
\gg cos(x)
ans =
    1.0000 0.0000 -1.0000 -0.0000 1.0000
>>
```
**1.** Compute the values of cos *x* and sin *x* for  $x = 0$ ,  $\pi/6$ ,  $\pi/4$ ,  $\pi/3$ , and  $\pi/2$ . Write a script and use fprintf.m to print a nice table to the command window.

# <span id="page-19-0"></span>**Lecture 6 | Line plots**

[View this lecture on YouTube](https://youtu.be/dw4a1SmUw7Q)

The most commonly used graphics function in MATLAB is the line plot. With  $x$  and  $y$ vectors containing the *x*-values and *y*-values of a sequence of points, a line plot connecting these points can be created using  $plot(x,y)$ . For example, we can plot a circle with the following script:

```
theta=linspace(0,2*pi);
x = cos(theta); y = sin(theta);plot(x,y);
```
The figure that is plotted looks like this:



Actually, this plot doesn't look exactly like a circle because the unit length of the x-axis is larger than the unit length of the y-axis. We can obtain a plot that looks like a circle by typing the additional command axis equal. We can also try to make the plot as pretty as possible. I can show you one possibility:

```
axis equal;
axis([-1.1 1.1 -1.1 1.1]);
ax = gca;ax.XTick = [-1 -0.5 0 0.5 1];ax.YTick = [-1 -0.5 0 0.5 1];xlabel('$x$', 'Interpreter', 'latex', 'FontSize',14);
ylabel('$y$', 'Interpreter', 'latex', 'FontSize',14);
title('Plot of a Circle', 'Interpreter', 'latex', 'FontSize',16);
```
Here, we use axis ([xmin xmax ymin ymax]) to set the limits on the axis. The command ax = gca returns what MATLAB calls the handle to the current axes for the current figure. Its use here allows us to specify the location of the tick marks on the *x*- and *y*-axes. To label the axes, I use the xlabel and ylabel commands, and rather than use the default font (as in  $x$ label('x')), I prefer to write mathematical symbols using latex, and to employ a larger font. The resulting graph now looks like



Matlab can also annotate the graph with text, plot unconnected symbols such as points, open circles, crosses, etc., and use various line types and colors. You will need to frequently consult the plot function in the Help Documentation to construct a graph to your liking. You will also need to use the hold on command if you want to add additional points or curves to an already drawn plot.

**1.** Generate the coordinates for a logarithmic spiral, with

$$
x = e^{k\theta} \cos \theta, \qquad y = e^{k\theta} \sin \theta.
$$

Choose  $k = 0.05$  and plot  $\theta$  values from  $-10\pi$  to  $10\pi$  using 2000 evenly spaced points.

**2.** Generate the coordinates of a leminscate, with

$$
x = \pm \cos \theta \sqrt{2 \cos(2\theta)}, \qquad y = \sin \theta \sqrt{2 \cos(2\theta)}.
$$

Plot  $\theta$  values from  $-\pi/4$  to  $\pi/4$  using 1000 evenly spaced points.

# <span id="page-22-0"></span>**Lecture 7 | Matrices**

[View this lecture on YouTube](https://youtu.be/hpQpsL0_WA8)

The name MATLAB stands for Matrix Laboratory, and was created to work with matrices. One can build matrices element-by-element:

```
>> A=[1 2 3; 4 5 6; 7 8 9]
A =1 2 3
   4 5 6
   7 8 9
\rightarrow
```
One can then access the matrix elements of A directly. For example, the element in row m, column n is accessed as  $A(m, n)$ . Submatrices of A can also be accessed. The submatrix consisting of rows i, j and columns l, m can be accessed as  $A([i \ j], [l \ m])$ . For example,

```
>> A(2, 3)ans =6
>> A([1 2], [2 3])ans =2 3
    5 6
>>
```
The colon operator can also be used to good effect. The first row of matrix A can be accessed using  $A(1,:)$ , and the first column by  $A(:,1)$ :

```
>> A=[1 2 3; 4 5 6; 7 8 9];>> A(1, :)ans =1 2 3
>> A(:, 1)
ans =
     1
     4
     7
>>
```
A matrix is actually stored in memory as a single array consisting of concatenated columns of the matrix. To view a matrix as a column vector, we can use  $A():$ . MATLAB's end can also used to denote the last index of a row or column, as in

```
>> A(1, 2:end)ans =
```
2 3

 $\gt$ 

One can also directly reassign matrix elements. For example,

```
>> A(1,1)=10A =10 2 3
    4 5 6
    7 8 9
\rightarrow
```
Matlab functions can also be used to build matrices. The m-by-n zero matrix is constructed from zeros(m,n) and the m-by-n ones matrix from ones(m,n):

```
\gg zeros(3, 2)ans =
    0 0
    0 0
    0 0
> ones (2, 3)ans =
    1 1 1
    1 1 1
>>
```
The n-by-n identity matrix is constructed from  $eye(n)$ , and a diagonal matrix is constructed by specifying a vector for the diagonal elements, as in diag( $[a \ b \ c]$ ). On the command line,

```
>> eye(2)
ans =
    1 0
    0 1
>> diag([1 2])
ans =1 0
    0 2
\gt
```
The MATLAB function  $diag$ .m can also be used to construct banded matrices. For example, the 3-by-3 tridiagonal matrix with 2's on the diagonal and -1's directly above and below the diagonal can be constructed using

```
\Rightarrow diag([2 2 2]) + diag([-1 -1],-1) + diag([-1 -1],1)
ans =
     2 -1 0-1 2 -10 -1 2
>>
```
A more general construction for an n-by-n matrix (here,  $n = 4$ ) can be done using the ones.m function:

```
>> n=4;\Rightarrow diag(2*ones(1,n)) + diag(-ones(1,n-1),-1) + diag(-ones(1,n-1),1)
ans =
     2 \t -1 \t 0 \t 0-1 2 -1 0
     0 -1 2 -10 0 -1 2
```

```
\gt
```
Matrix multiplication uses the \* symbol. For example, we can multiply two two-by-two matrices:

```
>> A=[1 \ 2;3 \ 4]; B=[-3 \ -4;4 \ 5]>> A*B
ans =
      5 6
      7 8
>>
```
We cannot matrix multiply two row vectors, because an m-by-n matrix can only be right multiplied by an n-by-p matrix (inner matrix dimensions must agree):

```
>> x=[1 2 3];
>> x*x
Error using *Inner matrix dimensions must agree.
```
>>

We can, however, multiply after taking the transpose of one of the vectors. For example, the scalar product is

```
>> x=[1 2 3];>> x*x'
ans =
      14
\gt
```
**1.** Construct the following matrix in MATLAB. Use the colon operator and the ones.m function:  $\overline{1}$ 

$$
A = \begin{pmatrix} 6 & 9 & 12 & 15 & 18 & 21 \\ 4 & 4 & 4 & 4 & 4 & 4 \\ 2 & 1 & 0 & -1 & -2 & -3 \\ -6 & -4 & -2 & 0 & 2 & 4 \end{pmatrix}
$$

Evaluate the following expressions without using MATLAB. Check your answers with MATLAB.

- *a*) A([1,3],[2,4])
- *b*)  $A(:,[1,4:6])$
- *c*) A([2,3],:)

**2.** Construct matrices of size *n*-by-*n* that have the following characteristics:

- *a*) Matrix A is a diagonal matrix with 1, 2, 3, . . . , *n* along the diagonal.
- *b*) Matrix B has ones along the superdiagonal (one above the diagonal), ones along the subdiagonal (one below the diagonal) and zeros everywhere else.
- *c*) Matrix C has threes on the diagonal, fours on the superdiagonal, ones on two above the diagonal and zeros everywhere else.

To construct these three matrices, write the MATLAB function specified by function  $[A, B, C] = \text{banded\_matrices}(n)$ 

# <span id="page-26-0"></span>**Lecture 8 | Logicals**

[View this lecture on YouTube](https://youtu.be/zUq5TNff5mo)

A logical variable takes the value of 1 for true and 0 for false. Below is a table of logical operators used to create logical variables:



For example,  $p=(0<1)$  will create a logical variable p and assign it the value of 1 because 0<1 is a true statement. Note that "equal to" is a double equal sign to distinguish it from the assignment operator.

Logical variables (with value 0 or 1) can be used in arithmetic expressions. Logical operators can also be used with vectors and matrices. The dimensions on both sides of a logical operator must agree, and the logical operator returns a vector or matrix of the same size after doing an element-by-element comparison. Here is a simple example:

```
>> x=[0 \ 1 \ 2 \ 0]; y=[0 \ 2 \ 4 \ 0];>> x==yans =1 0 0 1
>>
```
Logical operators have lower precedence than arithmetic operators, so for example  $p=(0<1)$ can be written as  $p=0<1$ , and  $(10/2) == 5$  can be written  $10/2 == 5$ . For clarity, you can always use parentheses. Logicals find their main use in conditional statements, which we discuss next.

- **1.** Evaluate the following expressions. Check your answer with Matlab.
	- *a*) 14>15/3
	- *b*) 8/2<5\*3+1>9
	- *c*) 8/(2<5)\*3+(1>9)
	- *d*) 2+4+3~=60/4-1

**2.** Given  $u=[4 -2 1]$  and  $v=[0 2 1]$ , evaluate the following logical expressions. Check your answer with MATLAB.

- *a*) u <= v
- *b*) u==v
- *c*) u<u+v
- *d*) (u<v)+u

# <span id="page-28-0"></span>**Lecture 9 | Conditionals**

[View this lecture on YouTube](https://youtu.be/jEIQqdcrkWQ)

MATLAB supports the if  $\dots$  end structure with its variants if  $\dots$  else  $\dots$  end and if ... elseif ... else ... end. An expression immediately follows the if or elseif keyword. Then follows one or more statements. The statements are executed only when the expression is true. An expression is true when it is nonempty and contains only nonzero elements. An expression can be a logical, integer or real variable or array. Most commonly, the expression contains logical operators and evaluates either as true (1) or false (0).

As an example, let's write a simple MATLAB function to evaluate  $\sin x/x$  with the correct limiting value when  $x = 0$ . Here it is:

```
function y = my\_sinc(x)if x \sim 0y=sin(x)/x;
else
   v=1;
end
```
Instead of writing if  $x \sim = 0$ , some MATLAB programmers might simply write if x because the following statement is executed only when  $x \sim 0$ . On the command line,

```
>> sin(0)/0
ans =NaN
>> my_sinc(0)
ans =1
```
**1.** Write two functions that return the two solutions to the quadratic equation  $ax^2 + bx +$  $c = 0$ . Use something like

function  $[p, q] = quadratic_f$ ormula(a, b, c)

For the first function, assign

$$
p = \frac{-b + \sqrt{b^2 - 4ac}}{2a}, \quad q = \frac{-b - \sqrt{b^2 - 4ac}}{2a}.
$$

Make sure the code works for some trial values of *a*, *b*, and *c*. Then try your code for  $a = 1$ ,  $b = -10^{12}$ ,  $c = 1$ . Note that  $q = 0$ , which is obviously not a correct root. This is an example of round-off error.

For the second function, assign

$$
p = \frac{-b + \sqrt{b^2 - 4ac}}{2a}, \quad q = \frac{2c}{-b + \sqrt{b^2 - 4ac}} \qquad (b < 0),
$$

and

$$
p = \frac{2c}{-b - \sqrt{b^2 - 4ac}}, \quad q = \frac{-b - \sqrt{b^2 - 4ac}}{2a}, \qquad (b \ge 0).
$$

Again, make sure the code works for some trial values of *a*, *b*, and *c*. Now try this code for  $a = 1$ ,  $b = -10^{12}$ ,  $c = 1$  to obtain more reasonable results.

# <span id="page-30-0"></span>**Lecture 10 | Loops**

```
View this lecture on YouTube
```
Most mathematical algorithms require iteration. The workhorse for iterations is the  $for$ ... end loop. The usage is

```
for variable = expression
     statements
```
end

It is common to make use of the colon operator in the expression. For example, suppose that you want to iterate the logistic map  $x_{n+1} = rx_n(1 - x_n)$  a fixed number of times. Part of your script may contain the code snippet

```
ntimes=10000;
for n=1:ntimes
     x = r * x * (1-x);
```
end

 $MATLAB$  also provides a while loop. The usage is

```
while expression
    statements;
end
```
The  $while$  loop continues as long as the expression is true. A  $while$  loop can be useful, for instance, when you want to iterate a sequence of commands until some event happens. For example, you may want to iterate a loop until some variable converges to a value. You can then end the loop when the change in the variable's value from the previous iteration falls below some specified threshold. For example,

```
tol=1.e-08;
error=2*tol;
while error>tol
    xold=x;
    x=x*x*(1-x);error=abs(x-xold)
```
end

MATLAB also has the commands continue and break. The continue command skips all remaining statements in the current loop iteration and proceeds to the top of the loop and the next iteration. The **break** command exits the loop entirely.

Loops are also commonly nested. Make sure that a different name for the control variable is used for each nested loop.

**1.** Write a function that returns the *n*th Fibonacci number  $F_n$  for integer input *n*. Use the recursion relation

 $F_n = F_{n-1} + F_{n-2}$ ;  $F_1 = 1$ ,  $F_2 = 1$ .

Handle zero and negative integers using

$$
F_0 = 0, \qquad F_{-n} = (-1)^{n+1} F_n.
$$

# <span id="page-32-0"></span>**Lecture 11 | Project I: Logistic map (Part A)**

[View this lecture on YouTube](https://youtu.be/mhHrjvmPnFY)

Your project for this week is to compute the bifurcation diagram for the logistic map. First, a little theory. The logistic map is a one-dimensional map, whose general form is given by

<span id="page-32-1"></span>
$$
x_{n+1} = f(x_n). \t\t(11.1)
$$

A solution of a map proceeds by iteration: starting with an initial value *x*0, one generates the sequence  $x_1$ ,  $x_2$ ,  $x_3$ , and so on.

We say that  $x_*$  is a fixed point of the map if  $x_* = f(x_*)$ . The stability of a fixed point can be determined by perturbation. We write  $x_n = x_* + \epsilon_n$ , and  $x_{n+1} = x_* + \epsilon_{n+1}$ , and substitute into [\(11.1\)](#page-32-1). Using a first-order Taylor series expansion, we obtain

$$
x_* + \epsilon_{n+1} = f(x_* + \epsilon_n) = f(x_*) + \epsilon_n f'(x_*) + \dots
$$
  
=  $x_* + \epsilon_n f'(x_*) + \dots$ ,

which for a small enough perturbation yields

$$
|\epsilon_{n+1}/\epsilon_n|=|f'(x_*)|.
$$

The fixed point  $x_*$  is called stable if  $|\epsilon_{n+1}| < |\epsilon_n|$  so that the perturbation decays. Therefore,

$$
x_*
$$
 is 
$$
\begin{cases} \text{ a stable fixed point if} & |f'(x_*)| < 1; \\ \text{ an unstable fixed point if} & |f'(x_*)| > 1. \end{cases}
$$

The logistic map that you will study is given by

$$
x_{n+1} = \mu x_n (1 - x_n),
$$

where you will assume that  $0 < \mu < 4$  and  $0 < x_0 < 1$ .

A fixed point  $x_*$  of the logistic map satisfies the quadratic equation  $x = \mu x(1 - x)$ , which has two solutions given by  $x_* = 0$  and  $x_* = 1 - 1/\mu$ . The stability of the fixed points are determined from the derivative  $f'(x) = \mu(1-2x)$  evaluated at the fixed points, and we find that  $x_* = 0$  is stable for  $0 < \mu < 1$  and  $x_* = 1 - 1/\mu$  is stable for  $1 < \mu < 3$ . For  $\mu > 3$ , there are no longer any stable fixed points and you will reveal the behavior of the map numerically.

**1.** We say that  $x_1$  and  $x_2$  are in the orbit of a period-2 cycle of a one-dimensional map  $f(x)$ if  $x_2 = f(x_1)$  and  $x_1 = f(x_2)$ , and  $x_1 \neq x_2$ . Determine the orbit of a period-2 cycle for the logistic map by solving the equation  $x = f(f(x))$ , with  $f(x) = \mu x(1 - x)$ . You will obtain a fourth-degree polynomial equation. Solve it by factoring out the known roots  $x = 0$  and  $x = 1 - \frac{1}{\mu}$ .

# <span id="page-34-0"></span>**Lecture 12 | Project I: Logistic map (Part B)**

[View this lecture on YouTube](https://youtu.be/wuWdpgYi78A)

The computational engine of your Matlab code is the iteration of the logistic map given by

$$
x_{n+1} = \mu x_n (1 - x_n).
$$

Your goal is to draw a diagram which illustrates the behavior of the iterates of the logistic map as a function of the parameter *µ*. You should discard the earliest iterations corresponding to transient behavior. The parameter  $\mu$  should vary over the parameter range  $2.4 \leq \mu \leq 4.$ 

Your main computation can contain one outer and two inner loops. Here is a reasonable outline:

```
Loop 1 Start at \mu = 2.4 and finish at \mu = 4.
   Set x = x_0.
   Loop 2 Iterate logistic map a fixed number of times (transient).
      Compute x
   Loop 2 (end)
   Loop 3 Iterate logistic map a fixed number of times (data).
      Compute and save x
   Loop 3 (end)
Loop 1 (end)
```
One will need to set some parameters and these can usually be adjusted after viewing a preliminary plot. Parameters include: (1) the resolution in  $\mu$ ; (2) the starting value  $x_0$ ; (3) the number of transient iterations; (4) the number data points for each  $\mu$ . For grading purposes, I will preset these parameters in the assessed code, but feel free to experiment with their values in MATLAB.

There are two possible approaches to graphing the bifurcation diagram. The simplest approach is to plot the iterates of the logistic map as points on a graph. This can serve to illustrate the bifurcation diagram but can not ultimately result in a high resolution image. The plotting of more and more data, rather than resulting in a better image, results in a blackened figure with all fine details obscured.

The best approach is to bin the iteration data in *x*, and to plot the bifurcation diagram by shading individual pixels of the image directly. To bin the data, you will need an additional parameter (5) the resolution in *x*. This portion of code is outside the scope of this exercise and will be provided to the students.

**1.** Compute the bifurcation diagram for the logistic map. The logistic map is given by

$$
x_{n+1} = \mu x_n (1 - x_n).
$$

and the bifurcation diagram illustrates the behavior of the iterates of the map as a function of the parameter  $\mu$ . You are tasked with writing the computational engine for this code. Here is a reasonable outline:

```
Loop 1 Start at \mu = 2.4 and finish at \mu = 4.
   Set x = x_0.
   Loop 2 Iterate logistic map a fixed number of times (transient).
      Compute x
   Loop 2 (end)
   Loop 3 Iterate logistic map a fixed number of times (data).
      Compute and save x
   Loop 3 (end)
Loop 1 (end)
```
## **Week II**

# **Root Finding**

*In this week's lectures, we learn about root finding. We begin by learning basic numerical methods: the bisection method, Newton's method and the secant method. We then learn how fast these methods converge to the root, defining the concept of order of convergence. The order of convergence of Newton's method is determined, and students are asked in the problems to find the order of convergence of the secant method. An interesting computation of fractals from Newton's method is demonstrated using* Matlab*, and we also discuss the* Matlab *functions that can perform root finding.*

*Your programming project will be to write a* Matlab *code using Newton's method to compute the Feigenbaum delta from the bifurcation diagram for the logistic map.*

## **Lecture 13 | Bisection method**

[View this lecture on YouTube](https://youtu.be/mzQFGOvH-mk)

The problem of root finding is to solve  $f(x) = 0$  for a root  $x = r$ . The bisection method is conceptually the simplest method and almost always works. It is, however, the slowest method and because of this should usually be avoided.

The bisection method requires finding values  $x_0$  and  $x_1$  such that  $x_0 < r < x_1$ . We say that  $x_0$  and  $x_1$  bracket the root. With  $f(r) = 0$ , we want  $f(x_0)$  and  $f(x_1)$  to be of opposite sign, so that  $f(x_0)f(x_1) < 0$ . We then assign  $x_2$  to be the midpoint of  $x_0$  and  $x_1$ , that is  $x_2 = (x_1 + x_0)/2$ , which we can write in the form

$$
x_2 = x_1 - \frac{x_1 - x_0}{2}.
$$

Here,  $-(x_1 - x_0)/2$  corrects the previous best value of the root, assumed to be  $x_1$ , and can be used as an estimate of the error. For the next iteration step, the sign of  $f(x_2)$  needs to be determined, and we proceed to compute  $x_3$  using either  $x_0$  and  $x_2$ , or  $x_1$  and  $x_2$ , whichever pair brackets the root. The iteration proceeds in this fashion and is typically stopped when the absolute value of the error estimate is smaller than some required precision.

*Example: Estimate*  $\sqrt{2} = 1.41421...$  *using*  $x_0 = 1$  *and*  $x_1 = 2$ *.* 

Now  $\sqrt{2}$  is the zero of the function  $f(x) = x^2 - 2$ , and  $f(x_0 = 1) = -1$  and  $f(x_1 = 2) = 2$ , so that the two initial guesses bracket the root. We iterate the algorithm. We have

$$
x_2 = (1+2)/2 = 3/2 = 1.5.
$$

Now,  $f(x_2) = 9/4 - 2 = 1/4 > 0$  so that  $x_0$  and  $x_2$  bracket the root. Therefore,

$$
x_3 = (1 + (3/2))/2 = \frac{5}{4} = 1.25.
$$

Now,  $f(x_3) = 25/16 - 2 = -7/16 < 0$  so that  $x_2$  and  $x_3$  bracket the root. Therefore,

$$
x_4 = ((3/2) + (5/4))/2 = \frac{11}{8} = 1.375,
$$

and so on.

**1.** Using the bisection method, estimate  $\sqrt{3} = 1.73205...$  Use  $x_0 = 1$  and  $x_1 = 2$  and iterate to the value of *x*4.

## **Lecture 14 | Newton's method**

[View this lecture on YouTube](https://youtu.be/cf_NK7NlWrs)

Newton's method is the fastest method, but requires analytical computation of the derivative of  $f(x)$ . If the derivative is known and computational speed is required, then this method should be used, although convergence to the desired root is not guaranteed.

Newton's method approximates the function  $y = f(x)$  by the tangent line to the curve at the point  $(x_n, y_n)$ , where  $y_n = f(x_n)$ . The slope of the tangent line is  $f'(x_n)$ , and we have

$$
y - f(x_n) = f'(x_n)(x - x_n).
$$

The next approximation to the root occurs where the line intercepts the *x*-axis. At this point,  $y = 0$ and  $x = x_{n+1}$ . We have  $-f(x_n) = f'(x_n)(x_{n+1}$ *xn*), or



$$
x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}.
$$

Newton's Method requires a guess for  $x_0$ , which should be chosen as close as possible to the root  $x = r$ .

Example: Estimate 
$$
\sqrt{2} = 1.41421...
$$
 using  $x_0 = 1$ .

Again, we solve  $f(x) = 0$ , where  $f(x) = x^2 - 2$ . To implement Newton's Method, we use  $f'(x) = 2x$ . Therefore, Newton's Method is the iteration

$$
x_{n+1} = x_n - \frac{x_n^2 - 2}{2x_n} = \frac{x_n^2 + 2}{2x_n}
$$

.

Choosing an initial guess  $x_0 = 1$ , we have

$$
x_1 = \frac{(1)^2 + 2}{2 \times 1} = \frac{3}{2} = 1.5, \qquad x_2 = \frac{\left(\frac{3}{2}\right)^2 + 2}{2 \times \frac{3}{2}} = \frac{17}{12} = 1.416667,
$$

$$
x_3 = \frac{\left(\frac{17}{12}\right)^2 + 2}{2 \times \left(\frac{17}{12}\right)} = \frac{577}{408} = 1.41426,
$$

and so on. Notice how much faster Newton's method converges compared to the bisection method.

**1.** Using Newton's method, estimate  $\sqrt{3} = 1.73205...$  Use  $x_0 = 1$  and iterate to the value of *x*3.

## **Lecture 15 | Secant method**

[View this lecture on YouTube](https://youtu.be/Zz8AYQ8c5-U)

The secant method is second fastest to Newton's method, and is most often used when it is not possible to take an analytical derivative of the function  $f(x)$ . To derive the secant method, we begin with Newton's method,

$$
x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)},
$$

and substitute an approximate numerical derivative for  $f'(x_n)$ :

$$
f'(x_n) \approx \frac{f(x_n) - f(x_{n-1})}{x_n - x_{n-1}}.
$$

The Secant method requires an initial guess for both  $x_0$  and  $x_1$ . The initial guesses need not bracket the root, but convergence of the method is not guaranteed.

*Estimate*  $\sqrt{2} = 1.41421...$  *using*  $x_0 = 1$  *and*  $x_1 = 2$ *.* 

Again, we solve  $f(x) = 0$ , where  $f(x) = x^2 - 2$ . The Secant Method iterates

$$
x_{n+1} = x_n - \frac{(x_n^2 - 2)(x_n - x_{n-1})}{x_n^2 - x_{n-1}^2}
$$
  
= 
$$
\frac{x_n x_{n-1} + 2}{x_n + x_{n-1}}.
$$

With  $x_0 = 1$  and  $x_1 = 2$ , we have

$$
x_2 = \frac{2 \times 1 + 2}{2 + 1} = \frac{4}{3} = 1.33333, \quad x_3 = \frac{\frac{4}{3} \times 2 + 2}{\frac{4}{3} + 2} = \frac{7}{5} = 1.4,
$$

$$
x_4 = \frac{\frac{7}{5} \times \frac{4}{3} + 2}{\frac{7}{5} + \frac{4}{3}} = \frac{58}{41} = 1.41463,
$$

and so on.

**1.** Using the secant method, estimate  $\sqrt{3} = 1.73205...$  Use  $x_0 = 1$  and  $x_1 = 2$  and iterate to  $x_4$ .

## **Lecture 16 | Order of convergence**

[View this lecture on YouTube](https://youtu.be/dHQVQG8NqR8)

Consider  $f(x) = 0$ , and suppose *r* is a root and  $x<sub>n</sub>$  is the *n*th approximation to this root. Define the error in the *n*th approximation as

$$
\epsilon_n=r-x_n.
$$

If for large *n*, we have the approximate relationship

$$
|\epsilon_{n+1}| = k |\epsilon_n|^p,
$$

with *k* a positive constant and  $p \geq 1$ , we say the root-finding numerical method is of order *p*. Larger values of *p* correspond to faster convergence to the root. The bisection method is of order one: the error is reduced by approximately a factor of 2 with each iteration so that

$$
|\epsilon_{n+1}| = \frac{1}{2} |\epsilon_n|.
$$

Newton's method is the fastest converging root-finding method and is of order two. The secant method is somewhat slower and is of order approximately 1.6. In the next lecture, we show how to derive the order of Newton's method using Taylor series expansions.

**1.** Complete the table below. Assume  $|\epsilon_{n+1}| = 0.5|\epsilon_n|^p$ .



# **Lecture 17 | Convergence of Newton's method**

[View this lecture on YouTube](https://youtu.be/IPP1p0NKdeo)

We derive the rate of convergence of Newton's method. Consider  $f(x) = 0$  with root *r*. Newton's method is given by

$$
x_{n+1} = x_n - \frac{f(x_n)}{f'(x_n)}.
$$

Defining  $\epsilon_n = r - x_n$ , and subtract both sides of Newton's iteration from *r*:

$$
r - x_{n+1} = r - x_n + \frac{f(x_n)}{f'(x_n)}
$$
, which we write as  $\epsilon_{n+1} = \epsilon_n + \frac{f(r - \epsilon_n)}{f'(r - \epsilon_n)}$ .

We Taylor series expand the functions  $f(r - \epsilon_n)$  and  $f'(r - \epsilon_n)$  for small  $\epsilon_n$ , using  $f(r) = 0$ . We have

$$
f(r-\epsilon_n)=-\epsilon_nf'(r)+\frac{1}{2}\epsilon_n^2f''(r)+\ldots, \qquad f'(r-\epsilon_n)=f'(r)-\epsilon_nf''(r)+\ldots
$$

The ellipsis  $(...)$  represents higher-order terms in  $\epsilon$  that we will eventually neglect. We will make use of the Taylor series  $1/(1-\epsilon)=1+\epsilon+\epsilon^2+\ldots$  , which converges for  $|\epsilon| < 1$ . We expand as follows:

$$
\epsilon_{n+1} = \epsilon_n + \frac{f(r - \epsilon_n)}{f'(r - \epsilon_n)} = \epsilon_n + \frac{-\epsilon_n f'(r) + \frac{1}{2} \epsilon_n^2 f''(r) + \dots}{f'(r) - \epsilon_n f''(r) + \dots}
$$
\n
$$
= \epsilon_n + \frac{-\epsilon_n + \frac{1}{2} \epsilon_n^2 \frac{f''(r)}{f'(r)} + \dots}{1 - \epsilon_n \frac{f''(r)}{f'(r)} + \dots} = \epsilon_n + \left(-\epsilon_n + \frac{1}{2} \epsilon_n^2 \frac{f''(r)}{f'(r)} + \dots\right) \left(1 + \epsilon_n \frac{f''(r)}{f'(r)} + \dots\right)
$$
\n
$$
= \epsilon_n + \left(-\epsilon_n + \epsilon_n^2 \left(\frac{1}{2} \frac{f''(r)}{f'(r)} - \frac{f''(r)}{f'(r)}\right) + \dots\right) = -\frac{1}{2} \frac{f''(r)}{f'(r)} \epsilon_n^2 + \dots
$$

We have shown that to leading order in *e*, we have

$$
|\epsilon_{n+1}| = k |\epsilon_n|^2
$$
, with  $k = \frac{1}{2} \left| \frac{f''(r)}{f'(r)} \right|$ ,

an expression that is valid as long as  $f'(r) \neq 0$  (a so-called simple root). Because the error decreases with each iteration as the square of the previous error, we say that Newton's method is of order two.

**1.** Take the following steps to determine the rate of convergence of the secant method, given by

$$
x_{n+1} = x_n - \frac{(x_n - x_{n-1})f(x_n)}{f(x_n) - f(x_{n-1})}.
$$

*a*) Let  $\epsilon_n = r - x_n$ , and subtract both sides of the secant method from *r* to obtain

$$
\epsilon_{n+1} = \epsilon_n + \frac{(\epsilon_{n-1} - \epsilon_n)f(r - \epsilon_n)}{f(r - \epsilon_n) - f(r - \epsilon_{n-1})}.
$$

*b*) Taylor series expand  $f(r - \epsilon_n)$  and  $f(r - \epsilon_{n-1})$  for small  $\epsilon$  using  $f(r) = 0$ . Obtain

$$
\epsilon_{n+1} = \epsilon_n + \frac{-\epsilon_n f'(r) + \frac{1}{2} \epsilon_n^2 f''(r) + \dots}{f'(r) - \frac{1}{2} (\epsilon_{n-1} + \epsilon_n) f''(r) + \dots}
$$

 $c$ ) For small  $\epsilon$ , use the Taylor series expansion

$$
\frac{1}{1-\epsilon} = 1 + \epsilon + \epsilon^2 + \dots
$$

to obtain

$$
|\epsilon_{n+1}| = \frac{1}{2} \left| \frac{f''(r)}{f'(r)} \right| |\epsilon_n||\epsilon_{n-1}|.
$$

*d*) Try  $|\epsilon_{n+1}| = k|\epsilon_n|^p$  and  $|\epsilon_n| = k|\epsilon_{n-1}|^p$  to obtain the equation  $p^2 = p + 1$ . Determine *p*.

# **Lecture 18 | Fractals from Newton's method**

#### [View this lecture on YouTube](https://youtu.be/vuL87WfRIQE)

The three cube roots of unity are the solutions of  $z^3 = 1$ . We can solve this equation using Euler's formula,  $exp(iθ) = cos θ + i sin θ$ . We write  $z<sup>3</sup> = exp(i2πn)$ , with *n* an integer, and take the cube root to find the solutions  $z = \exp(i2\pi n/3)$ . The three unique complex roots correspond to  $n = 0, 1$  and 2, and using the trig values of special angles, we find √ √

$$
r_1 = 1
$$
,  $r_2 = -\frac{1}{2} + \frac{\sqrt{3}}{2}i$ ,  $r_3 = -\frac{1}{2} - \frac{\sqrt{3}}{2}i$ .

These roots can be located on the unit circle in the complex plane, and are shown below.



We now define the root-finding problem to be  $f(z) = 0$ , where  $f(z) = z^3 - 1$ , with derivative  $f'(z) = 3z^2$ . Newton's method for determining the three complex roots of unity is given by the iteration

$$
z_{n+1} = z_n - \frac{f(z)}{f'(z)}.
$$

Here, we want to determine which initial values of  $z_0$  in the complex plane converge to which of the three cube roots of unity. So if for a given  $z_0$ , the iteration converges to  $r_1$ , say, we will color red the point  $z_0$  in the complex plane; if it converges to  $r_2$ , we will color  $z_0$  green; and if to  $r_3$ , blue.

In the next lecture, we show you how to construct a Matlab program that grids up a rectangle in the complex plane, and determines the convergence of each grid point *z*<sub>0</sub>. By coloring these grid points, we will compute a beautiful fractal.

**1.** What are the four fourth roots of unity?

# **Lecture 19 | Coding the Newton fractal**

[View this lecture on YouTube](https://youtu.be/_FrpXPbP-zk)

Our code first defines the function  $f(z)$  and its derivative to be used in Newton's method, and defines the three cube roots of unity:

```
f = \theta(z) z.^3-1; fp = \theta(z) 3*z.^2;
root1 = 1; root2 = -1/2 + 1i*sqrt(3)/2; root3 = -1/2 - 1i*sqrt(3)/2;
```
We let  $z = x + iy$ . The complex plane is represented as a two-dimensional grid and we define nx and ny to be the number of grid points in the *x*- and *y*-directions, respectively. We further define xmin and xmax to be the minimum and maximum values of *x*, and similarly for ymin and ymax. Appropriate values were determined by numerical experimentation.

```
nx=2000; ny=2000;
xmin=-2; xmax=2; ymin=-2; ymax=2;
```
The grid in *x* and *y* are defined using linspace.

x=linspace(xmin,xmax,nx); y=linspace(ymin,ymax,ny);

We then use meshgrid to construct the two-dimensional grid. Suppose that the *x*-*y* grid is defined by  $x=[x1 \ x2 \ x3]$  and  $y=[y1 \ y2 \ y3]$ . Then  $[x, y]$  =meshgrid( $x, y$ ) results in the matrices

$$
X = \begin{bmatrix} x1 & x2 & x3 \\ x1 & x2 & x3 \\ x1 & x2 & x3 \end{bmatrix}, \qquad Y = \begin{bmatrix} y1 & y1 & y1 \\ y2 & y2 & y2 \\ y3 & y3 & y3 \end{bmatrix},
$$

and we can grid the complex plane using

```
[X, Y]=meshgrid(x,y);
Z = X + 1i *Y;
```
We now iterate Newton's method nit times. These lines constitute the computational engine of the code.

```
nit=40;
for n=1:nit
   Z = Z - f(Z)./ fp(Z);
end
```
We next test to see which roots have converged. We use the logical variables  $z_1$ ,  $z_2$  and z<sub>3</sub> to mark the grid points that converge to one of the three roots. Grid points that have not converged are marked by  $z_4$ , and our convergence criteria is set by the variable eps. The function abs returns the modulus of a complex number.

```
eps=0.001;
Z1 = abs(Z-root1) < eps; Z2 = abs(Z-root2) < eps;Z3 = abs(Z-root3) < eps; Z4 = -(Z1+Z2+Z3);
```
Finally, we draw the fractal. The appropriate graphing function to use here is image, which can color pixels directly. We first open a figure and set our desired colormap.

Here, our map will be a four-by-three matrix, where row one of the matrix corresponds with the RGB triplet  $[1 \ 0 \ 0]$  that specifies red; row two of the matrix  $[0 \ 1 \ 0]$  specifies green; row three of the matrix  $[0 \ 0 \ 1]$  specifies blue; and row four of the matrix  $[0 \ 0 \ 0]$ specifies black. The numbers 1-2-3-4 in our image file will then be colored red-green-blueblack.

```
figure;
map = [1 0 0; 0 1 0; 0 0 1; 0 0 0]; colormap(map);
```
We construct the image file by combining the values of our four  $z$  matrices into a single  $z$ matrix containing elements 1, 2, 3, or 4.

 $Z = (Z1 + 2 \times Z2 + 3 \times Z3 + 4 \times Z4)$ ;

The graph is created in our final lines of code. We use the limits of  $x$  and  $y$  to specify the location of our image in the complex plane. One needs to be aware that the function image assumes that the first row of pixels is located at the top of the image. So by default, image inverts the y-axis direction by setting the 'YDir' property to 'reverse.' We need to undo this when plotting data from a computation because the usual convention is for the first row of pixels to be at the bottom of the image. We therefore set the 'YDir' property to 'normal.'

```
image([xmin xmax], [ymin ymax], Z); set(gca,'YDir','normal');
xlabel('$x$', 'Interpreter', 'latex', 'FontSize',14);
ylabel('$y$', 'Interpreter', 'latex', 'FontSize',14);
title('Fractal from $f(z)=z^3-1$', 'Interpreter', 'latex','FontSize', 16)
```


The resulting plot looks like

**1.** Determine the fractal that arises from using Newton's method to compute the four fourth roots of unity.

# **Lecture 20 | Root finding in** MATLAB

[View this lecture on YouTube](https://youtu.be/Bel9Z5INuOg)

Matlab has two commonly used functions for root finding. The first, roots.m, finds the *n* complex roots of an *n*th degree polynomial. This function is called with syntax  $r = \text{roots}(p)$ , where p is a vector containing  $n + 1$  polynomial coefficients, starting with the coefficient of  $x^n$ . For example, to find the roots of the cubic equation  $x^3-3x^2+4x-2=0$ , one types

```
>> p=[1 -3 4 -2];>> r=roots(p)
r =1.0000 + 1.0000i
    1.0000 - 1.0000i
    1.0000 + 0.0000i
\rightarrow
```
The function  $\text{roots}$ .m works in a roundabout way by finding the eigenvalues of a matrix using the MATLAB function  $eig.m.$  You can browse the help file for more details, or type edit roots.m to look at the function code directly.

The second Matlab function, fzero.m, finds a real root of a nonlinear function. The function fzero is typically called as  $r = f$ zero(f,x0), where fzero solves  $f(x) = 0$ , and x0 is either an initial guess for the root, or a two-dimensional vector whose components bracket the root. For fzero to work, the function must change sign at the root. For example, to solve  $x = \exp(-a * x)$ , where *a* is a parameter set equal to one-half, one can write the script

```
f = \theta(x, a) x-exp(-a*x); a=1/2;
x0=1:
r = fzero(\theta(x) f(x,a), x0);
```
Here, we have defined f using an anonymous function. The argument includes x and also parameters. The function fzero is passed the defined function as  $\mathfrak{g}(x) \in \{x, a\}$ .

Sometimes it may be easier to define *f* in a subfunction. Then a similar script would look like

```
a=1/2; x0=1;
r = fzero(\&(x) f(x,a), x0)function y = f(x, a)y=x-exp(-a*x);end
```
**1.** Consider a planet in an elliptical orbit about the sun. Assume the sun is fixed in space and at the origin of the coordinate system, and the planet is located at position  $\mathbf{r}(t) = x(t)\mathbf{i} + y(t)\mathbf{j}$ . Let  $T = 2\pi/\omega$  be the period of the planet's orbit, let *a* and *b* be onehalf the major and minor axes of the ellipse, and let  $e =$ √  $1 - b^2/a^2$  be the eccentricity of the ellipse, with  $0 \le e < 1$ . Then the coordinates of the planet at time *t* are given by

$$
x(t) = a (e - \cos E), \qquad y(t) = b \sin E.
$$

The eccentric anomaly,  $E = E(t)$ , is a solution of Kepler's equation, a transcendental equation given by

$$
E = \omega t + e \sin E.
$$

By solving Kepler's equation (using  $fzero.m$ ), you will compute and plot several different elliptical orbits of a planet.

The closest approach to the sun occurs when  $t = 0$  and  $E = 0$ . To construct your plots, choose units so that the distance of closest approach is one. At the closest approach,  $y = 0$ and  $x = a(e-1)$ , so choose the unit of length such that  $a(e-1) = -1$ . Solving for *a* and then *b* in terms of the dimensionless parameter *e*, we have

$$
a = \frac{1}{1 - e'}, \qquad b = \sqrt{\frac{1 + e}{1 - e}}.
$$

If we also choose units such that the period of an orbit is one, then  $\omega = 2\pi$ . On a single plot, show the four orbits corresponding to  $e = 0$ ,  $1/4$ ,  $1/2$  and  $3/4$ .

# **Lecture 21 | Project II: Feigenbaum delta (Part A)**

[View this lecture on YouTube](https://youtu.be/Od3a6kx3t2k)

The bifurcation diagram for the logistic map,

$$
x_{i+1} = \mu x_i (1 - x_i),
$$

showed a fixed point of the map bifurcating to a period-2 cycle at  $\mu = 3$ , and to a period-4 cycle near  $\mu = 3.45$ . A period-8 cycle can be observed soon after, and in fact, period doubling continues indefinitely, eventually accumulating at a value of  $\mu$  after which the period becomes infinite. This phenomena is called the period-doubling route to chaos, and is observed in many nonlinear systems.

Let  $\mu_n$  be the value of  $\mu$  at which the period-2<sup>n</sup> cycle bifurcates to a period-2<sup>n+1</sup> cycle. The Feigenbaum delta is defined as

$$
\delta=\lim_{n\to\infty}\frac{\mu_{n-1}-\mu_{n-2}}{\mu_n-\mu_{n-1}},
$$

and is a measure of the decreasing width in  $\mu$  of the period- $2^n$  regions as  $n$  increases.

The goal here is to accurately compute the Feigenbaum delta. Let  $N = 2<sup>n</sup>$ , and define the orbit of a period-*N* cycle to be  $x_0, x_1, \ldots, x_{N-1}$ . It can be shown that every period-*N* cycle contains one value of  $\mu$  with  $x = 1/2$  in the orbit. This orbit is superstable: iterations of the logistic map most rapidly converge to the orbit solution. If we define *m<sup>n</sup>* to be the value of  $\mu$  at which  $x_0 = 1/2$ , say, is in the orbit of the period-*N* cycle, then the easiest way to compute the Feigenbaum delta is to replace  $\mu_n$  by  $m_n$  in its definition.

We can determine analytically the first two values of *mn*. For the period-1 cycle, we solve  $x_0 = \mu x_0(1 - x_0)$ ; and taking  $x_0 = 1/2$  yields  $\mu = m_0 = 2$ . For the period-2 cycle, we solve

$$
x_1 = \mu x_0 (1 - x_0), \qquad x_0 = \mu x_1 (1 - x_1);
$$

and taking  $x_0 = 1/2$  reduces to  $\mu^3 - 4\mu^2 + 8 = 0$ . We can make use of MATLAB's roots .m: >> roots([1,-4,0,8])' ans = 3.2361 2.0000 -1.2361

>>

The root  $\mu = 2$  corresponds to the  $m_0$  solution. We could have divided this root out of the cubic equation to obtain  $m_1 = 1 +$  $\frac{1}{\sqrt{2}}$  $5 \approx 3.2361$ . Further values of  $m_n$  will be computed numerically.

- **1.** Determine the value of  $m_1$  as follows:
	- *a*) Show that the period-two fixed-point equations, given by

$$
x_1 = \mu x_0 (1 - x_0), \qquad x_0 = \mu x_1 (1 - x_1),
$$

with  $x_0 = 1/2$  reduces to  $\mu^3 - 4\mu^2 + 8 = 0$ .

*b*) Using long division, determine the quadratic polynomial obtained from

$$
\frac{\mu^3-4\mu^2+8}{\mu-2}.
$$

Show that the positive root of this quadratic is  $m_1 = 1 +$ √ 5.

# **Lecture 22 | Project II: Feigenbaum delta (Part B)**

[View this lecture on YouTube](https://youtu.be/Qbc2AdcHWRo)

We have defined the Feigenbaum delta as

$$
\delta = \lim_{n \to \infty} \frac{m_{n-1} - m_{n-2}}{m_n - m_{n-1}},
$$

where  $m_n$  is the value of  $\mu$  at which  $x_0 = 1/2$  is in the orbit of the period-*N* cycle, with  $N = 2^n$ . Analytically, we found  $m_0 = 2$  and  $m_1 = 1 + \sqrt{5} \approx 3.2361$ .

To determine more values of *mn*, we solve a root-finding problem. For each value of *n*, we determine the value of  $\mu$  such that the orbit of the logistic map starts with  $x_0 = 1/2$ and ends with  $x_N = 1/2$ . In other words, we solve for  $g(\mu) = 0$ , where

$$
g(\mu)=x_N-1/2.
$$

The roots are  $\mu = m_0, m_1, \ldots, m_n$ , and the largest root  $m_n$  will be the sort-after solution.

We use Newton's method. The derivative of  $g(\mu)$  with respect to  $\mu$  is given by  $g'(\mu)$  =  $x_N'$ . From the logistic map and its derivative, we obtain the coupled equations

$$
x_{i+1} = \mu x_i (1 - x_i), \qquad x'_{i+1} = x_i (1 - x_i) + \mu x'_i (1 - 2x_i),
$$

with initial values  $x_0 = 1/2$  and  $x'_0 = 0$ . Newton's method then solves for  $m_n$  by iterating

$$
\mu^{(j+1)} = \mu^{(j)} - \frac{x_N - 1/2}{x'_N},
$$

where both  $x_N$  and  $x'_N$  are obtained by iterating the coupled difference equations  $N$  times.

Newton's method requires an initial guess  $\mu^{(0)}$  for the root  $m_n.$  Using the definition of the Feigenbaum delta, we can approximate *m<sup>n</sup>* by

$$
m_n \approx m_{n-1} + \frac{m_{n-1} - m_{n-2}}{\delta_{n-1}},
$$

where *δn*−<sup>1</sup> is our current best approximation to the Feigenbaum delta, given by

$$
\delta_{n-1} = \frac{m_{n-2} - m_{n-3}}{m_{n-1} - m_{n-2}}.
$$

To start, we will take  $\delta_1 = 5$ , and compute  $\delta_n$  for  $n \geq 2$ . The successive values of  $\delta_n$  should converge to the Feigenbaum delta. In the next lecture, we will outline the construction of a MATLAB code.

# **Lecture 23 | Project II: Feigenbaum delta (Part C)**

[View this lecture on YouTube](https://youtu.be/c7msbbiXtLw)

Your second Matlab project is to compute the Feigenbaum delta to high accuracy. The first 30 decimal places of the Feigenbaum delta are known to be

*δ* = 4.669 201 609 102 990 671 853 203 821 578

Let's see how close you can get! Your computational engine should contain three nested loops. Here is a reasonable outline:

**Loop 1** Start at period-2<sup>*n*</sup> with  $n = 2$ , and increment *n* with each iteration Compute initial guess for  $m_n$  using  $m_{n-1}$ ,  $m_{n-2}$  and  $\delta_{n-1}$ **Loop 2** Iterate Newton's method, either a fixed number of times or until convergence Initialize logistic map **Loop 3** Iterate the logistic map  $2^n$  times Compute *x* and *x* 0 **Loop 3 (end)** One step of Newton's method **Loop 2 (end)** Save  $m_n$  and compute  $\delta_n$ **Loop 1 (end)**

You will find that your code runs fast, and convergence during Newton's method is quick. Unfortunately, precise calculation of the Feigenbaum delta is limited by round-off error occurring with double precision. Grading will be done on the computed values of *δn*.

#### *For ambitious students*:

One method to obtain a higher precision calculation of the Feigenbaum delta using Matlab is to replace all double precision calculations with variable precision arithmetic. The relevant MATLAB functions that you will need are vpa and digits, and you can use the MATLAB help pages to develop your code.

**1.** Compute the Feigenbaum delta from the logistic map. The logistic map is given by

$$
x_{i+1} = \mu x_i (1 - x_i),
$$

and the Feigenbaum delta is defined as

$$
\delta = \lim_{n \to \infty} \delta_n, \quad \text{where} \quad \delta_n = \frac{m_{n-1} - m_{n-2}}{m_n - m_{n-1}},
$$

and where  $m_n$  is the value of  $\mu$  for which  $x_0 = 1/2$  is in the orbit of the period-*N* cycle with  $N = 2^n$ . Here is a reasonable outline:

**Loop 1** Start at period-2<sup>*n*</sup> with  $n = 2$ , and increment *n* with each iteration

Compute initial guess for  $m_n$  using  $m_{n-1}$ ,  $m_{n-2}$  and  $\delta_{n-1}$ 

**Loop 2** Iterate Newton's method, either a fixed number of times or until convergence Initialize logistic map **Loop 3** Iterate the logistic map  $2^n$  times Compute *x* and *x* 0 **Loop 3 (end)**

One step of Newton's method

#### **Loop 2 (end)**

Save  $m_n$  and compute  $\delta_n$ 

#### **Loop 1 (end)**

Grading will be done on the computed values of  $\delta_n$  up to  $n = 11$ . Set  $\delta_1 = 5$ .

## **Week III**

## **Matrix Algebra**

*In this week's lectures, we learn about computational matrix algebra. When performing Gaussian elimination, round-off errors can ruin the computation and must be handled using the method of partial pivoting, where row interchanges are performed before each elimination step. The LU decomposition algorithm then includes permutation matrices. We also present the concept of operation counts, and teach the big-Oh notation for predicting the increase in computational time with large problem sizes. We show how to count operations for Gaussian elimination and forward and backward substitution. The power method for computing the largest eigenvalue and associated eigenvector of a matrix is also explained. Finally, we show how to use Gaussian elimination to solve a system of nonlinear differential equations using Newton's method.*

*Your programming project will be to write a* Matlab *code that applies Newton's method to a system of nonlinear equations.*

# **Lecture 24 | Gaussian elimination without pivoting**

#### [View this lecture on YouTube](https://youtu.be/S5dL9xOj0lU)

Gaussian elimination is the standard algorithm to solve a system of linear equations. However, a straightforward numerical implementation of Gaussian elimination without row or column interchanges (pivoting) can result in large errors because of the computer representation of real numbers. We illustrate here a clear but extreme example of what can go wrong. Be aware, though, that the accumulation of smaller round-off errors in large systems of linear equations can also lead to inaccurate results.

Define  $\epsilon = 2^{-52} \approx 2.2 \times 10^{-16}$  to be machine epsilon, that is, the distance between one and the next largest machine number. In computer arithmetic, recall that  $2 + \epsilon = 2$ , and  $4 - \epsilon = 4.$ 

Now, consider the system of equations given by

$$
\epsilon x_1 + 2x_2 = 4, \qquad x_1 - x_2 = 1.
$$

Since  $\epsilon$  < < 1, a close approximation to the solution can be found to be  $x_2 = 2$  and  $x_1 = 3$ .

Gaussian elimination on the augmented matrix without row exchanges (called pivoting) follows

$$
\begin{pmatrix} \epsilon & 2 & 4 \\ 1 & -1 & 1 \end{pmatrix} \rightarrow \begin{pmatrix} \epsilon & 2 & 4 \\ 0 & -\frac{2}{\epsilon} - 1 & -\frac{4}{\epsilon} + 1 \end{pmatrix}.
$$

Back substitution, assuming computer arithmetic, then results in

$$
x_2 = \frac{-\frac{4}{\epsilon} + 1}{-\frac{2}{\epsilon} - 1} = \frac{4 - \epsilon}{2 + \epsilon} = 2, \qquad x_1 = \frac{4 - 2x_2}{\epsilon} = \frac{4 - 4}{\epsilon} = 0.
$$

The result for  $x_1$  is wrong (zero instead of three). We can see where this error comes from if we solve the problem exactly:

$$
x_1 = \frac{4 - 2x_2}{\epsilon} = \frac{4 - 2(\frac{4 - \epsilon}{2 + \epsilon})}{\epsilon} = \frac{6\epsilon}{\epsilon(2 + \epsilon)} = \frac{6}{2 + \epsilon} \approx 3.
$$

In the exact calculation,  $x_2$  is slightly less than two, and a subtraction between two nearly identical numbers leads to the correct answer. But the roundoff errors, here given by  $4 - \epsilon = 4$  and  $2 + \epsilon = 2$ , causes the calculation to go completely awry.

In the next lecture, we will learn how to fix this problem.

**1.** Consider again the system of equations given by

$$
\epsilon x_1 + 2x_2 = 4, \qquad x_1 - x_2 = 1.
$$

The solution of these equations using Gaussian elimination without pivoting was found to be

$$
x_2 = \frac{-\frac{4}{\epsilon} + 1}{-\frac{2}{\epsilon} - 1}, \qquad x_1 = \frac{4 - 2x_2}{\epsilon}.
$$

Compute the value of  $x_2$  and  $x_1$  using MATLAB as a calculator. Now, repeat this calculation for the system of equations given by

$$
2\varepsilon x_1 + 2x_2 = 4, \qquad x_1 - x_2 = 1.
$$

# **Lecture 25 | Gaussian elimination with partial pivoting**

[View this lecture on YouTube](https://youtu.be/WU3tSs5q1MM)

Gaussian elimination with partial pivoting is the standard method for solving a linear system of equations on the computer. We consider again the system of equations given by

$$
\epsilon x_1 + 2x_2 = 4, \qquad x_1 - x_2 = 1,
$$

with a close approximation to the solution given by  $x_2 = 2$  and  $x_1 = 3$ .

The idea of partial pivoting is to perform row interchanges before each elimination step so that the pivot has the largest absolute value among all the rows. The augmented matrix of our system of equations is given by

$$
\begin{pmatrix} \epsilon & 2 & 4 \\ 1 & -1 & 1 \end{pmatrix}.
$$

To find the best pivot for the first column, we find the row with the largest absolute value in the pivot position. Here we only have two rows, but in general you would compare potential pivots in all *n* rows. Since the 1 in the second row has larger magnitude than the  $\epsilon$  in the first row, we interchange rows. We then proceed with the elimination step:

$$
\begin{pmatrix} 1 & -1 & 1 \ \epsilon & 2 & 4 \end{pmatrix} \rightarrow \begin{pmatrix} 1 & -1 & 1 \ 0 & 2+\epsilon & 4-\epsilon \end{pmatrix} \rightarrow \begin{pmatrix} 1 & -1 & 1 \ 0 & 2 & 4 \end{pmatrix},
$$

where the last step results from round-off error. Now, back substitution results in

$$
x_2 = \frac{4}{2} = 2, \qquad x_1 = 1 + x_2 = 3,
$$

a very close approximation to the exact solution.

This method is called *partial pivoting*. *Full pivoting* also implements column exchanges, but this unnecessarily slows down the computation and is usually not implemented.

**1.** Consider again the system of equations given by

$$
\epsilon x_1 + 2x_2 = 4, \qquad x_1 - x_2 = 1.
$$

The solution of these equations using Gaussian elimination with partial pivoting is found to be

$$
x_2 = \frac{4-\epsilon}{2+\epsilon}, \qquad x_1 = 1+x_2.
$$

Compute the value of  $x_2$  and  $x_1$  using MATLAB as a calculator. Now, repeat this calculation for the system of equations given by

$$
2\varepsilon x_1 + 2x_2 = 4, \qquad x_1 - x_2 = 1.
$$

# **Lecture 26 | LU decomposition with partial pivoting**

[View this lecture on YouTube](https://youtu.be/WgaFycuL8z0)

The LU decomposition can be generalized to include partial pivoting. For example, consider

$$
A = \begin{pmatrix} -2 & 2 & -1 \\ 6 & -6 & 7 \\ 3 & -8 & 4 \end{pmatrix}.
$$

We interchange rows to obtain the largest pivot; that is,

$$
A \rightarrow \begin{pmatrix} 6 & -6 & 7 \\ -2 & 2 & -1 \\ 3 & -8 & 4 \end{pmatrix} = P_{12}A, \quad \text{where } P_{12} = \begin{pmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{pmatrix}
$$

is the permutation matrix that interchanges rows one and two. The elimination step is then

$$
P_{12}A \rightarrow \begin{pmatrix} 6 & -6 & 7 \\ 0 & 0 & 4/3 \\ 0 & -5 & 1/2 \end{pmatrix} = M_2M_1P_{12}A, \quad \text{where } M_2M_1 = \begin{pmatrix} 1 & 0 & 0 \\ 1/3 & 1 & 0 \\ -1/2 & 0 & 1 \end{pmatrix}.
$$

The final step requires one more row interchange:

$$
M_2M_1P_{12}A \rightarrow \begin{pmatrix} 6 & -6 & 7 \\ 0 & -5 & 1/2 \\ 0 & 0 & 4/3 \end{pmatrix} = P_{23}M_2M_1P_{12}A = U.
$$

Since the permutation matrices corresponding to a single row exchange are their own inverses, we can write our result as

$$
(P_{23}M_2M_1P_{23})P_{23}P_{12}A = U.
$$

Multiplication of  $M_2M_1$  on the left by  $P_{23}$  interchanges rows two and three, while multiplication on the right by  $P_{23}$  interchanges columns two and three. That is,

$$
P_{23}\begin{pmatrix}1&0&0\\1/3&1&0\\-1/2&0&1\end{pmatrix}P_{23}=\begin{pmatrix}1&0&0\\-1/2&0&1\\1/3&1&0\end{pmatrix}P_{23}=\begin{pmatrix}1&0&0\\-1/2&1&0\\1/3&0&1\end{pmatrix}.
$$

The net result on  $M_2M_1$  is an interchange of the nondiagonal elements 1/3 and  $-1/2$ . We can then multiply by the inverse of  $(P_{23}M_2M_1P_{23})$  to obtain

$$
P_{23}P_{12}A = (P_{23}M_2M_1P_{23})^{-1}U = LU,
$$

where

$$
L = \begin{pmatrix} 1 & 0 & 0 \\ 1/2 & 1 & 0 \\ -1/3 & 0 & 1 \end{pmatrix}, \qquad U = \begin{pmatrix} 6 & -6 & 7 \\ 0 & -5 & 1/2 \\ 0 & 0 & 4/3 \end{pmatrix}
$$

Instead of L, MATLAB will write this as

$$
A = (P_{12}P_{23}L)U,
$$

where the permutation matrices times the lower triangular matrix is called a "psychologically lower triangular matrix," and is given by

$$
P_{12}P_{23}L = \begin{pmatrix} -1/3 & 0 & 1 \\ 1 & 0 & 0 \\ 1/2 & 1 & 0 \end{pmatrix}.
$$

**1.** Let

$$
A = \begin{pmatrix} -3 & 2 & -1 \\ 6 & -6 & 7 \\ 3 & -4 & 4 \end{pmatrix}.
$$

Using Gaussian elimination with partial pivoting, find the (PL)U decomposition of A, where U is an upper triangular matrix and (PL) is a psychologically lower triangular matrix.

# **Lecture 27 | Operation counts**

[View this lecture on YouTube](https://youtu.be/3yLy9vUihTc)

To estimate the computational time of an algorithm, one counts the number of operations required (multiplications, additions, etc.). Usually, one is interested in how an algorithm scales with the size of the problem. One can easily time a small problem, and then estimate the computational time required for a much larger problem.

For example, suppose one is multiplying two full  $n \times n$  matrices. The calculation of each element requires *n* multiplications and *n* − 1 additions. There are *n*<sup>2</sup> elements to compute so that the total operation counts are  $n^3$  multiplications and  $n^2(n-1)$  additions.

For large *n*, we might want to know how much longer the computation will take if *n* is doubled. What matters most is the fastest-growing, leading-order term in the operation count. In this matrix multiplication example, the leading-order term grows like  $n^3$ , and we say that the algorithm scales like  $O(n^3)$ , which is read as "big Oh of n cubed." When using the big-Oh notation, we drop both lower-order terms and constant multipliers.

The big-Oh notation tells us how the computational time of an algorithm scales. For example, suppose that the multiplication of two large  $n \times n$  matrices took a computational time of  $T_n$  seconds. With the known operation count going like  $O(n^3)$ , we can write  $T_n = kn^3$  for some constant *k*. To determine how much longer the multiplication of two  $2n \times 2n$  matrices will take, we write

$$
T_{2n} = k(2n)^3 = 8kn^3 = 8T_n,
$$

so that doubling the size of the matrix is expected to increase the computational time by a factor of  $2^3 = 8$ .

Running Matlab on my computer, the multiplication of two  $4096 \times 4096$  matrices took about 1.1 sec. The multiplication of two  $8192 \times 8192$  matrices took about 7.9 sec, which is a bit more than 7 times longer. Timing of code in Matlab can be found using the built-in stopwatch functions tic and toc.

**1.** A genetic model of recombination is solved using a computational algorithm that scales like  $O(3^L)$ , where *L* is the number of loci modeled. If it takes 10 sec to compute recombination when  $L = 15$ , estimate how long it takes to compute recombination when  $L = 16$ .

# **Lecture 28 | Operation counts for Gaussian elimination**

[View this lecture on YouTube](https://youtu.be/WLqtv-fr0R4)

When counting the number of operations required for an algorithm, three well-known summation formulas will come in handy. They are

$$
\sum_{k=1}^{n} 1 = n, \qquad \sum_{k=1}^{n} k = \frac{1}{2}n(n+1), \qquad \sum_{k=1}^{n} k^{2} = \frac{1}{6}n(2n+1)(n+1).
$$

To determine the operation counts for Gaussian elimination used to solve a system of *n* equations and *n* unknowns, consider an elimination step with the pivot on the diagonal in the *i*th row and column. There are both  $n - i$  rows below the pivot and  $n - i$  columns to the right of the pivot. To perform elimination on one row, each matrix element to the right of the pivot must be multiplied by a factor and added to the row underneath. This must be done for all the rows. There are therefore  $(n - i)(n - i)$  multiplication-additions to be done for a pivot . I will count a multiplication-addition as one operation.

To find the total operation count, we need to perform elimination using *n* − 1 pivots, so that

op. counts 
$$
=
$$
  $\sum_{i=1}^{n-1} (n-i)^2 = (n-1)^2 + (n-2)^2 + ... (1)^2$   
 $=$   $\sum_{i=1}^{n-1} i^2 = \frac{1}{6} (n-1)(2n-1)n.$ 

The leading-order term is  $n^3/3$ , and we say that Gaussian elimination scales like  $O(n^3)$ . Since LU decomposition with partial pivoting is essentially Gaussian elimination, it will have the same scaling. Double the number of equations and unknowns, and LU decomposition will take approximately eight times longer.

**1.** Derive the following summation identities:

a)  

$$
\sum_{k=1}^{n} 1 = n;
$$
  
b)  

$$
\sum_{k=1}^{n} k = \frac{1}{2}n(n+1);
$$

*c*)

$$
\sum_{k=1}^{n} k^2 = \frac{1}{6} n(2n+1)(n+1).
$$

# **Lecture 29 | Operation counts for forward and backward substitution**

[View this lecture on YouTube](https://youtu.be/_IbtdLqeisg)

Once the LU decomposition of a matrix A is known, the solution of  $Ax = b$  can proceed by forward and backward substitution. Here, we compute the operation counts for backward substitution. Forward substitution will be similar. The matrix equation to be solved is of the form

$$
\begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1(n-1)} & a_{1n} \\ 0 & a_{22} & \cdots & a_{2(n-1)} & a_{2n} \\ \vdots & \vdots & \ddots & \vdots & \vdots & \vdots \\ 0 & 0 & \cdots & a_{(n-1)(n-1)} & a_{(n-1)n} \\ 0 & 0 & \cdots & 0 & a_{nn} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_{n-1} \\ x_n \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_{n-1} \\ b_n \end{pmatrix}.
$$

The solution for  $x_i$  is found after solving for  $x_j$  with  $j > i$ . The explicit solution is given by

$$
x_i = \frac{1}{a_{ii}} \left( b_i - \sum_{j=i+1}^n a_{ij} x_j \right).
$$

The solution for  $x_i$  requires  $n - i + 1$  multiplications and  $n - i$  additions, and we need to solve for *n* values of *x<sup>i</sup>* . Counting just multiplications, we have

op. counts 
$$
=\sum_{i=1}^{n} n - i + 1 = n + (n - 1) + \dots + 1
$$
  
 $=\sum_{i=1}^{n} i = \frac{1}{2}n(n + 1).$ 

The leading-order term is *n* <sup>2</sup>/2, and forward substitution has a similar scaling. The operation counts for a forward and backward substitution, therefore, scales like  $O(n^2)$ .

For *n* large, *n*<sup>2</sup> grows much slower than *n*<sup>3</sup>, and a forward and backward substitution should be substantially faster than Gaussian elimination. If the matrix is fixed and the right-hand side keeps changing — for example, in a time-dependent computation where the matrix is independent of time but the right-hand side changes over time — a fast computation would first find the LU decomposition of the matrix, and then solve the recurring systems of equations by forward and backward substitution.
**1.** Solve the following lower triangular system for  $x_i$  in terms of  $x_j$ ,  $j < i$ :

$$
\begin{pmatrix} a_{11} & 0 & \cdots & 0 & 0 \\ a_{21} & a_{22} & \cdots & 0 & 0 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ a_{(n-1)1} & a_{(n-1)2} & \cdots & a_{(n-1)(n-1)} & 0 \\ a_{n1} & a_{n2} & \cdots & a_{n(n-1)} & a_{nn} \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_{n-1} \\ x_n \end{pmatrix} = \begin{pmatrix} b_1 \\ b_2 \\ \vdots \\ b_{n-1} \\ b_n \end{pmatrix}.
$$

Count the total number of multiplication-additions required for a complete solution. **[Solutions to the Problems](#page-187-0)**

# **Lecture 30 | Eigenvalue power method**

#### [View this lecture on YouTube](https://youtu.be/_PDyi5BVY-E)

The power method is a fast iterative algorithm to find the largest magnitude eigenvalue of a matrix and its associated eigenvector. Suppose that A is an *n*-by-*n* matrix with *n* distinct real eigenvalues, ordered such that  $|\lambda_1| > |\lambda_2| > \cdots > |\lambda_n|$ , and with corresponding linearly independent eigenvectors  $e_1, e_2, \ldots e_n$ . Then any vector  $x_0$  can be written as linear combination of the eigenvectors as

$$
x_0 = c_1e_1 + c_2e_2 + \dots + c_ne_n = \sum_{i=1}^n c_i e_i.
$$

With  $x_1 = Ax_0$ , we have

$$
x_1 = A \sum_{i=1}^n c_i e_i = \sum_{i=1}^n c_i A e_i = \sum_{i=1}^n c_i \lambda_i e_i.
$$

Repeated multiplication by A, with  $x_p = A^p x_0$ , yields

$$
x_p = \sum_{i=1}^n c_i \lambda_i^p e_i = \lambda_1^p \left( c_1 e_1 + \sum_{i=2}^n c_i (\lambda_i / \lambda_1)^p e_i \right).
$$

Since for  $i \geq 2$ ,

$$
\lim_{p\to\infty}(\lambda_i/\lambda_1)^p=0,
$$

a good approximation for large enough values of *p* is

$$
x_p \approx c_1 \lambda_1^p e_1.
$$

To find the dominant eigenvalue, we write

$$
\frac{x_p^T x_{p+1}}{x_p^T x_p} = \frac{\left(c_1 \lambda_1^p e_1^T\right) \left(c_1 \lambda_1^{p+1} e_1\right)}{\left(c_1 \lambda_1^p e_1^T\right) \left(c_1 \lambda_1^p e_1\right)} = \frac{c_1^2 \lambda_1^{2p+1} e_1^T e_1}{c_1^2 \lambda_1^{2p} e_1^T e_1} = \lambda_1.
$$

The corresponding eigenvector is found from  $x_p$ , which can be normalized as you like.

In practice, after many iterations underflow (when  $\lambda_1 < 1$ ) or overflow (when  $\lambda_1 > 1$ ) can occur, and it is best to normalize the vectors  $x_p$  after each iteration. A simple method of normalization is

$$
x_p \to \frac{x_p}{(x_p^Tx_p)^{1/2}}.
$$

**1.** The two largest (in absolute value) eigenvalues of an *n*-by-*n* matrix with real eigenvalues are  $\lambda_1 = 1$  and  $\lambda_2 = 1/2$ . Give a rough estimate of how many iterations of the power method is required for convergence to an error of less than  $10^{-8}$ .

# **Lecture 31 | Eigenvalue power method (example)**

[View this lecture on YouTube](https://youtu.be/nKd0lu3yThg)

*Use the power method to determine the dominant eigenvalue and corresponding eigenvector of the matrix*

$$
A = \begin{pmatrix} 6 & 5 \\ 4 & 5 \end{pmatrix}.
$$

For reference, the eigenvalues and eigenvectors of this matrix are given by

$$
\lambda_1 = 10
$$
,  $v_1 = \begin{pmatrix} 5/4 \\ 1 \end{pmatrix}$ ;  $\lambda_2 = 1$ ,  $v_2 = \begin{pmatrix} 1 \\ -1 \end{pmatrix}$ .

We apply here the power method without normalizing the vectors after each iteration. A computer solution should normalize vectors to avoid overflow or underflow. The assumed initial vector and first iteration is given by

$$
x_0 = \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \qquad x_1 = Ax_0 = \begin{pmatrix} 6 \\ 4 \end{pmatrix}.
$$

Continuing,

$$
x_2 = \begin{pmatrix} 6 & 5 \\ 4 & 5 \end{pmatrix} \begin{pmatrix} 6 \\ 4 \end{pmatrix} = \begin{pmatrix} 56 \\ 44 \end{pmatrix}; \qquad x_3 = \begin{pmatrix} 6 & 5 \\ 4 & 5 \end{pmatrix} \begin{pmatrix} 56 \\ 44 \end{pmatrix} = \begin{pmatrix} 556 \\ 444 \end{pmatrix}.
$$

Two more iterations give

$$
x_4 = \begin{pmatrix} 6 & 5 \\ 4 & 5 \end{pmatrix} \begin{pmatrix} 556 \\ 444 \end{pmatrix} = \begin{pmatrix} 5,556 \\ 4,444 \end{pmatrix};
$$
  $x_5 = \begin{pmatrix} 6 & 5 \\ 4 & 5 \end{pmatrix} \begin{pmatrix} 5,556 \\ 4,444 \end{pmatrix} = \begin{pmatrix} 55,556 \\ 44,444 \end{pmatrix}.$ 

The dominant eigenvalue is approximated from

$$
\lambda_1 \approx \frac{x_4^{\mathrm{T}} x_5}{x_4^{\mathrm{T}} x_4} = \frac{506,178,271}{50,618,272} = 9.99991 \approx 10;
$$

and the corresponding eigenvector is approximated by  $x<sub>5</sub>$ . Dividing by the second component,

$$
v_1 = \begin{pmatrix} 55,556/44,444 \\ 1 \end{pmatrix} \approx \begin{pmatrix} 1.25002 \\ 1 \end{pmatrix} \approx \begin{pmatrix} 5/4 \\ 1 \end{pmatrix}.
$$

**1.** Use the power method (without normalizing the vectors after each iteration) to determine the dominant eigenvalue and corresponding eigenvector of the matrix

$$
A = \begin{pmatrix} -5 & 6 \\ 5 & -4 \end{pmatrix}.
$$

# **Lecture 32 | Matrix algebra in** MATLAB

#### [View this lecture on YouTube](https://youtu.be/QimxbtqF8es)

In MATLAB, the solution to the matrix equation  $Ax = b$  is found by typing  $x = A \backslash b$ . The backslash operator can solve *n* equations and *n* unknowns, or a least-squares problem that has more equations than unknowns. For example, to solve the following system of three linear equations and three unknowns:

$$
\begin{pmatrix} -3 & 2 & -1 \ 6 & -6 & 7 \ 3 & -4 & 4 \ \end{pmatrix} \begin{pmatrix} x_1 \ x_2 \ x_3 \end{pmatrix} = \begin{pmatrix} -1 \ -7 \ -6 \end{pmatrix},
$$

#### we type

```
\Rightarrow A=[-3 2 -1; 6 -6 7; 3 -4 4]; b=[-1;-7;-6];
>> x=A\b
x =2
      2
     -1
>>
```
To solve the toy least-squares problem of three equations and two unknowns:

$$
\begin{pmatrix} 1 & 1 \ 1 & 2 \ 1 & 3 \end{pmatrix} \begin{pmatrix} x_1 \ x_2 \end{pmatrix} = \begin{pmatrix} 1 \ 3 \ 2 \end{pmatrix},
$$

we type

```
>> A=[1 1; 1 2; 1 3]; b=[1; 3; 2];
>> x=A\b
x =1.0000
     0.5000
```

```
\gt
```
For a problem with one matrix and many different right-hand-sides, one first finds  $A =$ LU by a function call, and then solves for x using  $x=U\setminus (L\setminus b)$ . For example, to solve again the system of three linear equations and three unknowns, we could type

```
\Rightarrow A=[-3 2 -1; 6 -6 7; 3 -4 4]; b=[-1; -7; -6];
>> [L U] = lu(A);>> x = U\(L\)x =
```

```
2
 2
-1
```
>>

MATLAB can also solve the eigenvalue problem  $Ax = \lambda x$ . The function eig.m can be used to find all the eigenvalues and eigenvectors of A. For example, to find all the eigenvalues of a two-by-two matrix A, we type

```
>> A=[0 1; 1 0];
>> lambda = eig(A)
lambda =
    -1
      1
\rightarrowTo also find the eigenvectors, we type
>> A=[0 1; 1 0];
>> [V, D] = eig(A)V =-0.7071 0.7071
    0.7071 0.7071
D =-1 0
     0 1
>>
```
The latter function call finds two matrices V and D that satisfy the matrix equation  $AV = VD$ .

**1.** Let

$$
A = \begin{pmatrix} -3 & 2 & -1 \\ 6 & -6 & 7 \\ 3 & -4 & 4 \end{pmatrix}.
$$

Use MATLAB to find the LU decomposition of A, where U is an upper triangular matrix and L is a psychologically lower triangular matrix.

2. Use MATLAB to find the eigenvalues and eigenvectors of the matrix

$$
A = \begin{pmatrix} -5 & 6 \\ 5 & -4 \end{pmatrix}.
$$

Normalize the eigenvectors so that their second component is one.

# **Lecture 33 | Systems of nonlinear equations**

[View this lecture on YouTube](https://youtu.be/wrQA1_nei7w)

A system of nonlinear equations can be solved using Newton's method and matrix algebra. We derive the method for a system of two equations and two unknowns. Suppose that we want to solve

$$
f(x,y) = 0, \qquad g(x,y) = 0,
$$

for the unknowns *x* and *y*. To solve, we construct two parallel sequences  $x_0, x_1, x_2, \ldots$ and  $y_0, y_1, y_2, \ldots$  that converge to the root. To construct these sequences, we Taylor series expand  $f(x_{n+1}, y_{n+1})$  and  $g(x_{n+1}, y_{n+1})$  about the point  $(x_n, y_n)$ . Using the partial derivatives  $f_x = \frac{\partial f}{\partial x}$ ,  $f_y = \frac{\partial f}{\partial y}$  and those for *g*, the two-dimensional Taylor series expansions, displaying up to the linear terms, are given by

$$
f(x_{n+1}, y_{n+1}) = f(x_n, y_n) + (x_{n+1} - x_n) f_x(x_n, y_n) + (y_{n+1} - y_n) f_y(x_n, y_n) + \dots,
$$
  

$$
g(x_{n+1}, y_{n+1}) = g(x_n, y_n) + (x_{n+1} - x_n) g_x(x_n, y_n) + (y_{n+1} - y_n) g_y(x_n, y_n) + \dots
$$

To obtain Newton's method, we set  $f(x_{n+1}, y_{n+1}) = g(x_{n+1}, y_{n+1}) = 0$ , and drop higherorder terms above linear. It is best to define the unknowns as

$$
\Delta x_n = x_{n+1} - x_n, \qquad \Delta y_n = y_{n+1} - y_n,
$$

so that the iteration equations will be given by

$$
x_{n+1} = x_n + \Delta x_n, \qquad y_{n+1} = y_n + \Delta y_n.
$$

The linear equations to be solved for ∆*x<sup>n</sup>* and ∆*yn*, in matrix form, are

$$
\begin{pmatrix} f_x & f_y \\ g_x & g_y \end{pmatrix} \begin{pmatrix} \Delta x_n \\ \Delta y_n \end{pmatrix} = - \begin{pmatrix} f \\ g \end{pmatrix},
$$

where *f*, *g*, *f<sub>x</sub>*, *fy*, *g<sub>x</sub>* and *g<sub>y</sub>* are all evaluated at the point  $(x_n, y_n)$ . The two-equation case is easily generalized to *n* equations. The matrix of partial derivatives is called the Jacobian matrix.

Each step of the iteration requires solving the *n*-by-*n* system of linear equations using Gaussian elimination.

**1.** The algorithm for solving the system of two equations and two unknowns,

$$
f(x, y) = 0, \qquad g(x, y) = 0,
$$

is given by the following two-step process.

1. Solve the linear system for  $\Delta x_n$  and  $\Delta y_n$  given by

$$
\begin{pmatrix} f_x & f_y \ g_x & g_y \end{pmatrix} \begin{pmatrix} \Delta x_n \\ \Delta y_n \end{pmatrix} = \begin{pmatrix} -f \\ -g \end{pmatrix}.
$$

2. Advance the iterative solution, using

$$
x_{n+1} = x_n + \Delta x_n, \qquad y_{n+1} = y_n + \Delta y_n.
$$

Write down the corresponding algorithm for three equations and three unknowns.

# **Lecture 34 | Systems of nonlinear equations (example)**

[View this lecture on YouTube](https://youtu.be/IaZPy6noXc8)

The Lorenz equations — a system of nonlinear odes that pioneered the study of chaos — are given by

$$
\dot{x} = \sigma(y - x), \qquad \dot{y} = x(r - z) - y, \qquad \dot{z} = xy - \beta z,
$$

where  $\sigma$ ,  $\beta$  and  $r$  are constant parameters. The fixed-point solutions of these equations satisfy  $\dot{x} = \dot{y} = \dot{z} = 0$ . Although the fixed-point solutions can be found analytically, here we will illustrate Newton's method. We solve

$$
\sigma(y - x) = 0
$$
,  $rx - y - xz = 0$ ,  $xy - \beta z = 0$ .

This is a root-finding problem, where we are solving  $f(x, y, z) = g(x, y, z) = h(x, y, z) = 0$ , and

$$
f(x,y,z) = \sigma(y-x), \qquad g(x,y,z) = rx - y - xz, \qquad h(x,y,z) = xy - \beta z.
$$

The Jacobian matrix is found from the partial derivatives, computed to be

$f_x = -\sigma$ ,	$f_y = \sigma$ ,	$f_z = 0$ ,
$g_x = r - z$ ,	$g_y = -1$ ,	$g_z = -x$ ,
$h_x = y$ ,	$h_y = x$ ,	$h_z = -\beta$ .

Using the Jacobian matrix, the iteration equation for Newton's method is given by

$$
\begin{pmatrix}\n-\sigma & \sigma & 0 \\
r - z_n & -1 & -x_n \\
y_n & x_n & -\beta\n\end{pmatrix}\n\begin{pmatrix}\n\Delta x_n \\
\Delta y_n \\
\Delta z_n\n\end{pmatrix} = -\begin{pmatrix}\n\sigma(y_n - x_n) \\
rx_n - y_n - x_n z_n \\
x_n y_n - \beta z_n\n\end{pmatrix},
$$

with

$$
x_{n+1} = x_n + \Delta x_n
$$
,  $y_{n+1} = y_n + \Delta y_n$ ,  $z_{n+1} = z_n + \Delta z_n$ .

**1.** The fixed-point solutions of the Lorenz equations satisfy

$$
\sigma(y - x) = 0
$$
,  $rx - y - xz = 0$ ,  $xy - \beta z = 0$ .

Find analytically three fixed-point solutions for  $(x, y, z)$  as a function of the parameters *σ*, *β* and *r*. What are the numerical values for the fixed points when  $r = 28$ ,  $σ = 10$  and  $\beta = 8/3?$ 

**2.** Complete a MATLAB code that uses Newton's method to determine the fixed-point solutions of the Lorenz equations. Solve using the parameters  $r = 28$ ,  $\sigma = 10$  and  $\beta = 8/3$ . Use as your three initial guesses  $x = y = z = 1$ ,  $x = y = z = 10$  and  $x = y = z = -10$ .

# **Lecture 35 | Project III: Fractals from the Lorenz equations**

[View this lecture on YouTube](https://youtu.be/4fV2fm6whEQ)

The fixed points of the Lorenz equations satisfy

$$
\sigma(y - x) = 0
$$
,  $rx - y - xz = 0$ ,  $xy - \beta z = 0$ .

Analytical solution of the fixed-point equations yields the three roots,

$$
(x,y,z)
$$
 =  $(0,0,0)$ ,  $(\pm \sqrt{\beta(r-1)}, \pm \sqrt{\beta(r-1)}, r-1)$ .

Implementation of Newton's method to determine the solution of the fixed-point equations requires iteration of

$$
\begin{pmatrix}\n-\sigma & \sigma & 0 \\
r - z_n & -1 & -x_n \\
y_n & x_n & -\beta\n\end{pmatrix}\n\begin{pmatrix}\n\Delta x_n \\
\Delta y_n \\
\Delta z_n\n\end{pmatrix} = -\begin{pmatrix}\n\sigma(y_n - x_n) \\
rx_n - y_n - x_n z_n \\
x_n y_n - \beta z_n\n\end{pmatrix},
$$

with

$$
x_{n+1} = x_n + \Delta x_n
$$
,  $y_{n+1} = y_n + \Delta y_n$ ,  $z_{n+1} = z_n + \Delta z_n$ .

We may suppose that an initial value given by  $(x_0, y_0, z_0)$  will converge to one of the three roots of the fixed-point equations. By mapping a grid of initial values to the converged root, an image of a fractal may result. Your third project is to compute an instance of one of these fractals.

**1.** Determine a fractal that arises from using Newton's method to compute the fixed-point solutions of the Lorenz equations. Use the parameter values  $r = 28$ ,  $\sigma = 10$  and  $\beta = 8/3$ . Initial values  $(x_0, z_0)$  are taken on a grid in the *x*-*z* plane with always  $y_0 = 3\sqrt{2}$ . For assessment purposes, the computational grid and the graphics code will be given in the Learner Template. To pass the assessment, every pixel in your figure needs to be colored correctly.

(Hint: Some grid points may require as many as 33 Newton iterations to converge while others may require as few as three. Unfortunately, if you uniformly use 33 Newton iterations at every grid point, the Matlab Grader may time out. You can accelerate your code by using a while loop instead of a for loop.)

## **Week IV**

## **Quadrature and Interpolation**

*In this week's lectures, we learn about numerical integration—also called quadrature—and function interpolation. We begin by learning the basics of quadrature, which include the elementary formulas for the trapezoidal rule and Simpson's rule, and how these formulas can be used to develop composite integration rules. We then discuss the more sophisticated method of Gaussian quadrature, and learn how to construct an adaptive integration routine in which the software itself determines the appropriate integration step size. Finally, we learn how to use the* Matlab *function* integral.m*.*

*In the second-half of this week we learn about interpolation. Given known points in the plane arising from a smooth function, a good interpolation routine will be able to estimate the values between these points. Linear interpolation is widely used, particularly when plotting data consisting of many points. The more sophisticated method of cubic splines may be a better choice, particularly if the points are more sparse.*

*Your programming project will be to write a* Matlab *code to compute the zeros of the Bessel functions. This requires combining both quadrature and root-finding routines.*

## **Lecture 36 | Midpoint rule**

[View this lecture on YouTube](https://youtu.be/9RZZgQl5TpU)

Elementary integration formulas are derived by integrating  $f(x)$  from 0 to  $h$ , and these formulas serve as building blocks for integrating over any interval. Here, we derive the midpoint rule.

We define the definite integral *I<sup>h</sup>* to be

$$
I_h = \int_0^h f(x) dx.
$$

To perform this integral, we Taylor series expand  $f(x)$  about the value  $x = h/2$ :

$$
f(x) = f(h/2) + (x - h/2)f'(h/2) + \frac{(x - h/2)^2}{2}f''(h/2) + \frac{(x - h/2)^3}{6}f'''(h/2) + \frac{(x - h/2)^4}{24}f''''(h/2) + \dots
$$

Integrating both sides from zero to *h*, and using the definition of *I<sup>h</sup>* , we have

$$
I_h = \int_0^h \left( f(h/2) + (x - h/2)f'(h/2) + \frac{(x - h/2)^2}{2} f''(h/2) + \frac{(x - h/2)^3}{6} f'''(h/2) + \frac{(x - h/2)^4}{24} f''''(h/2) + \dots \right) dx.
$$

To perform the integration, we let  $y = x - h/2$  so that  $dy = dx$ , and simplify according to whether each term in the integrand is even or odd. We have

$$
I_h = \int_{-h/2}^{h/2} \left( f(h/2) + y f'(h/2) + \frac{y^2}{2} f''(h/2) + \frac{y^3}{6} f'''(h/2) + \frac{y^4}{24} f''''(h/2) + \dots \right) dy
$$
  
=  $2 \int_0^{h/2} \left( f(h/2) + \frac{y^2}{2} f''(h/2) + \frac{y^4}{24} f''''(h/2) + \dots \right) dy.$ 

Using

$$
\int_0^{h/2} y^2 dy = h^3 / 24, \qquad \int_0^{h/2} y^4 dy = h^5 / 160,
$$

we obtain the midpoint rule, given by

$$
I_h = hf(h/2) + \frac{h^3}{24}f''(h/2) + \frac{h^5}{1920}f''''(h/2) + \dots
$$
  
=  $hf(h/2) + O(h^3)$ .

**1.** The area of a rectangle with base *h* and height  $f(h/2)$  is given by  $hf(h/2)$ . Draw a graph illustrating the midpoint rule.

**2.** Let  $f(x) = a + bx + cx^2$ , where *a*, *b* and *c* are constants. Prove by explicit calculation that

$$
\int_0^h f(x) \, dx = h f(h/2) + \frac{h^3}{24} f''(h/2).
$$

## **Lecture 37 | Trapezoidal rule**

[View this lecture on YouTube](https://youtu.be/K4x3aBY62hc)

The trapezoidal rule is one of the most useful integration formulas. From the Taylor series expansion of  $f(x)$  about  $x = h/2$ , we have

$$
f(0) = f(h/2) - \frac{h}{2}f'(h/2) + \frac{h^2}{8}f''(h/2) - \frac{h^3}{48}f'''(h/2) + \frac{h^4}{384}f''''(h/2) + ...,
$$
  

$$
f(h) = f(h/2) + \frac{h}{2}f'(h/2) + \frac{h^2}{8}f''(h/2) + \frac{h^3}{48}f'''(h/2) + \frac{h^4}{384}f''''(h/2) + ....
$$

Adding and multiplying by *h*/2 we obtain

$$
\frac{h}{2}(f(0)+f(h))=hf(h/2)+\frac{h^3}{8}f''(h/2)+\frac{h^5}{384}f''''(h/2)+\ldots
$$

We now substitute for  $hf(h/2)$  using the midpoint rule formula:

$$
\frac{h}{2}(f(0)+f(h))=\left(I_h-\frac{h^3}{24}f''(h/2)-\frac{h^5}{1920}f''''(h/2)\right)+\frac{h^3}{8}f''(h/2)+\frac{h^5}{384}f'''(h/2)+\ldots,
$$

and solving for *I<sup>h</sup>* , we obtain the trapezoidal rule,

$$
I_h = \frac{h}{2} (f(0) + f(h)) - \frac{h^3}{12} f''(h/2) - \frac{h^5}{480} f'''(h/2) + \dots
$$
  
=  $\frac{h}{2} (f(0) + f(h)) + O(h^3).$ 

A graphical representation of the trapezoidal rule is shown on the right, where the integral is approximated by the area of the drawn trapezoid.



**1.** Derive the trapezoidal rule by approximating  $f(x)$  by the straight line connecting the points  $(0, f(0))$  and  $(h, f(h))$ :

$$
f(x) \approx f(0) + \frac{f(h) - f(0)}{h}x.
$$

## **Lecture 38 | Simpson's rule**

[View this lecture on YouTube](https://youtu.be/Z6-W6AXqFFs)

To obtain Simpson's rule, we combine the midpoint and trapezoidal rule to eliminate the error term proportional to *h* 3 . The midpoint and trapezoidal rules are given by

$$
I_h = hf(h/2) + \frac{h^3}{24}f''(h/2) + \frac{h^5}{1920}f'''(h/2) + \dots
$$
  
\n
$$
I_h = \frac{h}{2}(f(0) + f(h)) - \frac{h^3}{12}f''(h/2) - \frac{h^5}{480}f''''(h/2) + \dots
$$

Multiplying the midpoint rule by two and adding it to the trapezoidal rule, we obtain

$$
3I_h = h\left(2f(h/2) + \frac{1}{2}\left(f(0) + f(h)\right)\right) + h^5\left(\frac{2}{1920} - \frac{1}{480}\right)f''''(h/2) + \dots,
$$

or

$$
I_h = \frac{h}{6} (f(0) + 4f(h/2) + f(h)) - \frac{h^5}{2880} f''''(h/2) + \dots
$$
  
=  $\frac{h}{6} (f(0) + 4f(h/2) + f(h)) + O(h^5).$ 

Usually, Simpson's rule is written by considering the three consecutive points 0, *h* and 2*h*. Transforming  $h \to 2h$ , we obtain the more standard result

$$
I_{2h} = \frac{h}{3}(f(0) + 4f(h) + f(2h)) - \frac{h^5}{90}f'''(h) + \dots
$$
  
=  $\frac{h}{3}(f(0) + 4f(h) + f(2h)) + O(h^5).$ 

**1.** Derive Simpson's rule by approximating  $f(x)$  by a quadratic polynomial connecting the points  $(0, f(0))$ ,  $(h, f(h))$  and  $(2h, f(2h))$ .

- *a*) Let  $g(x) = a + bx + cx^2$ . Determine the values of *a*, *b* and *c* such that  $g(x)$  passes through the points  $(0, f(0))$ ,  $(h, f(h))$  and  $(2h, f(2h))$ .
- *b*) Use  $f(x) \approx g(x)$  to derive Simpson's rule.

# **Lecture 39 | Composite quadrature rules**

#### [View this lecture on YouTube](https://youtu.be/6x28YeHIkc4)

We apply our elementary integration formulas to compute

$$
I = \int_{a}^{b} f(x) \, dx.
$$

Suppose that the function  $f(x)$  is known at the points  $a = x_0, x_1, \ldots, x_n = b$ . Define

$$
f_i = f(x_i), \quad h_i = x_{i+1} - x_i.
$$

Then the integral *I* may be decomposed as

$$
\int_a^b f(x)dx = \sum_{i=0}^{n-1} \int_{x_i}^{x_{i+1}} f(x)dx = \sum_{i=0}^{n-1} \int_0^{h_i} f(x_i+s)ds,
$$

where the last equality arises from the change-of-variables  $s = x - x_i$ . Applying the trapezoidal rule to the integral, we have

$$
\int_a^b f(x)dx = \frac{1}{2}\sum_{i=0}^{n-1} h_i (f_i + f_{i+1}),
$$

which is often useful for integrating experimental data.

If the points are evenly spaced, we have  $h_i = h = (b - a)/n$ , and the composite trapezoidal rule becomes

$$
\int_a^b f(x)dx = \frac{h}{2} (f_0 + 2f_1 + \cdots + 2f_{n-1} + f_n).
$$

The first and last terms have a multiple of one, all other terms have a multiple of two, and the entire sum is multiplied by *h*/2.

Similarly, the composite Simpson's rule for evenly space points is found to be

$$
\int_{a}^{b} f(x)dx = \frac{h}{3} (f_0 + 4f_1 + f_2) + \frac{h}{3} (f_2 + 4f_3 + f_4) + \dots + \frac{h}{3} (f_{n-2} + 4f_{n-1} + f_n)
$$
  
=  $\frac{h}{3} (f_0 + 4f_1 + 2f_2 + 4f_3 + 2f_4 + \dots + 4f_{n-1} + f_n).$ 

Note that *n* must be even for this scheme to work. The first and last terms have a multiple of one, the odd indexed terms have a multiple of 4, the even indexed terms have a multiple of 2, and the entire sum is multiplied by *h*/3.

For composite rules, the error for each elementary interval sums so that the global error increases by a factor of 1/*h*. The trapezoidal rule and Simpson's rule have global error  $h^2$  and  $h^4$ , respectively.

**1.** Simpson's 3/8 rule has elementary formula given by

$$
\int_0^{3h} f(x) dx = \frac{3h}{8} (f(0) + 3f(h) + 3f(2h) + f(3h)).
$$

Suppose that *f*(*x*) is known at the equally spaced points  $a = x_0, x_1, \ldots, x_n = b$ , and *n* is a multiple of three. Let  $f_i = f(x_i)$  and  $h = x_{i+1} - x_i$ . Find the formula for the composite Simpson's 3/8 rule.

## **Lecture 40 | Gaussian quadrature**

[View this lecture on YouTube](https://youtu.be/w2xjlPwYock)

A general Gaussian quadrature rule is of the form

$$
\int_a^b W(x)f(x)\,dx \approx \sum_{i=1}^n w_i f(x_i),
$$

where  $W(x)$  is called the weight function,  $w_i$  are called the weights, and  $x_i$  are called the nodes. The weights and nodes are chosen so that the integral is exact for polynomial functions  $f(x)$  of degree less than or equal to  $2n - 1$ . Variations include Chebyshev-Gauss, Laguerre-Gauss, and Hermite-Gauss quadrature. The definite integrals associated with these Gaussian quadratures are given, respectively, by

$$
\int_{-1}^{1} \frac{1}{\sqrt{1-x^2}} f(x) \, dx, \qquad \int_{0}^{\infty} e^{-x} f(x) \, dx, \qquad \int_{-\infty}^{\infty} e^{-x^2} f(x) \, dx.
$$

The most straightforward example is the Legendre-Gauss quadrature rule, with the twopoint formula given by

$$
\int_{-1}^{1} f(x) dx \approx w_1 f(x_1) + w_2 f(x_2).
$$

To determine the weights and the nodes, we require the rule to be exact for the polynomial basis functions  $f(x) = 1$ , x,  $x^2$  and  $x^3$ . Substituting in these functions yields the four equations with four unknowns given by

$$
2 = w_1 + w_2, \quad 0 = w_1 x_1 + w_2 x_2, \quad 2/3 = w_1 x_1^2 + w_2 x_2^2, \quad 0 = w_1 x_1^3 + w_2 x_2^3.
$$

If we argue from symmetry that  $x_1 = -x_2$  and  $w_1 = w_2$ , then a solution is readily found to be  $w_1 = w_2 = 1$ , and  $x_1 = -1/\sqrt{3}$  and  $x_2 = 1/\sqrt{3}$ .

**1.** Determine the weights and nodes of the three-point Legendre-Gauss quadrature rule. You may assume from symmetry that  $x_1 = -x_3$ ,  $w_1 = w_3$ , and  $x_2 = 0$ .

## **Lecture 41 | Adaptive quadrature**

[View this lecture on YouTube](https://youtu.be/U4NUXAwwR8E)

Adaptive quadrature lets the algorithm itself determine the integration step size required to meet a user-specified precision. Even more importantly, the step size need not be constant over the entire region of integration. Using the trapezoidal rule, we sketch how adaptive quadrature is implemented.

We begin integration of  $I = \int_a^b f(x) dx$  at what is called Level 1. Let  $h = b - a$ . Then integration from *a* to *b* using the trapezoidal rule is given by

$$
I = \frac{h}{2} \left( f(a) + f(b) \right) - \frac{h^3}{12} f''(\xi),
$$

where the Taylor remainder theorem states that *ξ* is a number between *a* and *b*. We now add a point *c* midway between *a* and *b* and apply the composite trapezoidal rule to obtain

$$
I = \frac{h}{4} \left( f(a) + 2f(c) + f(b) \right) - \frac{(h/2)^3}{12} f''(\xi_l) - \frac{(h/2)^3}{12} f''(\xi_r),
$$

where *ξ<sup>l</sup>* is a number between *a* and *c* and *ξ<sup>r</sup>* is a number between *c* and *b*.

We can define our two approximations to the integral by

$$
S_1 = \frac{h}{2}(f(a) + f(b)),
$$
  $S_2 = \frac{h}{4}(f(a) + 2f(c) + f(b)),$ 

and our two associated errors by

$$
E_1 = -\frac{h^3}{12}f''(\xi), \qquad E_2 = -\frac{h^3}{2^3 \cdot 12} \big(f''(\xi_l) + f''(\xi_r)\big).
$$

If we make the approximation that  $f''(\xi) \approx f''(\xi_l) \approx f''(\xi_r)$ , then  $E_1 \approx 4E_2$ . Then since  $S_1 + E_1 = S_2 + E_2$ , we can eliminate  $E_1$  to find an estimate for  $E_2$ :

$$
|E_2| \approx |S_2 - S_1|/3.
$$

If  $|E_2|$  is less than some specified tolerance, then we accept  $S_2$  as *I*. Otherwise, we proceed to Level 2.

The computation at Level 2 starts with the two integration intervals *a* to *c* and *c* to *b*, and repeats the above procedure independently on both intervals. Integration can be stopped on either interval provided the estimated error on that interval is less than onehalf the specified tolerance (since the sum of the errors on both intervals must be less than the tolerance). Otherwise, either interval can proceed to Level 3, and so on, until the global error estimate is less than the tolerance.

**1.** Consider  $I = \int_0^h f(x) dx$  with  $f(x) = x^3$ . Using the trapezoidal rule, compute  $S_1$ ,  $S_2$ ,  $E_1$  and  $E_2$  and show that  $E_1 = 4E_2$ .

# **Lecture 42 | Quadrature in** MATLAB

[View this lecture on YouTube](https://youtu.be/wbfzKRM_6lI)

The MATLAB function integral.m performs adaptive quadrature by combining a sevenpoint Legendre-Gauss quadrature rule with a fifteen-point Kronrod rule. We will skip the details here.

If *S* is the computed value of the integral and *I* is the (unknown) exact value, then the absolute error is defined as  $|S - I|$  and the relative error is defined as  $|(S - I)/I|$ . The user can specify both an absolute error tolerance and a relative error tolerance, with default values given by  $10^{-10}$  and  $10^{-6}$ , respectively. MATLAB will refine the integration intervals until either the absolute or relative error tolerance is lower than the default (or specified) values.

In practice, one begins a calculation with the default error tolerances and if the accuracy of the solution is unsatisfactory, then the error tolerances can be lowered. Consult the Matlab help page to learn how to do this. A good rule of thumb is to lower each error tolerance by a factor of ten until you are satisfied. The absolute error tolerance usually comes into play when the integral is close to zero.

The default call is  $I=$ integral(f,xmin,xmax), where f is the integrand and xmin and  $x$ max are the limits of integration. The passing of  $f$  to integral.m is similar to what we have done before.

*Compute the following integral:*

$$
I = \int_0^\infty e^{-ax^2} (\ln x)^2 dx, \qquad a = \pi.
$$

The following MATLAB commands compute the integral:

```
>> f = \theta(x, a) \exp(-a*x.^2) . *log(x) . ^2;\Rightarrow a = pi; I = integral(\theta(x) f(x,a), 0, Inf)
I =1.8245
>>
```
**1.** Consider the Fresnel integrals, defined by

$$
C(t) = \int_0^t \cos(\frac{1}{2}\pi x^2) dx, \qquad S(t) = \int_0^t \sin(\frac{1}{2}\pi x^2) dx.
$$

Write a script using integral.m to plot a Cornu spiral, which is a smooth curve of *C*(*t*) versus *S*(*t*). Plot your solution over the range  $-8 \le t \le 8$ .

## **Lecture 43 | Interpolation**

[View this lecture on YouTube](https://youtu.be/RpxoN9-i7Jc)

Given the *y*-values of a function sampled at an ordered sequence of *x*-values, i.e., (*x*0, *y*0),  $(x_1, y_1), \ldots, (x_n, y_n)$ , we want to estimate the *y*-value for some other value of *x*. When  $x_0 \leq x \leq x_n$ , the problem is called interpolation; when  $x < x_0$  or  $x > x_n$ , the problem is called extrapolation and is much more perilous. The problem of interpolation is basically one of drawing a smooth curve that passes through the given  $n + 1$  points.

It is possible to interpolate  $n + 1$  known points by a unique polynomial of degree *n*. With only two points, the polynomial is a linear function; with only three points, the polynomial is a quadratic function, and so on. Although low-order polynomials are sometimes helpful when interpolating only a few points, high-order polynomials tend to over-oscillate and not be useful.

In this course, we will learn about the more widely-used piecewise polynomial interpolation. The two most popular are piecewise linear interpolation and cubic spline interpolation. The first makes use of linear polynomials, and the second cubic polynomials. Piecewise linear interpolation is the default interpolation typically used when plotting data, and we start here.

Suppose the interpolating function is  $y = g(x)$ , and there are  $n + 1$  points to interpolate. We construct the function  $g(x)$  out of *n* local linear polynomials. We write

$$
g(x) = g_i(x), \quad \text{for } x_i \leq x \leq x_{i+1},
$$

where

$$
g_i(x) = a_i(x - x_i) + b_i
$$
, for  $i = 0, 1, ..., n - 1$ .

We now require  $y = g_i(x)$  to pass through the endpoints  $(x_i, y_i)$  and  $(x_{i+1}, y_{i+1})$ . We have

$$
y_i = b_i
$$
,  $y_{i+1} = a_i(x_{i+1} - x_i) + b_i$ .

The solution for the coefficients of  $g_i(x)$  is easily determined to be

$$
a_i = \frac{y_{i+1} - y_i}{x_{i+1} - x_i}, \qquad b_i = y_i.
$$

Although piecewise linear interpolation is widely used, particularly in plotting routines, it suffers from discontinuities in the derivative. This results in a function which doesn't look smooth if the points are too widely spaced. One may increase the number of function evaluations, but sometimes the computational cost to do so is too high. In the next two lectures we will learn about cubic spline interpolation that uses piecewise cubic polynomials.

- **1.** Consider the points  $(0,0)$ ,  $(1,1)$  and  $(2,1)$ .
	- *a*) Find the quadratic polynomial that interpolates these points. What are the interpolated *y*-values at  $x = 1/2$  and  $x = 3/2$ ?
	- *b*) Find the two piecewise linear polynomials that interpolate these points. What are the interpolated *y*-values at  $x = 1/2$  and  $x = 3/2$ ?
	- *c*) Use MATLAB to plot the three points and the two interpolating functions.

# **Lecture 44 | Cubic spline interpolation (Part A)**

[View this lecture on YouTube](https://youtu.be/LaolbjAzZvg)

For cubic spline interpolation, we define piecewise cubic polynomials by

$$
g_i(x) = a_i(x - x_i)^3 + b_i(x - x_i)^2 + c_i(x - x_i) + d_i
$$
, for  $i = 0$  to  $n - 1$  and  $x_i \le x \le x_{i+1}$ .

Requiring the global interpolating function  $y = g(x)$  to go through all  $n + 1$  points and be continuous results in the two constraints

$$
g_i(x_i) = y_i
$$
,  $g_i(x_{i+1}) = y_{i+1}$ , for  $i = 0$  to  $n - 1$ .

To achieve a smooth interpolation, we further require the first and second derivatives of  $g(x)$  to be continuous. This results in the additional constraints

$$
g'_{i}(x_{i+1}) = g'_{i+1}(x_{i+1}),
$$
  $g''_{i}(x_{i+1}) = g''_{i+1}(x_{i+1}),$  for  $i = 0$  to  $n - 2$ .

There are *n* piecewise cubic polynomials  $g_i(x)$  and each cubic polynomial has four free coefficients, for a total of  $4n$  coefficients. The number of constraining equations is  $2n +$  $2(n-1) = 4n - 2$  so that we are missing two equations to obtain a unique solution. These equations usually derive from some extra conditions imposed on  $g_0(x)$  and  $g_{n-1}(x)$ . We will discuss this in the next lecture.

We define now for convenience

$$
h_i = x_{i+1} - x_i
$$
,  $\eta_i = y_{i+1} - y_i$ , for  $i = 0$  to  $n - 1$ .

Calculating the derivatives and substituting in the end points, our constraints lead directly to the equations

$$
d_i = y_i, \t a_i h_i^3 + b_i h_i^2 + c_i h_i = \eta_i, \t for i = 0 to n - 1;
$$
  
\n
$$
3a_i h_i^2 + 2b_i h_i + c_i = c_{i+1}, \t 6a_i h_i + 2b_i = 2b_{i+1}, \t for i = 0 to n - 2.
$$

It will be useful later for us to include a definition of the coefficient *bn*, which is now missing. We simply extend the fourth of the above equations to  $i = n - 1$  (and thus add one more equation and one more coefficient). How to solve these equations will be the focus of our next lecture.

**1.** Let  $y = f(x)$  have known values  $(x_0, y_0)$ ,  $(x_1, y_1)$ , ...,  $(x_n, y_n)$ , and define piecewise cubic polynomials by

 $g_i(x) = a_i(x - x_i)^3 + b_i(x - x_i)^2 + c_i(x - x_i) + d_i$ , for  $i = 0$  to  $n - 1$  and  $x_i \le x \le x_{i+1}$ .

Suppose that the endpoint slopes  $f'(x_0) = y'_0$  and  $f'(x_n) = y'_n$  are known. From these two extra conditions, determine two extra constraints on the *a*, *b* and *c* coefficients.

# **Lecture 45 | Cubic spline interpolation (Part B)**

[View this lecture on YouTube](https://youtu.be/4VpE9Tbie14)

We now solve for the *a*, *b*, *c*, and *d* coefficients. The equations are

$$
d_i = y_i, \t a_i h_i^3 + b_i h_i^2 + c_i h_i = \eta_i, \t 3a_i h_i + b_i = b_{i+1}, \t for i = 0 to n - 1;
$$
  

$$
3a_i h_i^2 + 2b_i h_i + c_i = c_{i+1}, \t for i = 0 to n - 2.
$$

We can use the third equation to eliminate *a<sup>i</sup>* :

$$
a_i = \frac{1}{3h_i} (b_{i+1} - b_i), \text{ for } i = 0 \text{ to } n - 1.
$$

Then the second equation becomes

$$
c_i = \frac{1}{h_i} \left( \eta_i - a_i h_i^3 - b_i h_i^2 \right) = \frac{\eta_i}{h_i} - \frac{1}{3} h_i \left( b_{i+1} + 2 b_i \right), \quad \text{for } i = 0 \text{ to } n-1.
$$

Finally, eliminating the  $a$  and  $c$  coefficients from the last equation results, after some algebra, in

$$
\frac{1}{3}h_i b_i + \frac{2}{3}(h_i + h_{i+1})b_{i+1} + \frac{1}{3}h_{i+1} b_{i+2} = \frac{\eta_{i+1}}{h_{i+1}} - \frac{\eta_i}{h_i}, \quad \text{for } i = 0 \text{ to } n-2,
$$

which is a system of  $n - 1$  equations for the  $n + 1$  unknowns  $b_0, b_1, \ldots, b_n$ . We can represent these equations in matrix form as

$$
\begin{pmatrix}\n... & ... & ... & ... & \text{missing} & ... & ... \\
\frac{1}{3}h_0 & \frac{2}{3}(h_0 + h_1) & \frac{1}{3}h_1 & ... & 0 & 0 & 0 \\
... & ... & ... & ... & \frac{1}{3}h_{n-2} & \frac{2}{3}(h_{n-2} + h_{n-1}) & \frac{1}{3}h_{n-1} \\
... & ... & ... & ... & \text{missing} & ... & ...\n\end{pmatrix}\n\begin{pmatrix}\nb_0 \\
b_1 \\
\vdots \\
b_{n-1} \\
b_n\n\end{pmatrix}
$$
\n
$$
= \n\begin{pmatrix}\n\frac{\eta_1}{h_1} - \frac{\eta_0}{h_0} \\
\vdots \\
\frac{\eta_{n-1}}{h_{n-1}} - \frac{\eta_{n-2}}{h_{n-2}} \\
\vdots \\
\frac{\eta_{n-1}}{h_{n-1}} - \frac{\eta_{n-2}}{h_{n-2}}\n\end{pmatrix}
$$

where the first and last equations are still to be determined. Once we solve for the *b*coefficients, all the other coefficients can be found and the interpolating polynomial is determined.

The default way to specify the first and last equations is called the not-a-knot condition, which assumes that

$$
g_0(x) = g_1(x), \qquad g_{n-2}(x) = g_{n-1}(x).
$$

Consider the first constraint. Now, two cubic polynomials are equal if at some fixed *x*, the polynomials and their first three derivatives are equal. The cubic spline continuity conditions already ensure that  $g_0(x_1) = g_1(x_1)$ ,  $g'_0(x_1) = g'_1(x_1)$  and  $g''_0(x_1) = g''_1(x_1)$ . Therefore, the two polynomials are equal provided  $g_0'''(x_1) = g_1'''(x_1)$ , or  $a_0 = a_1$ . In terms of the *b* coefficients, this condition is given by

$$
h_1b_0 - (h_0 + h_1)b_1 + h_0b_2 = 0,
$$

which we can use as our missing first equation. A similar argument made for the second constraint yields for the missing last equation

$$
h_{n-1}b_{n-2} - (h_{n-2} + h_{n-1})b_{n-1} + h_{n-2}b_n = 0.
$$

**1.** Consider the points  $(0,0)$ ,  $(1,1)$ ,  $(2,1)$  and  $(3,2)$ . Using the not-a-knot condition, determine the four-by-four matrix equation for the *b* coefficients. Solve for the *b*'s as well as the *a*'s, *c*'s and *d*'s, and thus find the cubic spline interpolant. Plot your result. You may use MATLAB to assist in your algebra.
# **Lecture 46 | Interpolation in** MATLAB

[View this lecture on YouTube](https://youtu.be/0JrYTW-N2BA)

The MATLAB function  $interp1.m$  performs interpolation of a tabulated function. It is mainly called using  $yy=interpl(x, y, xx, method)$ ; where x and y are vectors containing the tabulated *x*-*y* function values and xx is a vector on whose values you want the function to be interpolated, with yy the returned interpolated values.

Several interpolation methods are possible, but the two we have discussed in this week assign method the values 'linear' or 'spline', where cubic spline interpolation uses the not-a-knot condition. There is also another MATLAB function spline.m that can use other conditions.

As an example, we show how to use  $\text{interpl.m}$  to interpolate the points on a sine curve using either piece-wise linear or cubic spline interpolation.

```
x=0:pi/4:2*pi; y=sin(x);plot(x,y,'ok'); hold on;
xx =linspace(0,2*pi,1000);
plot(xx,sin(xx),'k');
yy=interp1(x,y,xx,'linear');
plot(xx,yy,'b');
yy=interp1(x,y,xx,'spline');
plot(xx,yy, 'r');
```


1. Load the following MATLAB MAT-files:

data1.mat data2.mat

The file data1.mat contains the variables  $x1$  and  $y1$ . The file data2.mat contains the variables  $x2$  and  $y2$ . Use cubic spines to interpolate both  $y1$  and  $y2$  to a uniform x-grid. Plot  $y=y1+y2$  versus *x*. Your result should look like  $y = \sin x$ .

# **Lecture 47 | Project IV: Bessel functions and their**

### **zeros**

[View this lecture on YouTube](https://youtu.be/NoblEFn49sw)

The Bessel function of order *n* can be defined by the definite integral

$$
J_n(x) = \frac{1}{\pi} \int_0^{\pi} \cos (x \sin \theta - n\theta) d\theta, \quad \text{for } n = 0, 1, 2, \dots
$$

The zeros of the Bessel function play an important role in the solution of partial differential equations in cylindrical coordinates. The goal of this project is to compute the first five positive roots,  $j_{n,k}$  ( $k = 1, 2, ..., 5$ ) of the first six Bessel functions  $J_n(x)$  ( $n = 0, 1, ..., 5$ ) using quadrature and root-finding.



Before you start computing the roots, I would recommend plotting the Bessel functions  $J_n(x)$  versus *x* as I have done above. From the plot, you can easily visualize the behavior of the roots. To plot the Bessel functions, you can make use of the Matlab function integrate.m.

To find the roots of the Bessel functions, you can also make use of the MATLAB function fzero.m. Initial guesses for the roots can be obtained from your Bessel function plots, or from [Wolfram MathWorld](https://mathworld.wolfram.com/BesselFunctionZeros.html) (input by hand initial guesses accurate to no more than one decimal place). Generate your roots using two for loops, and output the roots in tabular form similar to that of [Wolfram MathWorld.](https://mathworld.wolfram.com/BesselFunctionZeros.html)

**1.** The Bessel function of order *n*, for  $n = 0, 1, 2, \ldots$ , can be defined by the definite integral

$$
J_n(x) = \frac{1}{\pi} \int_0^{\pi} \cos(x \sin \theta - n\theta) d\theta.
$$

Compute the first five positive roots  $j_{n,k}$ , ( $k = 1, 2, ..., 5$ ), of the first six Bessel functions  $J_n(x)$ ,  $(n = 0, 1, \ldots, 5)$ .

### **Week V**

### **Ordinary Differential Equations**

*In this week's lectures, we learn about the numerical integration of odes. The most basic method is called the Euler method, and it is a single-step, first-order method. The Runge-Kutta methods extend the Euler method to multiple steps and higher order, with the advantage that larger timesteps can be made. We show how to construct a family of second-order Runge-Kutta methods, and introduce you to the widely-used fourth-order Runge-Kutta method. These methods are easily adopted for solving systems of odes. We will show you how to use the* Matlab *function* ode45*, and how to solve a two-point boundary value ode using the shooting method.*

*Your programming project will be to write a* Matlab *code to compute the numerical simulation of the gravitational two-body problem.*

### **Lecture 48 | Euler method**

[View this lecture on YouTube](https://youtu.be/WMQ2Cac4sqw)



The Euler method is the most straightforward method to integrate a differential equation. The first-order differential equation  $\dot{x} = f(t, x)$ , with initial condition  $x(t_0) = x_0$ , provides the slope  $f(t_0, x_0)$  of the tangent line to the solution curve  $x = x(t)$  at the point  $(t_0, x_0)$ . With a small step size  $\Delta t = t_1 - t_0$ , the initial condition  $(t_0, x_0)$  can be marched forward to ( $t_1, x_1$ ) along the tangent line (see figure) using  $x_1 = x_0 + \Delta t f(t_0, x_0)$ . The point ( $t_1, x_1$ ) then becomes the new initial condition and is marched forward to  $(t_2, x_2)$  along a newly determined tangent line with slope given by  $f(t_1, x_1)$ . This iterative method is usually written as

$$
x_{n+1} = x_n + \Delta t f(t_n, x_n).
$$

For small enough ∆*t*, the numerical solution should converge to the exact solution of the ode, when such a solution exists.

The Euler Method has a local error, that is, the error incurred over a single time step, of  $O(\Delta t^2)$ . The global error, however, comes from integrating out to a time T. If this integration takes *N* time steps, then the global error is the sum of *N* local errors. Since  $N = T/\Delta t$ , the global error is given by  $O(\Delta t)$ , and it is customary to call the Euler Method a first-order method. We next consider some common second-order methods.

**1.** Let  $\dot{x} = b$ , with initial condition  $x(0) = x_0$  and *b* a constant. With  $t = n\Delta t$ , show that the Euler method results in the exact solution

$$
x(t) = x_0 + bt.
$$

## **Lecture 49 | Modified Euler method**

#### [View this lecture on YouTube](https://youtu.be/d3sZa3T2TJw)

The modified Euler method, also called Heun's method or the predictor-corrector method, is a second-order method. The idea is to average the value of *x*˙ at the beginning and end of each time step. Hopefully, we could write

$$
x_{n+1} = x_n + \frac{1}{2} \Delta t \big( f(t_n, x_n) + f(t_n + \Delta t, x_{n+1}) \big).
$$

The obvious problem with this formula is that the unknown value  $x_{n+1}$  appears on the right-hand-side. We can, however, estimate this value, in what is called the predictor step. For the predictor step, we use the Euler method to find

$$
x_{n+1}^p = x_n + \Delta t f(t_n, x_n).
$$

Then the corrector step becomes

$$
x_{n+1} = x_n + \frac{1}{2}\Delta t \big(f(t_n, x_n) + f(t_n + \Delta t, x_{n+1}^p)\big).
$$

The Modified Euler method is usually coded as

$$
k_1 = \Delta t f(t_n, x_n), \qquad k_2 = \Delta t f(t_n + \Delta t, x_n + k_1),
$$
  

$$
x_{n+1} = x_n + \frac{1}{2}(k_1 + k_2).
$$

The modified Euler method is one of a family of methods called second-order Runge-Kutta methods, which we derive in the next lecture.

**1.** Let  $\dot{x} = bt$ , with initial condition  $x(0) = x_0$  and *b* a constant. With  $t = n\Delta t$ , show that the Modified Euler method results in the exact solution

$$
x(t) = x_0 + \frac{1}{2}bt^2.
$$

### **Lecture 50 | Runge-Kutta methods**

[View this lecture on YouTube](https://youtu.be/C4UymmEi3Kw)

To illustrate the derivation of Runge-Kutta methods, we derive here the complete family of second-order methods. We march the solution of  $\dot{x} = f(t, x)$  forward by writing

$$
k_1 = \Delta t f(t_n, x_n), \qquad k_2 = \Delta t f(t_n + \alpha \Delta t, x_n + \beta k_1),
$$
  

$$
x_{n+1} = x_n + ak_1 + bk_2,
$$

where  $\alpha$ ,  $\beta$ ,  $\alpha$  and  $\beta$  are constants that define particular second-order methods. We will constrain the values of these constants by computing the Taylor series of  $x_{n+1}$  in two ways.

First, we compute the Taylor series for  $x_{n+1}$  directly:

$$
x_{n+1} = x(t_n + \Delta t) = x(t_n) + \Delta t \dot{x}(t_n) + \frac{1}{2} (\Delta t)^2 \ddot{x}(t_n) + O(\Delta t^3).
$$

Now,  $\dot{x}(t_n) = f(t_n, x_n)$ . The second derivative is more tricky and requires partial derivatives. We have

$$
\ddot{x}(t_n) = \left. \frac{d}{dt} f(t, x(t)) \right|_{t=t_n} = f_t(t_n, x_n) + \dot{x}(t_n) f_x(t_n, x_n) = f_t(t_n, x_n) + f(t_n, x_n) f_x(t_n, x_n).
$$

Putting all the terms together, we obtain

<span id="page-117-0"></span>
$$
x_{n+1} = x_n + \Delta t f(t_n, x_n) + \frac{1}{2} (\Delta t)^2 \left( f_t(t_n, x_n) + f(t_n, x_n) f_x(t_n, x_n) \right) + O(\Delta t^3).
$$
 (50.1)

Second, we compute the Taylor series for  $x_{n+1}$  from the Runge-Kutta formula. We start with

$$
x_{n+1} = x_n + a\Delta t f(t_n, x_n) + b\Delta t f(t_n + \alpha \Delta t, x_n + \beta \Delta t f(t_n, x_n)) + O(\Delta t^3);
$$

and the Taylor series that we need is

$$
f(t_n+\alpha\Delta t,x_n+\beta\Delta t f(t_n,x_n))=f(t_n,x_n)+\alpha\Delta t f_t(t_n,x_n)+\beta\Delta t f(t_n,x_n)f_x(t_n,x_n)+O(\Delta t^2).
$$

The Taylor-series for  $x_{n+1}$  from the Runge-Kutta method is therefore given by

<span id="page-117-1"></span>
$$
x_{n+1} = x_n + (a+b)\Delta t f(t_n, x_n) + (\Delta t)^2 (\alpha b f_t(t_n, x_n) + \beta b f(t_n, x_n) f_x(t_n, x_n)) + O(\Delta t^3).
$$
\n(50.2)

Comparing [\(50.1\)](#page-117-0) and [\(50.2\)](#page-117-1), we find three constraints for the four constants:

$$
a + b = 1
$$
,  $ab = 1/2$ ,  $bb = 1/2$ .

# **Lecture 51 | Second-order Runge-Kutta methods**

[View this lecture on YouTube](https://youtu.be/wIiCaSzMJFw)

The family of second-order Runge-Kutta methods that solve  $\dot{x} = f(t, x)$  is given by

$$
k_1 = \Delta t f(t_n, x_n), \qquad k_2 = \Delta t f(t_n + \alpha \Delta t, x_n + \beta k_1),
$$

$$
x_{n+1} = x_n + ak_1 + bk_2,
$$

where we have derived three constraints for the four constants *α*, *β*, *a* and *b*:

$$
a + b = 1
$$
,  $ab = 1/2$ ,  $\beta b = 1/2$ .

The modified Euler method corresponds to  $\alpha = \beta = 1$  and  $a = b = 1/2$ . The function *f*(*t*, *x*) is evaluated at the times *t* = *t*<sub>*n*</sub> and *t* = *t*<sub>*n*</sub> +  $\Delta$ *t*, and we have

$$
k_1 = \Delta t f(t_n, x_n), \qquad k_2 = \Delta t f(t_n + \Delta t, x_n + k_1),
$$
  

$$
x_{n+1} = x_n + \frac{1}{2}(k_1 + k_2).
$$

The midpoint method corresponds to  $\alpha = \beta = 1/2$ ,  $a = 0$  and  $b = 1$ . In this method, the function *f*(*t*, *x*) is evaluated at the times *t* = *t*<sub>*n*</sub> and *t* = *t*<sub>*n*</sub> +  $\Delta t/2$  and we have

$$
k_1 = \Delta t f(t_n, x_n), \qquad k_2 = \Delta t f\left(t_n + \frac{1}{2}\Delta t, x_n + \frac{1}{2}k_1\right),
$$

$$
x_{n+1} = x_n + k_2.
$$

**1.** Construct Ralston's method, which is a second-order Runge-Kutta method corresponding to  $\alpha = \beta = 3/4$ ,  $a = 1/3$  and  $b = 2/3$ .

**2.** Consider the ode given by

$$
\frac{dy}{dx} = f(x),
$$

with *y*(0) as the initial value. Use the second-order Runge-Kutta methods given by the midpoint rule and the modified Euler method to derive two elementary quadrature formulas.

# **Lecture 52 | Higher-order Runge-Kutta methods**

[View this lecture on YouTube](https://youtu.be/40NvJNrmHhw)

Higher-order Runge-Kutta methods can also be derived, but require substantially more algebra. For example, the general form of the third-order method is given by

$$
k_1 = \Delta t f(t_n, x_n), \qquad k_2 = \Delta t f(t_n + \alpha \Delta t, x_n + \beta k_1), \qquad k_3 = \Delta t f(t_n + \gamma \Delta t, x_n + \delta k_1 + \epsilon k_2),
$$

$$
x_{n+1} = x_n + ak_1 + bk_2 + ck_3,
$$

with constraints on the constants *α*, *β*, *γ*, *δ*, *e*, *a*, *b* and *c*. The fourth-order method has stages  $k_1$ ,  $k_2$ ,  $k_3$  and  $k_4$ . The fifth-order method requires at least six stages. The table below gives the order of the method and the minimum number of stages required.



Because the fifth-order method requires two more stages than the fourth-order method, the fourth-order method has found some popularity. The general fourth-order method with four stages has 13 constants and 11 constraints. A particularly simple fourth-order method that has been widely used in the past by physicists is given by

$$
k_1 = \Delta t f(t_n, x_n), \qquad k_2 = \Delta t f\left(t_n + \frac{1}{2}\Delta t, x_n + \frac{1}{2}k_1\right),
$$
  

$$
k_3 = \Delta t f\left(t_n + \frac{1}{2}\Delta t, x_n + \frac{1}{2}k_2\right), \qquad k_4 = \Delta t f\left(t_n + \Delta t, x_n + k_3\right);
$$
  

$$
x_{n+1} = x_n + \frac{1}{6}\left(k_1 + 2k_2 + 2k_3 + k_4\right).
$$

**1.** Consider the ode given by

$$
\frac{dy}{dx} = f(x),
$$

with *y*(0) as the initial value. Use the standard fourth-order Runge-Kutta method to derive Simpson's rule.

# **Lecture 53 | Higher-order odes and systems**

[View this lecture on YouTube](https://youtu.be/wu1sS1g5ivs)

Our numerical methods can be easily adapted to solve higher-order odes, which are equivalent to a system of first-order odes. As an example, consider the second-order ode given by

$$
\ddot{x} = f(t, x, \dot{x}).
$$

This second-order ode can be rewritten as a system of two first-order odes by defining  $u = \dot{x}$ . We then obtain

$$
\dot{x} = u, \qquad \dot{u} = f(t, x, u).
$$

This trick also works for higher-order odes. For example, the third-order ode,

$$
\ddot{x} = f(t, x, \dot{x}, \ddot{x}),
$$

becomes

$$
\dot{x} = u, \qquad \dot{u} = v, \qquad \dot{v} = f(t, x, u, v).
$$

We can easily generalize the Runge-Kutta method to solve a system of first-order odes. As an example, we solve the following system of two first-order odes,

$$
\dot{x} = f(t, x, y), \quad \dot{y} = g(t, x, y).
$$

The generalization of the modified Euler method, say, would be

$$
k_1 = \Delta t f(t_n, x_n, y_n),
$$
  
\n
$$
l_1 = \Delta t g(t_n, x_n, y_n),
$$
  
\n
$$
k_2 = \Delta t f(t_n + \Delta t, x_n + k_1, y_n + l_1),
$$
  
\n
$$
l_2 = \Delta t g(t_n + \Delta t, x_n + k_1, y_n + l_1);
$$
  
\n
$$
x_{n+1} = x_n + \frac{1}{2} (k_1 + k_2),
$$
  
\n
$$
y_{n+1} = y_n + \frac{1}{2} (l_1 + l_2),
$$

where each stage requires computation of both a *k* and a *l*.

**1.** Write down the modified Euler method for the system of equations given by

$$
\dot{x} = f(t, x, y, z),
$$
  $\dot{y} = g(t, x, y, z),$   $\dot{z} = h(t, x, y, z).$ 

## **Lecture 54 | Adaptive Runge-Kutta methods**

#### [View this lecture on YouTube](https://youtu.be/6bCBXvsD7gw)

An adaptive ode solver automatically finds the best integration step-size ∆*t* at each time step. The Dormand-Prince method, which is implemented in Matlab's most widely used solver,  $\circ$ de45.m, determines the step size by comparing the results of fourth- and fifthorder Runge-Kutta methods. This solver requires six function evaluations per time step, and saves computational time by constructing both fourth- and fifth-order methods using the same function evaluations.

Suppose the fifth-order method finds  $x_{n+1}$  with local error  $O(\Delta t^6)$ , and the fourthorder method finds  $X_{n+1}$  with local error  $O(\Delta t^5)$ . Let *ε* be the requested error tolerance, and let *e* be the actual error. We can estimate *e* from the difference between the fifth- and fourth-order methods; that is,

$$
e=|x_{n+1}-X_{n+1}|.
$$

Now, we know that *e* is of O(∆*t* 5 ). Let ∆*τ* be the estimated step size required for the requested error *ε*. Then, by the scaling of the errors, we have

$$
e/\varepsilon = (\Delta t)^5/(\Delta \tau)^5,
$$

or solving for ∆*τ*,

$$
\Delta \tau = \Delta t \left(\frac{\varepsilon}{e}\right)^{1/5}.
$$

On the one hand, if the actual error is greater than the desired error,  $e > \varepsilon$ , then we reject the integration step and redo the time step using the smaller value ∆*τ*. On the other hand, if the actual error is less that the requested error,  $e < \varepsilon$ , then we accept  $x_{n+1}$  and increase the next time step to the larger value ∆*τ*. In practice, one usually includes a safety factor when computing ∆*τ*, i.e.,

$$
\Delta \tau = S \Delta t \left(\frac{\varepsilon}{e}\right)^{1/5},
$$

with  $S = 0.9$ , say, to prevent the wastefulness of too many failed time steps.

**1.** Using the Dormand-Prince method, suppose that a user requests an error tolerance of  $\varepsilon = 10^{-6}$ , and suppose the time step attempted was  $\Delta t = 0.01$  and that  $e = |x_{n+1} - x_{n+1}|$  $X_{n+1}| = 1.1 \times 10^{-6}$ . Is the current time step accepted? What time step will be used next? Assume a safety factor of 0.9.

## **Lecture 55 | Integrating odes in** MATLAB (Part A)

[View this lecture on YouTube](https://youtu.be/uyXXWxx_zIo)

The general purpose MATLAB ode solver is  $\triangleleft 45$ , and we discuss how to use it here. A single first-order ode is easily integrated using a technique similar to what we have already learned for both root finding and quadrature. For example, to numerically solve the logistic equation  $\dot{x} = rx(1 - x)$  and plot *x* versus *t*, we can code

```
dxdt = ( (x, r) r * x * (1-x);
r=1; x0=0.01; tspan=[0 10];
[t, x] = ode45(\theta(t, x) \text{dxdt}(x, r), tspan, x0);plot(t,x);
```
The right-hand side of the ode is defined as the function handle dxdt and this is passed as the first argument to ode45. Here, the right-hand side of the ode depends on *x* and a single parameter *r*, but in general, it can also depend on *t* and have any number of parameters. The integrator ode45, however, expects a function that has the independent variable as the first argument, and the dependent variables (written as a vector) as the second argument. To satisfy this requirement, we give  $\triangleleft$  e45 the handle to an anonymous function, so that its first argument looks like  $\mathcal{C}(t,x)$  dxdt  $(x, r)$ . For another example, if the right-hand side of the ode also depends on *t* and another parameter *s*, we may give as the first argument to ode45 the expression  $\mathfrak{g}(t,x)$  dxdt  $(t,x,r,s)$ .

For a more complicated system of odes, it may be more convenient to define a subfunction. For example, to solve the Lotka-Volterra equations,  $\dot{x} = rx(1-y)$ ,  $\dot{y} = \frac{1}{x}$  $\frac{1}{r}y(x-1)$ and plot *x* and *y* versus *t* as well as *y* versus *x*, we can code

```
r=1; x0=1.1; y0=1; tspan=[0 6*pi];
[t,xy]=ode45(@(t,xy) lv(xy,r),tspan,[x0 y0]);
figure(1); plot(t,xy); xlabel('t'); ylabel('x, y');
figure(2); plot(xy(:,1),xy(:,2)); xlabel('x'); ylabel('y');
function dxydt = lv(xy, r)x=xy(1); y=xy(2);dxydt=[rx*(1-y); (1/r)*y*(x-1)];
end
```
Some comments on the syntax. First, the independent variable *t* is returned as a column vector and the dependent variables  $x, y$  are returned as the two columns of a matrix. Second, the function  $\text{ode}45 \text{.m}$  knows we are solving a system of two first-order odes because we give it a two-component initial condition. Third, the right-hand sides of the differential equations dxydt must be written as a column vector. If they are written as a row vector, you will get an error message.

## **Lecture 56 | Integrating odes in** MATLAB (Part B)

[View this lecture on YouTube](https://youtu.be/YdhN_gGuPiA)

The graphs output by MATLAB from the Lotka-Volterra code are shown below:



A careful inspection of the figures reveals two anomalies. First, in the plot of *x* and *y* versus *t*, there are not enough data points in the output for a smooth plot, particularly at the maxima and minima of the functions where the linear interpolation of the data is evident. Second, in the phase-space plot of *y* versus *x*, it is unclear whether the solution is exactly periodic since overlapping curves are evident.

We can modify our call to  $\text{ode}45$  to improve these plots. The integrator  $\text{ode}45$  can do its own interpolation at little expense if you define tspan to have more than two elements. For example, we can replace  $tspan=[0 \ 6*pi]$  by the code  $tspan=lim space(0,6*pi,1000)$ .

After rerunning the code, you will observe that although the plot of *x* and *y* versus *t* is now smooth, the plot of *y* versus *x* still shows overlapping curves. The problem here is with the default error tolerance. There are two tolerances that are set: RelTol and AbsTol. Roughly speaking, RelTol sets the number of significant digits in the computed solution, and AbsTol sets the error tolerance as the solution approaches zero. The default values of these two tolerances are  $\text{ReLU} = 1 \cdot e^{-3}$ , corresponding to approximately three significant digits, and  $AbsTol=1.e-6$ . Here, the phase-space plot may be improved by decreasng RelTol. The appropriate change to the Matlab program is done in the following two lines of code:

options = odeset('RelTol',1.e-4);  $[t,xy] =$ ode45(@(t,xy) lv(xy,r),tspan,[x0; y0],options);

The improved plots are shown on the next page. This code runs fast, but it should be noted that for more computationally intensive calculations, smaller values for the error tolerances will result in longer computational times.



**1.** The Lorenz equations are a system of nonlinear odes that pioneered the study of chaos. The Lorenz equations are given by

$$
\dot{x} = \sigma(y - x), \qquad \dot{y} = x(r - z) - y, \qquad \dot{z} = xy - \beta z,
$$

where *σ*, *β* and *r* are constants. Edward Lorenz studied the solution for  $σ = 10$ ,  $β = 8/3$ and  $r = 28$ , and the result is now known as the Lorenz attractor, an example of what is now more generally known as a strange attractor. Compute the Lorenz attractor and plot *z* versus *x* and *y*. Use the MATLAB function plot3.m. Remove the transient before plotting.

# **Lecture 57 | Shooting method for boundary value problems**

[View this lecture on YouTube](https://youtu.be/qIfxydBEdzg)

We consider the general second-order ode of the form

$$
\frac{d^2y}{dx^2} = f(x, y, dy/dx),
$$

with the two-point boundary values  $y(x_0) = y_0$  and  $y(x_f) = y_f$ . To develop a solution method, we first formulate the second-order ode as a system of two first-order odes. We have

$$
\frac{dy}{dx} = z, \qquad \frac{dz}{dx} = f(x, y, z).
$$

One initial condition is known,  $y(x_0) = y_0$ , but the second initial condition for  $z(x_0)$  is unknown. The aim of the shooting method is to determine the value of  $z(x_0)$  that results in  $y(x_f) = y_f$ .

This is a root-finding problem. Let the variable  $\zeta$  denote the unknown value  $z(x_0)$ . We define  $F = F(\xi)$  as

$$
F(\xi) = y(x_f) - y_f,
$$

where  $y(x_f)$  is the value obtained by integrating the differential equations to  $x = x_f$  using the initial conditions  $y(x_0) = y_0$  and  $z(x_0) = \xi$ . The solution to  $F(\xi) = 0$  determines the value of  $\zeta = z(x_0)$  that solves the two-point boundary value problem. This method is called *shooting* because the slope of the function  $y = y(x)$  at  $x = x_0$  plays the role of aiming the gun, with the intention to hit the target at  $y(x_f) = y_f$ .

To determine the value of  $\xi$  that solves  $F(\xi) = 0$ , we can make use of the MATLAB function fzero, together with the differential equation solver ode45.

**1.** The dimensionless, unforced pendulum equation is given by

$$
\ddot{\theta} + \alpha \dot{\theta} + \sin \theta = 0,
$$

where  $\alpha$  is the only free parameter.

Consider initial conditions with the mass at the bottom,  $\theta(0) = 0$ . Using the shooting method, determine the smallest positive value of  $\dot{\theta}(0)$  such that the mass becomes exactly balanced at the top ( $\theta = \pi$ ). Plot this value of  $\dot{\theta}(0)$  versus  $\alpha$  for  $0 \le \alpha \le 2$ .

## **Lecture 58 | Project V: Two-body problem (Part A)**

[View this lecture on YouTube](https://youtu.be/lDKnPVCVzSQ)

Consider two masses  $m_1$  and  $m_2$  with position vectors  $\mathbf{r}_1$  and  $\mathbf{r}_2$ . Newton's second law and the universal law of gravitation yields for the governing equations

$$
m_1\ddot{\mathbf{r}}_1 = \mathbf{F}, \qquad m_2\ddot{\mathbf{r}}_2 = -\mathbf{F}, \quad \text{where} \quad \mathbf{F} = -Gm_1m_2\frac{\mathbf{r}_1 - \mathbf{r}_2}{|\mathbf{r}_1 - \mathbf{r}_2|^3},
$$

where *G* is the gravitational constant. The two-body problem can be reduced to a onebody problem by a change of coordinates. Let **r** be the relative position vector and **R** be the coordinate of the center of mass; that is,

$$
\mathbf{r}=\mathbf{r_1}-\mathbf{r_2},\qquad \mathbf{R}=\frac{m_1\mathbf{r_1}+m_2\mathbf{r_2}}{m_1+m_2}.
$$

Solving for **r<sup>1</sup>** and **r<sup>2</sup>** in terms of **r** and **R** gives

$$
r_1 = R + \frac{m_2r}{m_1 + m_2}, \qquad r_2 = R - \frac{m_1r}{m_1 + m_2}.
$$

Now, changing coordinates in Newton's equations from **r<sup>1</sup>** and **r<sup>2</sup>** to **r** and **R** results in

$$
m_1\ddot{\mathbf{R}} + \frac{m_1m_2\ddot{\mathbf{r}}}{m_1 + m_2} = \mathbf{F}, \qquad m_2\ddot{\mathbf{R}} - \frac{m_1m_2\ddot{\mathbf{r}}}{m_1 + m_2} = -\mathbf{F};
$$

and adding and subtracting these two equations yields the result

$$
\ddot{\mathbf{R}} = 0, \qquad \mu \ddot{\mathbf{r}} = \mathbf{F},
$$

where  $\mu$  is called the reduced mass and is given by

$$
\mu = \frac{m_1 m_2}{m_1 + m_2}.
$$

The center-of-mass coordinate system is an inertial coordinate system (one moving at constant velocity) with the center of mass fixed at the origin  $(\mathbf{R} = 0)$ . The remaining equation for the relative coordinate has the form of a one-body problem, and can be written as

$$
\ddot{\mathbf{r}} = -k\frac{\mathbf{r}}{r^3}, \qquad \text{where} \quad k = G(m_1 + m_2).
$$

Once **r** is determined, the coordinates of the two masses in the center of mass coordinate system can be found from

$$
\mathbf{r_1} = \left(\frac{m_2}{m_1 + m_2}\right)\mathbf{r}, \qquad \mathbf{r_2} = -\left(\frac{m_1}{m_1 + m_2}\right)\mathbf{r}.
$$

**1.** Consider the two-body problem where the solution for the relative coordinates is a circular orbit of unit radius, that is,

$$
\mathbf{r} = \cos(\omega t)\mathbf{i} + \sin(\omega t)\mathbf{j}.
$$

Sketch the orbits of  $m_1$  and  $m_2$  for (a)  $m_1 = m_2$  and (b)  $m_1 = 3m_2$ .

# **Lecture 59 | Project V: Two-body problem (Part B)**

#### [View this lecture on YouTube](https://youtu.be/1i4dE3cp_44)

The equivalent one-body problem can be simplified further. The position vector **r** and the velocity vector **˙r** form a plane, and since the acceleration vector points in the direction of the position vector, the one body motion will be restricted to this plane. We can place our coordinate system so that **r** lies in the *x*-*y* plane with  $z = 0$ , and we can write  $\mathbf{r} = x\mathbf{i} + y\mathbf{j}$ .

The relevant units are length and time, and they may be nondimensionalized using the constant parameter *k* (which has units  $l^3/t^2$ ) and the distance of closest approach of the two masses (the minimum value of *r*). The dimensionless governing equations, in component form, are then given by

$$
\ddot{x} = -\frac{x}{(x^2 + y^2)^{3/2}},
$$
  $\ddot{y} = -\frac{y}{(x^2 + y^2)^{3/2}}.$ 

We can further orient the *x*-*y* axes so that the solution  $\mathbf{r} = \mathbf{r}(t)$  is symmetric about the *x*-axis, and the minimum value of *r* occurs at  $x = -1$  and  $y = 0$ , where symmetry also requires  $\dot{x} = 0$ . Setting our initial conditions at the point of minimum  $r$ , we can write

$$
x(0) = -1
$$
,  $y(0) = 0$ ,  $\dot{x}(0) = 0$ ,  $\dot{y}(0) = \sqrt{1+e}$ ,

where we have parameterized  $\dot{y}(0)$  using  $e$ . The parameter  $e$  is called the eccentricity of the orbit, and some further analysis can show that closed elliptical orbits correspond to  $0 < e < 1$  and open hyperbolic orbits correspond to  $e > 1$ . The orbit with  $e = 0$  is circular and the orbit with  $e = 1$  is parabolic. Furthermore, the period of a closed orbit is given by  $T = 2\pi/(1-e)^{3/2}.$ 

For this project, you will write a code which solves the two-body problem. Given the initial conditions with specified value of  $e$ , the solution for  $\mathbf{r} = \mathbf{r}(t)$  for several periods of an orbit can be computed using the MATLAB function  $\text{ode45.m}$  by solving a system of four first-order differential equations. The motion of the corresponding two masses can then be simulated after specifying the ratio of the masses.

**1.** By solving a system of differential equations, determine the orbit of two masses using Newton's law and the universal law of gravitation. Display an animation of the orbit.

### **Week VI**

### **Partial Differential Equations**

*In this week's lectures, we learn how to find numerical solutions of some partial differential equations. This is a vast topic, and research areas such as computational fluid dynamics have many specialized solution methods.*

*Here, we only provide a taste of this subject. We divide the numerical solutions of pdes into boundary value problems and initial value problems, and apply the finite difference method of solution. We first show how to solve the Laplace equation, a boundary value problem. Two methods are illustrated: a direct method where the solution is found by Gaussian elimination; and an iterative method, where the solution is approached asymptotically. Second, we show how to solve the onedimensional diffusion equation, an initial value problem. The Crank-Nicolson method of solution is derived. We also show how to apply the Von Neumann stability analysis to determine the stability of our time-integration schemes.*

Your programming project will be to write a MATLAB code that computes the solution of the two*dimensional diffusion equation using the Crank-Nicolson method.*

# **Lecture 60 | Boundary and initial value problems**

#### [View this lecture on YouTube](https://youtu.be/tMO28AakkZ8)

A pde for which a numerical solution is sought most generally divides into two fundamental types. For a boundary value problem, the function or its derivatives are known on the boundary of a domain, and we want to solve for the function inside the domain. Usually, the solution is in some type of equilibrium or steady state. For an initial value problem, the function is known at  $t = 0$  throughout its spatial domain, and we want to know how the function evolves in time. Sometimes, it is possible to obtain the solution of a boundary value problem by evolving an initial value problem to a steady state. But it is often faster to solve the boundary value problem directly.

Schematics for a two-dimensional boundary value problem and one-dimensional initial value problem are shown below:



The black and red circles represent given boundary and initial values, respectively. The open circles represent the unknown solution.

# **Practice Quiz | Classify partial differential equations**

Classify the following pde problems:

- **1.** The steady flow of a fluid past an infinite cylinder.
	- *a*) Boundary value problem
	- *b*) Initial value problem
- **2.** The turbulent flow of a fluid past an infinite cylinder.
	- *a*) Boundary value problem
	- *b*) Initial value problem
- **3.** The steady-state salt concentration in a reservoir.
	- *a*) Boundary value problem
	- *b*) Initial value problem
- **4.** The transient salt concentration in a reservoir after a heavy rain.
	- *a*) Boundary value problem
	- *b*) Initial value problem
- **5.** The electrostatic potential surrounding two conductors held at constant potential.
	- *a*) Boundary value problem
	- *b*) Initial value problem
- **6.** The transitory electrostatic potential when a conductor is moved.
	- *a*) Boundary value problem
	- *b*) Initial value problem

#### **[Solutions to the Practice quiz](#page-206-1)**

# **Lecture 61 | Central difference approximation**

#### [View this lecture on YouTube](https://youtu.be/Tfo12ylAMso)

Finite difference methods are often used to solve pdes, and we need first to derive finite difference formulas for derivatives. To determine a formula for the derivative of a function *y* = *y*(*x*), consider the Taylor series for *y*(*x* + *h*) and *y*(*x* − *h*) about *x*:

$$
y(x+h) = y(x) + hy'(x) + \frac{1}{2}h^2y''(x) + \frac{1}{6}h^3y'''(x) + \frac{1}{24}h^4y''''(x) + ...,
$$
  

$$
y(x-h) = y(x) - hy'(x) + \frac{1}{2}h^2y''(x) - \frac{1}{6}h^3y'''(x) + \frac{1}{24}h^4y''''(x) + ....
$$

The standard definitions of the derivatives give the first-order approximations,

$$
y'(x) = \frac{y(x+h) - y(x)}{h} + O(h), \qquad y'(x) = \frac{y(x) - y(x-h)}{h} + O(h),
$$

where here we use the big-Oh notation, as in  $O(h^p)$ , to denote omitted terms that go to zero no slower than  $h^p$  as  $h$  goes to zero.

The more widely-used second-order approximation for the first derivative is called the central-difference approximation and is given by

$$
y'(x) = \frac{y(x+h) - y(x-h)}{2h} + O(h^2).
$$

The second-order approximation for the second derivative is also a central-difference approximation and can be found from considering

$$
y(x+h) + y(x-h) = 2y(x) + h2y''(x) + \frac{1}{12}h4y'''(x) + \dots,
$$

from which we obtain

$$
y''(x) = \frac{y(x+h) - 2y(x) + y(x-h)}{h^2} + O(h^2).
$$

**1.** Using Taylor series approximations for  $y(x + 2h)$ ,  $y(x + h)$ ,  $y(x - h)$  and  $y(x - 2h)$ , derive a central difference approximation for the first derivative  $y'(x)$  that is accurate to  $O(h^4)$ .

## **Lecture 62 | Discrete Laplace equation**

[View this lecture on YouTube](https://youtu.be/5-QAgZ81gg0)

As an example of a boundary value problem and the finite difference method, we discretize the two-dimensional Laplace equation,

$$
\left(\frac{\partial^2}{\partial x^2}+\frac{\partial^2}{\partial y^2}\right)\Phi=0,
$$

on the rectangular domain  $[0, L_x] \times [0, L_y]$ . We form a two-dimensional grid with  $N_x$  and  $N_y$  intervals in the *x*- and *y*-direction, respectively, and  $n_x = N_x + 1$  and  $n_y = N_y + 1$  grid points. With grid spacing  $\Delta x = L_x/N_x$  and  $\Delta y = L_y/N_y$ , the *x* and *y* coordinates of the grid points are given by

$$
x_i = (i-1)\Delta x, \quad i = 1, 2, ..., n_x;
$$
  $y_j = (j-1)\Delta y, \quad j = 1, 2, ..., n_y.$ 

Furthermore, we denote the value of  $\Phi(x_i, y_j)$  by  $\Phi_{i,j}$ . The central difference approximations for the second derivatives at the point  $(x_i, y_j)$  are given by

$$
\frac{\partial^2 \Phi}{\partial x^2} \approx \frac{1}{(\Delta x)^2} \left( \Phi_{i+1,j} - 2 \Phi_{i,j} + \Phi_{i-1,j} \right), \qquad \frac{\partial^2 \Phi}{\partial y^2} \approx \frac{1}{(\Delta y)^2} \left( \Phi_{i,j+1} - 2 \Phi_{i,j} + \Phi_{i,j-1} \right).
$$

We consider here only square grids with  $\Delta x = \Delta y$ . Combining the formulas for the second derivatives and rearranging terms results in the discrete Laplace equation given by

$$
4\Phi_{i,j} - \Phi_{i+1,j} - \Phi_{i-1,j} - \Phi_{i,j+1} - \Phi_{i,j-1} = 0.
$$

The five terms in this equation are shown in the figure on the right, where the red colored point is the central position at which the derivatives have been computed. A direct solution of the Laplace equation writes this difference equation in matrix form, with boundary conditions incorporated into the equations, and solves using Gaussian elimination.



**1.** Show that the solution of the discrete Laplace equation at grid point (*i*, *j*) on a uniform grid is just the average value of the solution at the neighboring four grid points.

### **Lecture 63 | Natural ordering**

[View this lecture on YouTube](https://youtu.be/0Wq2U_6ibSw)

To convert the discrete Laplace equation into matrix form, we will need to place the function values Φ*i*,*<sup>j</sup>* into a column vector. The commonly used ordering, called natural ordering, starts at the bottom left of the grid and moves left-to right along rows. This results in the column vector

$$
\Phi = [\Phi_{1,1}, \Phi_{2,1}, \ldots, \Phi_{n_x,1}, \Phi_{1,2}, \Phi_{2,2}, \ldots, \Phi_{n_x,2} \ldots, \Phi_{n_x,n_y}]^T,
$$

which has a total of  $n_x n_y$  components. The double index then maps into a single index, with mapping  $(i, j) \rightarrow k = i + (j - 1)n_x$ , and the discrete Laplace equation transforms from

> 4Φ*i*,*<sup>j</sup>* − Φ*i*+1,*<sup>j</sup>* − Φ*i*−1,*<sup>j</sup>* − Φ*i*,*j*+<sup>1</sup> − Φ*i*,*j*−<sup>1</sup> = 0 to  $4\Phi_k - \Phi_{k+1} - \Phi_{k-1} - \Phi_{k+n_k} - \Phi_{k-n_k} = 0.$

This equation is valid when *k* labels an interior point.

Boundary values for Φ need to satisfy prescribed boundary conditions. Here, we consider Dirichlet boundary conditions, where the values of  $\Phi$  itself are specified on the boundaries of the domain. The coordinates  $(i, j)$  on the boundaries are Bottom (B):  $j = 1$ ; Left (L):  $i = 1$ ; Top (T):  $j = n_y$ ; and Right (R):  $i = n_x$ ; and the ranges of the *k* indices (using the MATLAB colon operator), are given by

(B) 
$$
1 : n_x
$$
;  
\n(C)  $1 : n_x : 1 + (n_y - 1)n_x$ ;  
\n(D)  $1 : n_x : 1 + (n_y - 1)n_x$ ;  
\n(E)  $n_x : n_x : n_x n_y$ .

In the next lecture, we learn how to construct the matrix equation for Φ.
**1.** What are the *k* coordinates for the four corners of a rectangular box with  $n_x$  and  $n_y$ grid points in the *x*- and *y*-directions, respectively?

## **Lecture 64 | Matrix formulation**

[View this lecture on YouTube](https://youtu.be/NcVifhGXwrs)

To solve the discrete Laplace equation, given by

$$
4\Phi_k - \Phi_{k+1} - \Phi_{k-1} - \Phi_{k+n_x} - \Phi_{k-n_x} = 0,
$$

with  $\Phi_k$  specified on the boundaries, we must first write it as a matrix equation  $A\Phi = b$ , with A an  $n_x n_y$ -by- $n_x n_y$  square matrix.

When *k* indexes an internal grid point, the discrete Laplace equation goes into row *k* of matrix A with a zero in row *k* of vector b. When *k* indexes a boundary point, row *k* of matrix A will be row *k* of the identity matrix, and the known value of  $\Phi_k$  will go into row *k* of b. These boundary rows, then, serve the purpose of assigning the known boundary values to Φ.

For illustration, we construct the matrix equation for a three-by-three grid, shown at right, where the *k* indexing is labeled. Here we see that there is only one internal grid point (red) and eight boundary points (yellow), though usually there are many more internal points than boundary points. The discrete Laplace equation is only used in row five of the matrix, and the other rows are given by the corresponding rows of the identity matrix. When *k* indexes a boundary point, we assume that  $\Phi_k = b_k$  is known.



The matrix equation is given by

$$
\begin{pmatrix}\n1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & 0 & 1 & 0 & 0 & 0 & 0 & 0 & 0 \\
0 & -1 & 0 & -1 & 4 & -1 & 0 & -1 & 0 \\
0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 1 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 1\n\end{pmatrix}\n\begin{pmatrix}\n\Phi_1 \\
\Phi_2 \\
\Phi_3 \\
\Phi_4 \\
\Phi_5 \\
\Phi_6 \\
\Phi_7 \\
\Phi_8\n\end{pmatrix} =\n\begin{pmatrix}\nb_1 \\
b_2 \\
b_3 \\
b_4 \\
b_5 \\
b_6 \\
b_7 \\
b_8 \\
b_9\n\end{pmatrix}
$$

In general, the  $n_x n_y$ -by- $n_x n_y$  matrix A can be constructed using the MATLAB function spdiags, suitable for a sparse banded matrix. First A is constructed with five diagonals corresponding to the five terms in the discrete Laplace equation, and then the rows corresponding to boundary points can be replaced by the corresponding rows of the identity matrix. The right-hand side b has zeros in rows corresponding to interior points, and the boundary values of  $\Phi$  in rows corresponding to boundary points.

We will discuss the MATLAB code that solves the Laplace equation in the next Lecture.

**1.** Construct the matrix equation for the discrete Laplace equation on a four-by-four grid. When *k* indexes a boundary point, assume that  $\Phi_k = b_k$  is known.

**2.** On a rectangular grid with  $n_x$  and  $n_y$  grid points, how many interior points are there and how many boundary points? What percentage of grid points are boundary points when  $n_x = n_y = 100$ , and what percentage when  $n_x = n_y = 1000$ ?

# **Lecture 65 |** Matlab **solution of the Laplace equation (direct method)**

#### [View this lecture on YouTube](https://youtu.be/kl3c7XGRN9g)

We solve the Laplace equation in a rectangle with side lengths  $Lx$  and  $Ly$ . The resolution is set by the parameters  $Nx$  and  $Ny$  with  $nx=Nx+1$  and  $ny=Ny+1$  grid points in each direction. The accuracy of the solution is improved by increasing the values of  $Nx$  and  $Ny$ . The beginning of the code defines the rectangle and the grid:

```
Lx=1; Ly=1; %rectangle dimensions
Nx=100; Ny=100; %number of intervals in x, y directions
nx=Nx+1; ny=Ny+1; %number of gridpoints in x,y directions
dx=Lx/Nx; dy=Ly/Ny; %grid length in x,y directions
x=(0:Nx)*dx; y=(0:Ny)*dy; %x, y values on the grid
```
Next, the A matrix is constructed. The *k* indices of the boundary points are defined in the vector boundary\_index, the five diagonals are placed in the A matrix using the Matlab spdiags function, and the rows associated with the boundary points are replaced by the corresponding rows of the identity matrix using speye:

```
boundary_index=[1:nx, 1:nx:1+(ny-1)*nx, 1+(ny-1)*nx:nx*ny, nx:nx:nx*ny];
diagonals = [4 \times \text{ones}(nx \times ny, 1), -ones(nx \times ny, 4)];
A=spdiags(diagonals,[0 -1 1 -nx nx], nx *ny, nx *ny;
I=speye(nx*ny);
A(boundary_index,:)=I(boundary_index,:);
```
We next construct the right-hand side of the matrix equation. We use the boundary conditions  $\Phi = 0$  on the bottom, left and right sides of the square, and  $\Phi = 4x(1-x)$ at the top. Since  $0 \le x \le 1$ ,  $\Phi$  will equal zero on the corners and one at the midpoint of the top. The simplest and clearest way to construct  $\triangleright$  is to first use  $(i, j)$  indexing and then turn  $\mathbf b$  into a column vector using the MATLAB reshape function (or, more simply, the colon operator):

```
b=zeros(nx,ny); %interior rows are zero
b(:,1)=0; % bottomb(1,:)=0; %letb(:, ny) = 4 \times x. *(1-x); % topb(nx,:)=0; \text{'right}b=reshape(b,nx*ny,1); \frac{1}{2} and \frac{1}{2} column vector using natural ordering [same as b=b(:)]
```
We solve for  $\Phi$  and return to  $(i, j)$  indexing using reshape:

Phi=A\b; %solution by Gaussian elimination Phi=reshape(Phi,nx,ny); %back to (i,j) indexing

Finally, we plot the solution. We use the MATLAB meshgrid to set up the grid points, and choose the contour levels we want to plot in the vector  $v$ . The contour plot is made using the Matlab function contour. We use axis equal so that the graph has the same aspect ratio as the problem domain.

To understand the required data structures of Phi, meshgrid and contour, suppose that the *x*-*y* grid is defined by  $x=[x1 \ x2 \ x3]$  and  $y=[y1 \ y2 \ y3]$ . Then  $[X, Y]$ =meshgrid(x, y) results in the matrices

$$
X = \begin{bmatrix} x1 & x2 & x3 \\ x1 & x2 & x3 \\ x1 & x2 & x3 \end{bmatrix}, \quad Y = \begin{bmatrix} y1 & y1 & y1 \\ y2 & y2 & y2 \\ y3 & y3 & y3 \end{bmatrix},
$$

which are used to specify every point on the grid. The  $(i, j)$  data structure for  $\Phi$ , where the first index gives the *x*-location and the second index gives the *y*-location, is given by

$$
\text{Phi} = \begin{bmatrix} \text{Phi111} & \text{Phi12} & \text{Phi13} \\ \text{Phi21} & \text{Phi22} & \text{Phi23} \\ \text{Phi31} & \text{Phi32} & \text{Phi33} \end{bmatrix}.
$$

So when calling contour, to match the values of  $Phi$  with the locations in the  $X$ ,  $Y$  grid, we need to use the transpose, that is,



The code which plots the contours is then written as

```
[X, Y] =meshqrid(x, y);
v=[0.8 0.6 0.4 0.2 0.1 0.05 0.01]; %contour levels
contour(X,Y,Phi',v,'ShowText','on'); %requires transpose
axis equal;
set(gca,'YTick',[0 0.2 0.4 0.6 0.8 1]);
set(gca,'XTick',[0 0.2 0.4 0.6 0.8 1]);
xlabel('$x$','Interpreter','latex','FontSize',14 );
ylabel('$y$','Interpreter','latex','FontSize',14);
title('Solution of the Laplace equation','Interpreter','latex','FontSize',16);
```
The resulting plot looks like this:



**1.** Using the direct method, solve the Laplace equation inside a unit square. Set the boundary conditions to be zero on the left and bottom sides, and to go from zero to one across the top, and from one to zero down the right side. Model these boundary conditions as

 $\Phi = x(2-x)$  for  $y=1$ ;  $\Phi = y(2-y)$  for  $x=1$ .

# **Lecture 66 | Jacobi, Gauss-Seidel and SOR methods**

#### [View this lecture on YouTube](https://youtu.be/QpzOttega9s)

Iterative methods are often used for solving systems of equations arising from pdes. Here, we explain the related iterative methods that go by the names of Jacobi, Gauss-Seidel, and Successive Over Relaxation (SOR) methods.

For illustration, we consider again the discrete Laplace equation given by

$$
4\Phi_{i,j} - \Phi_{i+1,j} - \Phi_{i-1,j} - \Phi_{i,j+1} - \Phi_{i,j-1} = 0.
$$

The Jacobi method simply solves this equation iteratively for Φ*i*,*<sup>j</sup>* . Adding superscripts to denote the iteration step, we iterate the following equation until convergence:

$$
\Phi_{i,j}^{(n+1)} = \frac{1}{4} \left( \Phi_{i+1,j}^{(n)} + \Phi_{i-1,j}^{(n)} + \Phi_{i,j+1}^{(n)} + \Phi_{i,j-1}^{(n)} \right).
$$

Historically, the FORTRAN language required two different arrays to implement the Jacobi method: the first to store the current solution and the second to store the updated solution. When only one array was used, Φ*i*−1,*<sup>j</sup>* and Φ*i*,*j*−<sup>1</sup> were updated before Φ*i*,*<sup>j</sup>* and the method was instead called Gauss-Seidel. Mathematically, it looked like

$$
\Phi_{i,j}^{(n+1)} = \frac{1}{4} \left( \Phi_{i+1,j}^{(n)} + \Phi_{i-1,j}^{(n+1)} + \Phi_{i,j+1}^{(n)} + \Phi_{i,j-1}^{(n+1)} \right).
$$

The SOR method first rewrites the Jacobi (or Gauss-Seidel) method by adding and subtracting Φ (*n*)  $\int_{i,j}^{(n)}$  on the right-hand side to obtain

$$
\Phi_{i,j}^{(n+1)} = \Phi_{i,j}^{(n)} + \frac{1}{4} \left( \Phi_{i+1,j}^{(n)} + \Phi_{i-1,j}^{(n)} + \Phi_{i,j+1}^{(n)} + \Phi_{i,j-1}^{(n)} - 4\Phi_{i,j}^{(n)} \right).
$$

In this form, we see that the Jacobi method updates the value of Φ*i*,*<sup>j</sup>* at each iteration. We can either magnify or diminish this update by multiplying the correction by *λ*. After a bit of algebra, we find

$$
\Phi_{i,j}^{(n+1)} = (1-\lambda)\Phi_{i,j}^{(n)} + \frac{\lambda}{4} \left( \Phi_{i+1,j}^{(n)} + \Phi_{i-1,j}^{(n)} + \Phi_{i,j+1}^{(n)} + \Phi_{i,j-1}^{(n)} \right).
$$

When 1 < *λ* < 2, the method is called *Successive Over Relaxation*, or SOR, and convergence may be accelerated. For some nonlinear equations, however, it may be necessary to choose  $0 < \lambda < 1$  for stability reasons, and perhaps a better name for the method would be *Successive Under Relaxation*. When under relaxing, the convergence of the iteration will be slowed, but that is the price that sometimes must be paid.

**1.** The Jacobi, Gauss-Seidel and SOR methods can also be used to solve a system of linear equations. Consider the system of equations given by

$$
a_{11}x_1 + a_{12}x_2 + a_{13}x_3 = b_1,
$$
  
\n
$$
a_{21}x_1 + a_{22}x_2 + a_{23}x_3 = b_2,
$$
  
\n
$$
a_{31}x_1 + a_{32}x_2 + a_{33}x_3 = b_3.
$$

By solving the *i*th equation for *x<sup>i</sup>* , write down the Jacobi iteration method for this system.

# **Lecture 67 | Red-black ordering**

[View this lecture on YouTube](https://youtu.be/giTZ89q-Bpk)

A final variant of the classical iteration methods uses red-black ordering of the grid points. In this algorithm, the grid is viewed as a checkerboard with alternating red and black grid points, as shown below.



An updating of Φ*i*,*<sup>j</sup>* is done in two passes: in the first pass, Φ*i*,*<sup>j</sup>* is updated only at the red grid points; in the second pass, only on the black grid points. Because of the structure of the discrete Laplace equation, the updated values of  $\Phi$  on the red squares depend only on the values of  $\Phi$  on the black squares, and the updated values of  $\Phi$  on the black squares depend only on the values of  $\Phi$  on the red squares. All updating terms, then, are at the same level of iteration. This could result in faster and more stable convergence, or may be useful when implementing parallel programming algorithms.

# **Lecture 68 |** Matlab **solution of the Laplace equation (iterative method)**

[View this lecture on YouTube](https://youtu.be/vz-enxt57yY)

The code begins in the same way as the direct method:

```
Lx=1; Ly=1; %rectangle dimensions
Nx=100; Ny=100; %number of intervals in x,y directions
nx=Nx+1; ny=Ny+1; %number of gridpoints in x,y directions
dx=Lx/Nx; dy=Ly/Ny; %grid length in x,y directions
x=(0:Nx)*dx; y=(0:Ny)*dy; %x, y values on the grid
```
Next, some additional parameters and index vectors are defined:

```
eps=1.e-6; %convergence criteria for each value of Phi
index_x=2:nx-1; index_y=2:ny-1; %internal grid points
```
The solution matrix  $\Phi$  is initialized with zero at all internal grid points, and with the boundary values fixed:

```
Phi=zeros(nx, ny); % matrix with solution and boundary conditions
Phi(:,1)=0; Phi(1,:)=0; Phi(:, ny)=4*x.*(1-x); Phi(nx,:)=0; %boundary grid points BLTR
```
We now iterate the Jacobi method until each internal value of  $\Phi$  changes less than eps with each iteration.

```
Phi_old=Phi;
error=2*eps; ncount=0;
while (error > eps)
   ncount=ncount+1;
   Phi(index_x,index_y)=0.25*(Phi(index_x+1,index_y) \ldots+Phi(index_x-1,index_y)+Phi(index_x,index_y+1)+Phi(index_x,index_y-1));
   error=max(abs(Phi(:)-Phi_old(:)));
   Phi_old=Phi;
end
```
Finally, we plot the solution as we did for the direct method.

```
[X, Y] =meshqrid(x, y);
v=[0.8 0.6 0.4 0.2 0.1 0.05 0.01];
[C, h] =contour(X, Y, Phi', v); clabel(C, h);
axis equal;
set(gca, 'YTick', [0 0.2 0.4 0.6 0.8 1]);
set(gca, 'XTick', [0 0.2 0.4 0.6 0.8 1]);
xlabel('$x$','Interpreter','latex','FontSize',14 );
ylabel('$y$','Interpreter','latex','FontSize',14);
title('Solution of the Laplace equation','Interpreter','latex','FontSize',16);
```
**1.** Using the Jacobi method, solve the Laplace equation inside a unit square. Set the boundary conditions to be zero on the left and bottom sides, and to go from zero to one across the top, and from one to zero down the right side. Model these boundary conditions as

 $\Phi = x(2 - x)$  for  $y = 1$ ;  $\Phi = y(2 - y)$  for  $x = 1$ .

# **Lecture 69 | Explicit methods for solving the diffusion equation**

[View this lecture on YouTube](https://youtu.be/MEqYYjGcnww)

The initial value problem considered here is the one-dimensional diffusion equation for  $u = u(x, t)$ , given by

$$
\frac{\partial u}{\partial t} = D \frac{\partial^2 u}{\partial x^2},
$$

with boundary conditions  $u(-L_x,t) = u(L_x,t) = 0$  and initial conditions  $u(x,0) = u_0(x)$ .

We divide the line into  $N_x$  intervals using  $n_x = N_x + 1$  grid points. With grid spacing  $\Delta x = 2L_x/N_x$ , the *x* coordinates of the grid points are given by

$$
x_j = -L_x + (j-1)\Delta x
$$
, for  $j = 1, 2, ..., n_x$ .

Time is also discretized, and with time step  $\Delta t$ , the solution is computed at the times

$$
t_l = (l-1)\Delta t, \quad \text{for } l = 1, 2, \ldots.
$$

We will denote the value of  $u(x_j, t_l)$  by  $u_j^l$ . Here, we use the convention that subscripts on *u* denote the spacial grid point and superscripts on *u* denote the time step.

The Forward Time Centered Space (FTCS) discretization, with the relevant grid points shown on the right, uses the second-order central difference approximation for the second derivative and the first-order Euler method for the time integration.



The discrete diffusion equation then becomes

$$
u_j^{l+1} = u_j^l + \frac{\Delta t D}{(\Delta x)^2} \left( u_{j+1}^l - 2u_j^l + u_{j-1}^l \right), \quad \text{for } j = 2, 3, ..., n_x - 1 \text{ and } l = 1, 2, ...
$$

This type of method is called explicit because the solution at the time step *l* + 1 can be written explicitly in terms of the solution at the earlier time step *l*. Higher-order timestepping methods can also be formulated using Runge-Kutta methods. We shall see in the next lecture, however, that stability issues may be more important than the order of the method.

**1.** Use the second-order Runge-Kutta method known as the modified Euler method to write a two-step process for solving the one-dimensional diffusion equation.

**2.** Consider the one-dimensional advection equation given by

$$
\frac{\partial u}{\partial t} = -c \frac{\partial u}{\partial x}.
$$

Using the second-order central difference approximation for the spatial derivative and the first-order Euler method for the time integration, derive the FTCS scheme for the advection equation.

# **Lecture 70 | Von Neumann stability analysis**

[View this lecture on YouTube](https://youtu.be/QUiUGNwNNmo)

We will analyze the stability of the FTCS scheme for the diffusion equation, given by

$$
u_j^{l+1} = u_j^l + \frac{\Delta t D}{(\Delta x)^2} \left( u_{j+1}^l - 2u_j^l + u_{j-1}^l \right).
$$

We look for solutions of the form

$$
u_j^l = \xi^l e^{ikj\Delta x},
$$

where  $i =$ √ −1; *k* is called the wavenumber of the mode and can be any value; *ξ* (here raised to the *l*th power) is the unknown parameter that we will try to determine. If we find that  $|\xi| > 1$  for any value of the wavenumber *k*, then we say that the scheme is unstable.

Substitution into the FTCS scheme for the diffusion equation results in

$$
\xi^{l+1}e^{ikj\Delta x} = \xi^l e^{ikj\Delta x} + \frac{\Delta t D}{(\Delta x)^2} \left( \xi^l e^{ik(j+1)\Delta x} - 2\xi^l e^{ikj\Delta x} + \xi^l e^{ik(j-1)\Delta x} \right).
$$

Dividing by *ξ l e ikj*∆*x* leads to

$$
\xi = 1 + \frac{\Delta t D}{(\Delta x)^2} \left( e^{ik\Delta x} - 2 + e^{-ik\Delta x} \right).
$$

Since  $e^{ik\Delta x} + e^{-ik\Delta x} = 2\cos(k\Delta x)$ , we obtain

$$
\xi = 1 + \frac{2\Delta t D}{(\Delta x)^2} (\cos (k\Delta x) - 1).
$$

The value of  $|\xi|$  is largest when  $\cos(k\Delta x) = -1$ , and the scheme is unstable when

$$
1 - \frac{4\Delta t D}{(\Delta x)^2} < -1, \qquad \text{or} \qquad \Delta t > \frac{(\Delta x)^2}{2D}.
$$

When a numerical scheme is unstable, the solution blows up in time, which is obviously an unphysical solution of the diffusion equation. Using the FTCS scheme and other explicit methods, stability considerations limit the maximum time step that can be used. We will see that implicit methods will remove this constraint.

**1.** Analyze the stability of the FTCS scheme for the advection equation, given by

$$
u_j^{l+1} = u_j^l - \frac{c\Delta t}{2\Delta x} \left( u_{j+1}^l - u_{j-1}^l \right).
$$

# **Lecture 71 | Implicit methods for solving the diffusion equation**

[View this lecture on YouTube](https://youtu.be/DgyGoLxWmcg)

Instead of computing the spatial derivative at the time *t<sup>l</sup>* , we compute it at the forward time  $t_{l+1}$  (see the grid points on the right). The discrete diffusion equation is now



$$
u_j^{l+1} = u_j^l + \frac{\Delta t D}{(\Delta x)^2} \left( u_{j+1}^{l+1} - 2u_j^{l+1} + u_{j-1}^{l+1} \right).
$$

This method is called implicit because finding the solution at the time step  $l + 1$  requires solving a system of equations.

We can once again perform a stability analysis. With our ansatz

$$
u_j^l = \xi^l e^{ikj\Delta x},
$$

the implicit discrete diffusion equation results in

$$
\xi^{l+1}e^{ikj\Delta x} = \xi^l e^{ikj\Delta x} + \frac{\Delta t D}{(\Delta x)^2} \left( \xi^{l+1}e^{ik(j+1)\Delta x} - 2\xi^{l+1}e^{ikj\Delta x} + \xi^{l+1}e^{ik(j-1)\Delta x} \right).
$$

Similar to before, dividing by *ξ l e ikj*∆*x* then leads to

$$
\xi = 1 + \frac{2\Delta t D\xi}{(\Delta x)^2} (\cos(k\Delta x) - 1),
$$

where there is now a factor of *ξ* on the right-hand side. Solving for *ξ*, we obtain

$$
\xi = \frac{1}{1 + \frac{2\Delta t D}{(\Delta x)^2} (1 - \cos(k\Delta x))},
$$

which has magnitude always less than or equal to one. The implicit method is unconditionally stable.

**1.** Consider the one-dimensional advection equation given by

$$
\frac{\partial u}{\partial t} = -c \frac{\partial u}{\partial x}.
$$

- *a*) By computing the spatial derivative at the advanced time step  $t_{l+1}$ , derive the implicit discrete advection equation.
- *b*) Analyze its stability.

# **Lecture 72 | Crank-Nicolson method for the diffusion equation**

[View this lecture on YouTube](https://youtu.be/f_JZRjt8AZ4)

The most popular discretization of the one-dimensional diffusion equation is called the Crank-Nicolson method and is a stable implicit method that also has the advantage of being second-order in both space and time. To obtain Crank-Nicolson, we average the previous spacial discretizations at the times  $t_l$  and  $t_{l+1}$  to obtain a central difference approximation for the time derivative of *u* at the time  $t_{l+1/2}$ . The discrete diffusion equation becomes

$$
u_j^{l+1} = u_j^l + \frac{\Delta t D}{2(\Delta x)^2} \left( \left( u_{j+1}^l - 2u_j^l + u_{j-1}^l \right) + \left( u_{j+1}^{l+1} - 2u_j^{l+1} + u_{j-1}^{l+1} \right) \right).
$$

If we define  $\alpha = \Delta t D / (\Delta x)^2$  and multiply the equation by two, we can write the Crank-Nicolson method as

$$
-\alpha u_{j+1}^{l+1} + 2(1+\alpha)u_j^{l+1} - \alpha u_{j-1}^{l+1} = \alpha u_{j+1}^l + 2(1-\alpha)u_j^l + \alpha u_{j-1}^l,
$$

which leads to a tridiagonal matrix equation. For example, with spatial boundary conditions given by  $u_1^l = u_{n_x}^l = 0$ , and for  $n_x = 4$ , the matrix equation including the boundary conditions is given by

$$
\begin{pmatrix} 1 & 0 & 0 & 0 \ -\alpha & 2(1+\alpha) & -\alpha & 0 \ 0 & -\alpha & 2(1+\alpha) & -\alpha \ 0 & 0 & 0 & 1 \ \end{pmatrix} \begin{pmatrix} u_1^{l+1} \\ u_2^{l+1} \\ u_3^{l+1} \\ u_4^{l+1} \end{pmatrix} = \begin{pmatrix} 0 \\ \alpha u_1^l + 2(1-\alpha)u_2^l + \alpha u_3^l \\ \alpha u_2^l + 2(1-\alpha)u_3^l + \alpha u_4^l \\ 0 \end{pmatrix}.
$$

Although  $u_1^l$  and  $u_4^l$  are zero on the right-hand side, we include them here for ease of coding. The  $n_x$ -by- $n_x$  matrix is independent of time step, and the right-hand side is continuously updated as the solution evolves.

**1.** Consider the one-dimensional advection equation given by

$$
\frac{\partial u}{\partial t} = -c \frac{\partial u}{\partial x}.
$$

The explicit Lax scheme for the advection equation is given by

$$
u_j^{l+1} = \frac{1}{2}(u_{j+1}^l + u_{j-1}^l) - \frac{c\Delta t}{2\Delta x}(u_{j+1}^l - u_{j-1}^l).
$$

Analyze its stability and derive the Courant-Friedrichs-Lewy (CFL) stability criterion, which is widely used in fluid turbulence simulations.

# **Lecture 73 |** Matlab **solution of the diffusion equation**

#### [View this lecture on YouTube](https://youtu.be/S3DXGvrdx1w)

We solve the one-dimensional diffusion equation for  $|x| \leq L_x$  with *u* equal zero on the boundaries. We will use the Crank-Nicolson method. The length is set by the parameter  $Lx$ , and the number of spatial intervals is set by the parameter  $Nx$ . The number of grid points is  $nx=Nx+1$ . The spatial grid size is  $dx$  and the time-step is  $dt$ . Here we set the time-step at the stability limit for the FTCS scheme.

```
D=1; Lx=1; %Diffusion constant; Domain length
nsteps=10000; %number of time steps
nout=500; %plot every nout time steps
Nx=500; %number of intervals
nx=Nx+1;%number of gridpoints in x direction including boundaries
dx=2*Lx/Nx; %grid size in x
x=-Lx + (0:Nx)*dx; %x values on the grid
dt=(dx)^2/(2*D); %borderline stability of FTCS scheme
```
Next, the nx-by-nx matrix is constructed using spdiags. The top and bottom boundary rows are replaced by the corresponding rows of the identity matrix:

```
alpha=dt*D/dx^2;
diagonals = [2*(1+a1pha)*ones(nx,1), -alpha*ones(nx,2)];
A=spdiags(diagonals, [0 -1 1], nx, nx);
I=speye(nx);
A([1 nx],:)=I([1 nx],:); %boundaries
```
We next initialize the solution and plot the initial conditions. We choose a sharply peaked Gaussian centered at  $x = 0$ . The values of *u* on the boundaries will be close enough to zero and will not change. sigma=Lx/16;

```
u=1/(sigma*sqrt(2*pi))*exp(-0.5*(x/sigma).^2); u=u';
plot(x,u); hold on;
xlabel('$x$','Interpreter','latex','FontSize',14);
ylabel('$u(x, t)$','Interpreter','latex','FontSize',14);
title('Solution of the diffusion equation','Interpreter','latex','FontSize',16);
```
Finally, we evolve the diffusion equation by constructing the right-hand sides and solving using the backslash operator. The solution is plotted every nout time steps.

```
for m=1:nsteps
   b=[0; [alpha*u(1:nx-2) + 2*(1-alpha)*u(2:nx-1) + alpha*u(3:nx)]; 0];u=A\b;
   if mod(m, nout) == 0, plot(x, u), end
end
```
See the next page for the resulting plot.



**1.** *a*) Derive a second-order method for the *x*-derivative at boundary points. When *x* is a boundary point on the left, use the Taylor series

$$
y(x+h) = y(x) + hy'(x) + \frac{1}{2}h^2y''(x) + O(h^3),
$$
  

$$
y(x+2h) = y(x) + 2hy'(x) + 2h^2y''(x) + O(h^3).
$$

When  $x$  is a boundary point on the right, use the Taylor series

$$
y(x - h) = y(x) - hy'(x) + \frac{1}{2}h^2y''(x) + O(h^3),
$$
  

$$
y(x - 2h) = y(x) - 2hy'(x) + 2h^2y''(x) + O(h^3).
$$

- *b*) No-flux boundary conditions sets *∂u*/*∂x* equal to zero on the boundaries. Using the results of Part a), determine boundary conditions for  $u_1^l$  and  $u_{n_x}^l$ .
- **2.** Solve the one-dimensional diffusion equation for  $|x| \leq L$  using the Crank-Nicolson method. Assume no flux boundary conditions. Use the results of 1 (b).

# **Lecture 74 | Project VI: Twodimensional diffusion equation**

#### [View this lecture on YouTube](https://youtu.be/-DJJqgr6HdE)

Your final project is to solve the two-dimensional diffusion equation. You will combine both the direct method for the Laplace equation and the Crank-Nicolson method.

The two-dimensional diffusion equation for  $u = u(x, y, t)$  on a square of side 2*L* is given by

$$
\frac{\partial u}{\partial t} = D \left( \frac{\partial^2 u}{\partial x^2} + \frac{\partial^2 u}{\partial y^2} \right),
$$

where we will assume boundary conditions  $u(\pm L, \gamma, t) = u(x, \pm L, t) = 0$ .

We form a two-dimensional spatial grid with *N* intervals and  $n = N + 1$  grid points in both the *x*- and *y*-directions. With grid spacing *h* = 2*L*/*N*, the *x* and *y* coordinates of the grid points are given by

$$
x_i = -L + (i - 1)h
$$
,  $y_j = -L + (j - 1)h$ ,

for  $i, j = 1, 2, \ldots, n$ . With time step given by  $\Delta t$ , we have

$$
t_l = (l-1)\Delta t, \quad \text{for } l = 1, 2, \ldots;
$$

and we denote the value of  $u(x_i, y_j, t_l)$  by  $u_{i,j}^l$ .

The Crank-Nicolson method applied to the two-dimensional diffusion equation results in

$$
u_{i,j}^{l+1} = u_{i,j}^l + \frac{\Delta t D}{2h^2} \left( \left( u_{i+1,j}^l + u_{i-1,j}^l + u_{i,j+1}^l + u_{i,j-1}^l - 4u_{i,j}^l \right) + \left( u_{i+1,j}^{l+1} + u_{i-1,j}^{l+1} + u_{i,j+1}^{l+1} + u_{i,j-1}^{l+1} - 4u_{i,j}^{l+1} \right) \right).
$$

We define  $\alpha = \Delta t D / h^2$ , multiply by two and rearrange terms to obtain

$$
2(1+2\alpha)u_{i,j}^{l+1} - \alpha \left(u_{i+1,j}^{l+1} + u_{i-1,j}^{l+1} + u_{i,j+1}^{l+1}\right) = 2(1-2\alpha)u_{i,j}^{l} + \alpha \left(u_{i+1,j}^{l} + u_{i-1,j}^{l} + u_{i,j+1}^{l} + u_{i,j-1}^{l}\right).
$$

Finally, applying natural ordering with the mapping  $(i, j) \rightarrow k = i + (j - 1)n$ , we obtain

the equation

$$
2(1+2\alpha)u_k^{l+1} - \alpha \left(u_{k+1}^{l+1} + u_{k-1}^{l+1} + u_{k+n}^{l+1} + u_{k-n}^{l+1}\right)
$$
  
= 2(1-2\alpha)u\_k^l + \alpha \left(u\_{k+1}^l + u\_{k-1}^l + u\_{k+n}^l + u\_{k-n}^l\right).

This is a matrix equation with a time-independent banded matrix with five diagonals. The right-hand side is time dependent. Since you will need to solve a large  $n^2$ -by- $n^2$  matrix equation at each time step, the LU decomposition of the matrix can be performed once and then used to speed-up the code.

**1.** Solve the two-dimensional diffusion equation on a square with *u* equal to zero on the boundaries.

# **Problem solutions and MATLAB learner templates**

**1.** A binary number rounded to six places after the binary point will round up if the truncated number is greater than  $2^{-7}$  and will round down if the truncated number is less than 2−<sup>7</sup> . We have

$$
\frac{1}{3} = 0.010101\overline{01} \approx 0.010101, \quad \frac{1}{5} = 0.001100\overline{1100} \approx 0.001101, \quad \frac{1}{6} = 0.001010\overline{10} \approx 0.001011,
$$

$$
\frac{1}{7} = 0.001001\overline{001} \approx 0.001001, \quad \frac{1}{9} = 0.000111\overline{000111} \approx 0.000111.
$$

The numbers  $1/3$ ,  $1/7$  and  $1/9$  round down and the numbers  $1/5$  and  $1/6$  round up.

### **Solutions to the [Problems for Lecture 2](#page-11-0)**

**1.** Using decimal notaton, the number 1 corresponds to  $s = 0$ ,  $e = 1023$  and  $f = 0$ . The binary for 1023 can be found from  $1023 = 2^{10} - 1$ , so that  $e = 100\,0000\,0000 000\,0000\,0001 = 011\,1111\,111$ , resulting in the double precision format for 1 to be

$$
0011\,1111\,1111\,0000\ldots 0000.
$$

This number is usually written in hex, which replaces every four bits by a numeral representing the decimal numbers  $0 - 15$ . The numerals for  $0 - 9$  are retained, and those for 10 − 15 are written using the letters a, b, ..., f. We would therefore write the computer number representing 1 as

#### 3ff0 0000 0000 0000.

The number  $1/2$  corresponds to  $s = 0$ ,  $e = 1022$  and  $f = 0$ . The binary for 1022 is  $e = 011 1111 1110$ , resulting in the double precision format for  $1/2$  to be

#### 0011 1111 1110 0000 . . . 0000,

which in hex is given by

$$
3fe0\,0000\,0000\,0000.
$$

To determine the computer number for 1/3, we first write

$$
\frac{1}{3} = \frac{1}{4} \left( 1 + \frac{1}{3} \right).
$$

We can observe that  $e = 1021$  and  $f = 1/3$ , which in binary becomes  $e = 01111111101$ and  $f = 0101 \ldots 0101$ , where we have rounded down. We would therefore write the computer number representing 1/3 as

#### 0011 1111 1101 0101 0101 . . . 0101,

which in hex is given by

3fd5 5555 5555 5555.

**2.** The largest number corresponds to  $s = 0$ ,  $e = 2^{11} - 2 - 1023 = 1023$ , and  $1.f =$  $1 + (1 - 2^{-52})$ . Therefore realmax =  $2^{1023}(2-2^{-52}) = 1.7977e+308.$ 

**3.** The smallest positive normal number corresponds to  $s = 0$ ,  $e = 1$ , and  $1.f = 1 + 0$ . Therefore

 $\text{realmin} = 2^{-1022} = 2.2251e-308.$ 

**4.** The number one is represented as  $(-1)^{0} \times 2^{0} \times 1.0$ . The next largest number is  $(-1)^{0} \times 2^{0} \times (1+2^{-52}$ . Therefore, eps =  $2^{-52}$  = 2.2204e-16.

### **Solutions to the [Problems for Lecture 3](#page-13-0)**

#### **1.**

```
a) (sqrt(5)-1)/2
  0.6180
```
- *b*) (((sqrt(5)+1)/2)^10 ((sqrt(5)-1)/2 )^10)/sqrt(5) 55.0000 (Try subtracting 55 from this number.)
- *c*) 2^5/(2^5-1) 1.0323
- *d*)  $(1-1/2^5)^{-1}$ 1.0323
- *e*) exp(1)^3 20.0855
- *f*) log(exp(1)^3) 3.0000 (Try subtracting 3 from this number.)
- *g*) sin(pi/6) 0.5000 (Try subtracting 0.5 from this number.)
- *h*) cos(pi/6) 0.8660
- *i*)  $\sin(pi/6)^2 + \cos(pi/6)^2$

1 (Notice the format, and try subtracting 1 from this number.)

**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
function F = Fibonacci(n)% Assign Phi, phi, and F below. Do not change these variable names.
Phi =
phi =
F =end
```
Here is a test of the finished function:

```
>> Fibonacci(30)
ans =832040
\rightarrow
```
## **Solutions to the [Problems for Lecture 5](#page-18-0)**

**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template. Script was saved as trig\_values.m.

```
% Assign x below to a row vector. Do not change this variable name.
x =fprintf('x: '); fprintf('%7.4f ', x); fprintf('\n');
fprintf('cos(x):'); fprintf('%7.4f ', cos(x)); fprintf('\n');
fprintf('sin(x):'); fprintf('%7.4f ', sin(x)); fprintf('\n');
>> trig_values
x: 0.0000 0.5236 0.7854 1.0472 1.5708
cos(x): 1.0000 0.8660 0.7071 0.5000 0.0000
sin(x): 0.0000 0.5000 0.7071 0.8660 1.0000
\gt
```
### **Solutions to the [Problems for Lecture 6](#page-21-0)**

**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
k = 0.05;% Assign theta, x, and y below. Do not change these variable names.
theta =
x =y =% Graphics
plot(x,y)
axis equal
```

```
xlabel('$x$', 'Interpreter', 'latex', 'FontSize', 14)
ylabel('$y$', 'Interpreter', 'latex', 'FontSize', 14)
title('Logarithmic Spiral','Interpreter','latex','FontSize', 16)
```
Your plot should look like this:



#### **2.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
% Assign theta, x, and y below. Do not change these variable names.
theta =
x =y =% graphics
plot(x,y,'r',-x,y,'r')
axis equal;
axis([-1.5 1.5 -1 1])xlabel('$x$', 'Interpreter', 'latex', 'FontSize',14);
ylabel('$y$', 'Interpreter', 'latex', 'FontSize',14);
title('Leminscate', 'Interpreter', 'latex', 'FontSize', 16)
```
#### Your plot should look like this:



**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
% Assign the matrix A. Use the ones.m function and the colon operator.
A =B = A([1,3], [2,4])C = A(:,[1,4:6])D = A([2, 3], :)
```
**2.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
function [A, B, C] = \text{banded_matrix}% Assign A, B and C below.
   A =B =C =end
```
## **Solutions to the [Problems for Lecture 8](#page-27-0)**

```
1.
>> 14>15/3
ans =
    1
>> 8/2<5*3+1>9
ans =
   0
\gg 8/(2<5)*3+(1>9)ans =
   24
>> 2+4+3-60/4-1ans =
    1
\,
```
#### **2.**

```
\Rightarrow u=[4 -2 1]; v=[0 2 1];
>> u \le vans =0 1 1
>> u==v
ans =
```

```
0 0 1
>> u<u+v
ans =0 1 1
\gg (u<v) +u
ans =4 -1 1\gt
```

```
1.
function [p, q] = quadratic_formula_1(a, b, c)
% standard quadratic formula
p=(-b+sqrt(c) \cdot 2 - 4 \cdot *a \cdot *c))./(2*a);
q=(-b-sqrt(b.^{2} - 4.*a.*c))./(2*a);
end
\gg [p, q] = quadratic_formula_1(1, -1, -6)
p =3
q =-2\gg [p, q] = quadratic_formula_1(1, 1, -6)
p =2
q =-3\Rightarrow [p, q] = quadratic_formula_1(1, -1e12, 1)
p =1.0000e+12
q =0
\gtfunction [p, q] = quadratic_formula_2(a, b, c)% better computational quadratic formula
if b<0
    p=(-b+sqrt(c) \cdot 2 - 4 \cdot *a \cdot *c)) \cdot / (2 * a);q=(2*c)./(-b+sqrt(b.^2 - 4.*a.*c));
else
    p=(2*c)./(-b-sqrt(b.^2 - 4.*a.*c));
    q=(-b-sqrt(b.^2 - 4.*a.*c))./(2*a);
end
\gg [p, q] = quadratic_formula_2(1, -1, -6)
p =3
```

```
q =-2\gg [p, q] = quadratic_formula_2(1, 1, -6)
p =2
q =-3\gg [p, q] = quadratic_formula_2(1, -1e12, 1)
p =1.0000e+12
q =1.0000e-12
\rightarrow
```
**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
function F = Fibonacci(n)% Returns the nth Fibonacci number using recursion relation
n=round(n); % n must be an integer
```
Here are some sample function calls:

```
>> Fibonacci(9)
ans =
    34
>> Fibonacci(-9)
ans =
    34
>> Fibonacci(30)
ans =
      832040
>> Fibonacci(-30)
ans =-832040
>>
```
### **Solutions to the [Problems for Lecture 11](#page-33-0)**

**1.** We solve

$$
x = f(f(x)) = f(\mu x(1-x)) = r(\mu x(1-x))(1 - \mu x(1-x)).
$$

Expanding, we obtain the quartic equation

$$
\mu^3 x^4 - 2\mu^3 x^3 + \mu^2 (1 + \mu) x^2 + (1 - \mu^2) x = 0.
$$

Two solutions corresponding to period-1 cycles are known:  $x = 0$  and  $x = 1 - 1/\mu$ . We factor out these two solutions, the second by using long division, and divide by  $\mu$ . The result is the quadratic equation

$$
\mu^2 x^2 - \mu(\mu + 1)x + (\mu + 1) = 0.
$$

The period-2 cycle, then, corresponds to the two roots of this quadratic equation; that is,

$$
x_1 = \frac{1}{2\mu} \left( (\mu + 1) + \sqrt{(\mu + 1)(\mu - 3)} \right), \quad x_2 = \frac{1}{2\mu} \left( (\mu + 1) - \sqrt{(\mu + 1)(\mu - 3)} \right).
$$

These roots are valid solutions for  $\mu \geq 3$ . Exactly at  $\mu = 3$ ,  $x_1 = x_2 = 2/3$ , which coincides with the value of the fixed point  $1 - 1/\mu = 2/3$ . At  $\mu = 3$ , we will see that the fixed point of the logistic map bifurcates to a period-2 cycle.

#### **Solutions to the [Problems for Lecture 12](#page-35-0)**

**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
mu_min=2.4; mu_max=4; %range of mu values
n_mu=500; %number of mu pixels
n_x=400; %number of x pixels
mu_edges=linspace(mu_min,mu_max,n_mu+1); %edges of mu pixels
mu=(mu_edges(1:n_mu)+mu_edges(2:n_mu+1))/2; %values of mu on which to perform computation
x_edges=linspace(0,1,n_x+1); %edges of x pixels
n_trans=200000; %transient iterations
n_data=100000; %number of x values per mu value
x_data=zeros(n_data,n_mu); %x-data used to construct figure
x_0=0.5; %initial condition
```
% WRITE THE COMPUTATIONAL ENGINE OF THE CODE. % USE THE ALREADY DEFINED PARAMETERS AND VARIABLES: n\_mu, mu, x\_0, n\_trans, n\_data. % YOUR FINAL RESULT WILL BE THE VARIABLE x\_data, and this variable will be assessed.

```
%%%%% bin data and plot image %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
x_histogram=zeros(n_x,n_mu); %binned values of x
for i=1:n_mu
x_histogram(:,i)=histcounts(x_data(:,i),x_edges);
x_histogram(:,i)=255*x_histogram(:,i)/max(x_histogram(:,i));
end
colormap(flipud(gray(256))); brighten(-0.8); cmap=colormap;
im=image([mu_edges(1) mu_edges(end)], [x_edges(1) x_edges(end)], x_histogram);
set(gca, 'YDir', 'normal');
xlabel('$\mu$','Interpreter','latex','FontSize',14);
ylabel('$x\;\;$','Interpreter','latex','FontSize',14);
title('Logistic Map Bifurcation Diagram','Interpreter','latex','FontSize',16)
```
**1.** Now  $\sqrt{3}$  is the zero of the function  $f(x) = x^2 - 3$ , and  $f(x_0 = 1) = -2$  and  $f(x_1 = 1)$  $2) = 1$ , so that the two initial guesses bracket the root. We iterate the algorithm. We have

$$
x_2 = 2 - \frac{2 - 1}{2} = \frac{3}{2} = 1.5.
$$

Now,  $f(x_2) = 9/4 - 3 = -3/4 < 0$  so that  $x_1$  and  $x_2$  bracket the root. Therefore,

$$
x_3 = \frac{3}{2} - \frac{\frac{3}{2} - 2}{2} = \frac{7}{4} = 1.75.
$$

Now,  $f(x_3) = 49/16 - 3 = 1/16 > 0$  so that  $x_2$  and  $x_3$  bracket the root. Therefore,

$$
x_4 = \frac{7}{4} - \frac{\frac{7}{4} - \frac{3}{2}}{2} = \frac{13}{8} = 1.624,
$$

and so on.

#### **Solutions to the [Problems for Lecture 14](#page-40-0)**

**1.** We solve  $f(x) = 0$ , where  $f(x) = x^2 - 3$ . We use  $f'(x) = 2x$ . Therefore, Newton's method iterates

$$
x_{n+1} = x_n - \frac{x_n^2 - 3}{2x_n}.
$$

Choosing an initial guess  $x_0 = 1$ , we have

$$
x_1 = 1 - \frac{-2}{2} = 2,
$$
  
\n
$$
x_2 = 2 - \frac{4-3}{4} = \frac{7}{4} = 1.75,
$$
  
\n
$$
x_3 = \frac{7}{4} - \frac{\frac{7^2}{4} - 3}{\frac{14}{4}} = \frac{388}{224} = 1.73214,
$$

already accurate to four significant digits.

### **Solutions to the [Problems for Lecture 15](#page-42-0)**

**1.** We solve  $f(x) = 0$ , where  $f(x) = x^2 - 3$ . The secant method iterates

$$
x_{n+1} = x_n - \frac{(x_n^2 - 3)(x_n - x_{n-1})}{x_n^2 - x_{n-1}^2}
$$
  
=  $x_n - \frac{(x_n^2 - 3)}{x_n + x_{n-1}}$ .

With  $x_0 = 1$  and  $x_1 = 2$ , we have

$$
x_2 = 2 - \frac{4 - 3}{3} = \frac{5}{3} = 1.66666,
$$
  
\n
$$
x_3 = \frac{5}{3} - \frac{\frac{25}{9} - 3}{\frac{5}{3} + 2} = \frac{57}{33} = 1.72727,
$$
  
\n
$$
x_4 = \frac{57}{33} - \frac{(\frac{57}{33})^2 - 3}{\frac{57}{33} + \frac{5}{3}} = \frac{1067}{616} = 1.73214,
$$

and so on.

### **Solutions to the [Problems for Lecture 16](#page-44-0)**

**1.** Use  $|\epsilon_{n+1}| = 0.5 |\epsilon_n|^p$ .



## **Solutions to the [Problems for Lecture 17](#page-46-0)**

**1.** We start with

$$
x_{n+1} = x_n - \frac{(x_n - x_{n-1})f(x_n)}{f(x_n) - f(x_{n-1})}.
$$

*a*) Let  $\epsilon_n = r - x_n$ . Then  $x_n = r - \epsilon_n$  and  $x_{n-1} = r - \epsilon_{n-1}$ , so that

$$
x_n - x_{n-1} = (r - \epsilon_n) - (r - \epsilon_{n-1})
$$
  
= 
$$
(\epsilon_{n-1} - \epsilon_n).
$$
Subtract both sides of the secant method from *r* to obtain

$$
\epsilon_{n+1} = \epsilon_n + \frac{(\epsilon_{n-1} - \epsilon_n) f(r - \epsilon_n)}{f(r - \epsilon_n) - f(r - \epsilon_{n-1})}.
$$

*b*) We Taylor series *f* for small  $\epsilon$ , using  $f(r) = 0$ , to obtain

$$
f(r - \epsilon_n) = -\epsilon_n f'(r) + \frac{1}{2} \epsilon_n^2 f''(r) + \dots,
$$
  

$$
f(r - \epsilon_{n-1}) = -\epsilon_{n-1} f'(r) + \frac{1}{2} \epsilon_{n-1}^2 f''(r) + \dots,
$$

so that

$$
f(r-\epsilon_n)-f(r-\epsilon_{n-1})=(\epsilon_{n-1}-\epsilon_n)f'(r)+\frac{1}{2}(\epsilon_n^2-\epsilon_{n-1}^2)f''(r)+\dots
$$
  

$$
=(\epsilon_{n-1}-\epsilon_n)\left(f'(r)-\frac{1}{2}(\epsilon_{n-1}+\epsilon_n)f''(r)+\dots\right),
$$

and the secant method becomes

$$
\epsilon_{n+1} = \epsilon_n + \frac{-\epsilon_n f'(r) + \frac{1}{2} \epsilon_n^2 f''(r) + \dots}{f'(r) - \frac{1}{2} (\epsilon_{n-1} + \epsilon_n) f''(r) + \dots}
$$

*c*) We rewrite the previous equation to obtain

$$
\epsilon_{n+1} = \epsilon_n - \epsilon_n \frac{1 - \frac{1}{2} \epsilon_n \frac{f''(r)}{f'(r)} + \dots}{1 - \frac{1}{2} (\epsilon_{n-1} + \epsilon_n) \frac{f''(r)}{f'(r)} + \dots}
$$

We now Taylor series expand the denominator to obtain

$$
\epsilon_{n+1} = \epsilon_n - \epsilon_n \left( 1 - \frac{1}{2} \epsilon_n \frac{f''(r)}{f'(r)} + \dots \right) \left( 1 + \frac{1}{2} \left( \epsilon_{n-1} + \epsilon_n \right) \frac{f''(r)}{f'(r)} + \dots \right)
$$
  
= 
$$
- \frac{1}{2} \frac{f''(r)}{f'(r)} \epsilon_n \epsilon_{n-1} + \dots,
$$

or to leading order

<span id="page-180-0"></span>
$$
|\epsilon_{n+1}| = \frac{1}{2} \left| \frac{f''(r)}{f'(r)} \right| |\epsilon_n||\epsilon_{n-1}|.
$$
 (1)

*d*) To determine the order of convergence, we look for a solution of the form

$$
|\epsilon_{n+1}| = k |\epsilon_n|^p, \qquad |\epsilon_n| = k |\epsilon_{n-1}|^p.
$$

Substitution into [\(1\)](#page-180-0) results in

$$
k^{p+1}|\epsilon_{n-1}|^{p^2} = \frac{k}{2} \left| \frac{f''(r)}{f'(r)} \right| |\epsilon_{n-1}|^{p+1}.
$$

Equating the power of  $\epsilon_{n-1}$  results in

$$
p^2 = p + 1.
$$

The rate of convergence of the secant method, given by  $p$ , is therefore the positive root of the quadratic equation  $p^2 - p - 1 = 0$ , or

$$
p = \frac{\sqrt{5} + 1}{2} \approx 1.618,
$$

which coincidentally is the famous irrational number that is called the golden ratio, and goes by the symbol Φ.

#### **Solutions to the [Problems for Lecture 18](#page-48-0)**

**1.** We are looking for the four solutions to  $z^4 = 1$ . Clearly, one and minus one are solutions. Since  $i^2 = (-i)^2 = -1$ , the four solutions, moving counterclockwise around the unit circle, are evidently

$$
r_1 = 1
$$
,  $r_2 = i$ ,  $r_3 = -1$ ,  $r_4 = -i$ .

Their location on the unit circle in the complex plane is shown below.



**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
% code the function and its derivative that will be used in Newton's method
f = (g(z)) add function here
fp = @ (z) add derivative here
% assign the four fourth roots of unity,
% starting at 1 and moving around the unit circle counterclockwise
root1 = 1;root2 = add second root here
root3 = add third root here
root4 = add fourth root here
nx=1000; ny=1000;
xmin=-2; xmax=2; ymin=-2; ymax=2;
x=linspace(xmin, xmax, nx); y=linspace(ymin, ymax, ny);
[X, Y]=meshqrid(x, y);
Z = X + 1 i * Y;nit=50;
for n=1:nit
    Z = Z - f(Z)./ fp(Z);
end
eps=0.001;
Z1 = abs(Z-root1) < eps; Z2 = abs(Z-root2) < eps;Z3 = abs(Z-root3) < eps; Z4 = abs(Z-root4) < eps;Z5 = -(Z1+Z2+Z3+Z4);figure;
map = [1 0 0; 0 1 0; 0 0 1; 1 1 0; 0 0 0];colormap(map); %[red;blue;green;yellow;black]
Z = (Z1 + 2 \times Z2 + 3 \times Z3 + 4 \times Z4 + 5 \times Z5);
image([xmin xmax], [ymin ymax], Z); set(gca,'YDir','normal');
xlabel('$x$', 'Interpreter', 'latex', 'FontSize',14);
ylabel('$y$', 'Interpreter', 'latex', 'FontSize',14);
title('Fractal from $f(z)=z^4-1$', 'Interpreter', 'latex','FontSize', 16)
```
#### Your plot should look like this:



**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
period=1; omega=2*pi/period;
e=[0,1/4,1/2,3/4]; color=['r','g','b','c'];
a=1./(1-e); b=sqrt((1+e) ./(1-e));t=linspace(0,period,1000);
x=zeros(length(t),length(e)); y=zeros(length(t),length(e));
for j=1:length(e)
    for i=1:length(t)
        E=fzero((\mathcal{C}(E) ..., 0); % add anonymous function for root finding.
                               % Make use of the variables e(j), t(i) and omega.
        x(i, j) = ; % assign x-coordinate. Make use of the variables a(j), e(j) and E.
        y(i, j) = j % assign y-coordinate. Make use of the variables b(j) and E.
    end
end
for j=1:length(e)
    plot(x(:,j),y(:,j),color(j)); axis equal; hold on;
end
plot(0,0,'xk') %mark the origin
xlabel('$x$', 'Interpreter', 'latex', 'FontSize',14)
ylabel('$y$', 'Interpreter', 'latex', 'FontSize',14)
legend('$e=0$','$e=1/4$','$e=1/2$','$e=3/4$','Interpreter','latex','Location','East')
title('Planetary Orbits', 'Interpreter', 'latex', 'FontSize',16)
```
Your plot should look like this:



#### **1.**

*a*) The period-two fixed point equations are given by

$$
x_1 = \mu x_0 (1 - x_0), \qquad x_0 = \mu x_1 (1 - x_1).
$$

We can eliminate  $x_1$  to obtain the single equation

$$
x_0 = \mu(\mu x_0(1 - x_0))(1 - \mu x_0(1 - x_0)).
$$

Substituting  $x_0 = 1/2$ , we have

$$
\frac{1}{2} = \mu(\frac{\mu}{4})(1 - \frac{\mu}{4}).
$$

Multiplying by 16 and expanding, we obtain  $\mu^3 - 4\mu^2 + 8 = 0$ .

b) The long division, complements of the LAT<sub>E</sub>X Polynom Package, is given by

$$
\begin{array}{r} x^2 - 2x - 4 \\
x - 2 \overline{\smash{\big)}\ x^3 - 4x^2} + 8 \\
\underline{-x^3 + 2x^2} \\
-2x^2 \\
\underline{2x^2 - 4x} \\
-4x + 8 \\
\underline{4x - 8} \\
0\n\end{array}
$$

The roots of

$$
\mu^2-2\mu-4=0
$$

are given by

$$
\mu = \frac{2 \pm \sqrt{4 + 16}}{2} = 1 \pm \sqrt{5},
$$

and the positive root yields  $m_1 = 1 +$ √ 5.

#### **Solutions to the [Problems for Lecture 23](#page-58-0)**

**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
% Compute the Feigenbaum delta
% Store approximate values in the row vector delta for assessment,
% where length(delta)= num_doublings and delta(2:num_doublings)
% are computed from the algorithm described in Lectures 21-23.
```

```
num_doublings=11; delta=zeros(1,num_doublings); delta(1)=5;
% Write your code here
```

```
% Output your results
fprintf('n delta(n)\n');
for n=1:num_doublings
   fprintf('%2g %18.15f\n',n,delta(n));
end
```
1. Using MATLAB as a calculator, we have

```
\Rightarrow x2=(-4/eps + 1)/(-2/eps - 1)
x^2 =2
>> x1 = (4-2*x2)/epsx1 =\Omega\Rightarrow x2=(-4/(2*eps) + 1)/(-2/(2*eps) - 1)
x2 =2.0000
>> x1=(4-2*x2)/(2*eps)
x1 =3
>>
```
# **Solutions to the [Problems for Lecture 25](#page-63-0)**

1. Using MATLAB as a calculator, we have

```
>> x2=(4-eps)/(2+eps)
x2 =2
>> x1=1+x2
x1 =
```

```
3
>> x2 = (4-2*eps) / (2+2*eps)x2 =2.0000
>> x1=1+x2x1 =3.0000
\gt
```
**1.** We begin with a row exchange:

$$
A = \begin{pmatrix} -3 & 2 & -1 \\ 6 & -6 & 7 \\ 3 & -4 & 4 \end{pmatrix} \rightarrow \begin{pmatrix} 6 & -6 & 7 \\ -3 & 2 & -1 \\ 3 & -4 & 4 \end{pmatrix} = P_{12}A.
$$

The first elimination step is

$$
P_{12}A \rightarrow \begin{pmatrix} 6 & -6 & 7 \\ 0 & -1 & 5/2 \\ 3 & -4 & 4 \end{pmatrix} = M_1P_{12}A, \quad \text{where } M_1 = \begin{pmatrix} 1 & 0 & 0 \\ 1/2 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}.
$$

The second elimination step is

$$
M_1P_{12}A \rightarrow \begin{pmatrix} 6 & -6 & 7 \\ 0 & -1 & 5/2 \\ 0 & -1 & 1/2 \end{pmatrix} = M_2M_1P_{12}A, \quad \text{where } M_2 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -1/2 & 0 & 1 \end{pmatrix}.
$$

No more row interchanges are needed and the third elimination step is

$$
M_2M_1P_{12}A \rightarrow \begin{pmatrix} 6 & -6 & 7 \\ 0 & -1 & 5/2 \\ 0 & 0 & -2 \end{pmatrix} = M_3M_2M_1P_{12}A = U, \text{ where } M_3 = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -1 & 1 \end{pmatrix}.
$$

Taking the inverse of the matrix multiplying A, we have

$$
A = P_{12}(M_3M_2M_1)^{-1}U = (P_{12}L)U,
$$

where the upper triangular and psychologically lower triangular matrices are given by

$$
U = \begin{pmatrix} 6 & -6 & 7 \\ 0 & -1 & 5/2 \\ 0 & 0 & -2 \end{pmatrix}, \qquad (P_{12}L) = P_{12} \begin{pmatrix} 1 & 0 & 0 \\ -1/2 & 1 & 0 \\ 1/2 & 1 & 1 \end{pmatrix} = \begin{pmatrix} -1/2 & 1 & 0 \\ 1 & 0 & 0 \\ 1/2 & 1 & 1 \end{pmatrix}.
$$

**1.** Let  $T_L$  be the time to compute recombination with *L* loci. We have

$$
T_L = k3^L,
$$

so that

$$
T_{L+1} = k3^{L+1} = 3k3^L = 3T_L.
$$

If  $L = 15$  takes 10 sec, then the time for  $L = 16$  is estimated to be  $3 \times 10$  sec = 30 sec.

#### **Solutions to the [Problems for Lecture 28](#page-70-0)**

- **1.**
- *a*) The summation denoted as  $\sum_{k=1}^{n} 1$  is just the addition of *n* ones and is equal to *n*.
- *b*) The simplest way to compute  $\sum_{k=1}^{n} k$  is to define this sum to be *X*, and to write *X* twice as follows:

$$
X = 1+ 2 + 3 + ... + n
$$
  

$$
X = n + (n - 1) + (n - 2) + ... + 1.
$$

Adding these two equations yields  $2X = n(n+1)$ , or  $X = n(n+1)/2$ .

*c*) The trick to compute  $\sum_{k=1}^{n} k^2$  is to use the expansion  $(k-1)^3 = k^3 - 3k^2 + 3k - 1$ , and to write  $k^3 - (k-1)^3 = 3k^2 - 3k + 1$ . Since  $\sum_{k=1}^n k^3 - (k-1)^3$  is what is called a telescoping series (each new term cancels with part of the preceding term), we have

$$
\sum_{k=1}^{n} k^3 - (k-1)^3 = n^3 = 3 \sum_{k=1}^{n} k^2 - 3 \sum_{k=1}^{n} k + \sum_{k=1}^{n} 1.
$$

Therefore,

$$
\sum_{k=1}^{n} k^2 = \frac{1}{3} \left( n^3 + \frac{3n(n+1)}{2} - n \right) = \frac{1}{6} n (2n+1)(n+1).
$$

#### **Solutions to the [Problems for Lecture 29](#page-72-0)**

**1.** The solution for  $x_i$  is found after solving for  $x_j$  with  $j < i$ . The explicit solution for  $x_i$  is given by

$$
x_i = \frac{1}{a_{ii}} \left( b_i - \sum_{j=1}^{i-1} a_{ij} x_j \right).
$$

The solution for  $x_i$  requires  $i$  multiplication-additions, and since this must be done for  $n$ such *i's*, we have

op. counts 
$$
=\sum_{i=1}^{n} i = \frac{1}{2}n(n+1).
$$

**1.** The two largest eigenvalues are  $\lambda_1 = 1$  and  $\lambda_2 = 1/2$  with corresponding eigenvectors *e*<sub>1</sub> and *e*<sub>2</sub>. Suppose  $x_0 = c_1e_1 + c_2e_2 + ...$  Then

$$
x_p = c_1 \lambda_1^p e_1 + c_2 \lambda_2^p e_2 + \dots = c_1 e_1 + (1/2)^p c_2 e_2 + \dots
$$

We estimate that to converge to an error of less than  $10^{-8}$ , we need to require

$$
(1/2)^p \approx 10^{-8},
$$

or  $p \approx 8/\log_{10} 2 \approx 27$ .

#### **Solutions to the [Problems for Lecture 31](#page-76-0)**

**1.** We apply the power method to

$$
A = \begin{pmatrix} -5 & 6 \\ 5 & -4 \end{pmatrix}.
$$

The assumed initial vector and first iteration is given by

$$
x_0 = \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \qquad x_1 = Ax_0 = \begin{pmatrix} -5 \\ 5 \end{pmatrix}.
$$

Continuing,

$$
x_2 = \begin{pmatrix} -5 & 6 \\ 5 & -4 \end{pmatrix} \begin{pmatrix} -5 \\ 5 \end{pmatrix} = \begin{pmatrix} 55 \\ -45 \end{pmatrix}; \qquad x_3 = \begin{pmatrix} -5 & 6 \\ 5 & -4 \end{pmatrix} \begin{pmatrix} 56 \\ 44 \end{pmatrix} = \begin{pmatrix} -545 \\ 455 \end{pmatrix}.
$$

Two more iterations give

$$
x_4 = \begin{pmatrix} -5 & 6 \\ 5 & -4 \end{pmatrix} \begin{pmatrix} -545 \\ 455 \end{pmatrix} = \begin{pmatrix} 5,455 \\ -4,545 \end{pmatrix}; \qquad x_5 = \begin{pmatrix} -5 & 6 \\ 5 & -4 \end{pmatrix} \begin{pmatrix} 5,455 \\ -4,545 \end{pmatrix} = \begin{pmatrix} -54,545 \\ 45,455 \end{pmatrix}.
$$

The dominant eigenvalue is approximated from

$$
\lambda_1 \approx \frac{x_4^{\text{T}} x_5}{x_4^{\text{T}} x_4} = \frac{-504,135,950}{50,414,050} \approx -9.99991 \approx -10;
$$

and the corresponding eigenvector is approximated by  $x_5$ . Dividing by the second component,

$$
v_1 = \begin{pmatrix} -54.545/45.455 \\ 1 \end{pmatrix} \approx \begin{pmatrix} -1.19998 \\ 1 \end{pmatrix} \approx \begin{pmatrix} -6/5 \\ 1 \end{pmatrix}.
$$

**1.** Complete your solution on Matlab. Here is the Learner Template:

```
% Define matrix A
% Define L to be the psychologically lower triangular matrix
% Define U to be the upper triangular matrix
```
**2.** Complete your solution on Matlab. Here is the Learner Template:

```
% Define matrix A
% Find eigenvalues, lambda1<lambda2
lambda1= ...
lambda2= ...
% Find eigenvectors associated with lambda1 and lambda2
% Normalize eigenvectors so that their second component is one
v1 = 1...
v2 = ...
```
## **Solutions to the [Problems for Lecture 33](#page-81-0)**

**1.** We write down the algorithm to solve

$$
f(x, y, z) = 0, \t g(x, y, z) = 0, \t h(x, y, z) = 0.
$$

1. Solve the linear system for  $\Delta x_n$ ,  $\Delta y_n$  and  $\Delta z_n$  given by

$$
\begin{pmatrix} f_x & f_y & f_z \ g_x & g_y & g_z \ h_x & h_y & h_z \end{pmatrix} \begin{pmatrix} \Delta x_n \\ \Delta y_n \\ \Delta z_n \end{pmatrix} = \begin{pmatrix} -f \\ -g \\ -h \end{pmatrix}.
$$

2. Advance the iterative solution, using

$$
x_{n+1} = x_n + \Delta x_n
$$
,  $y_{n+1} = y_n + \Delta y_n$ ,  $z_{n+1} = z_n + \Delta z_n$ .

## **Solutions to the [Problems for Lecture 34](#page-83-0)**

**1.** We solve

$$
\sigma(y - x) = 0
$$
,  $rx - y - xz = 0$ ,  $xy - \beta z = 0$ .

By inspection, the first fixed-point solution is  $(x, y, z) = (0, 0, 0)$ , and another two can be found. The first equation yields  $x = y$ . Eliminating *y* from the second equation results in  $x(r-1-z) = 0$ , with solutions  $x = 0$  or  $z = r-1$ . The  $x = 0$  solution corresponds to the

first fixed-point solution. Substituting  $y = x$  and  $z = r - 1$  into the third equation results in  $x^2 = \beta(r-1)$ . The square root results in the second and third fixed-point solutions. Our three solutions are

$$
(x,y,z)=(0,0,0), \quad (\sqrt{\beta(r-1)},\sqrt{\beta(r-1)},r-1), \quad (-\sqrt{\beta(r-1)},-\sqrt{\beta(r-1)},r-1).
$$

When  $r = 28$ ,  $\sigma = 10$  and  $\beta = 8/3$ , the numerical values are

$$
(x,y,z) = (0,0,0), \quad (8.48528, 8.48528, 27), \quad (-8.48528, -8.48528, 27).
$$

#### **2.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
r=28; sigma=10; b=8/3;
RelTol=1.e-06; AbsTol=1.e-09;
for nroot=1:3
    if nroot==1, x=1; y=1; z=1; end
    if nroot==2, x=10; y=10; z=10; end
    if nroot==3, x=-10; y=-10; z=-10; end
    error=Inf;
    while error > max(RelTol*max(abs([x,y,z])),AbsTol)
        J= % DEFINE THE JACOBIAN MATRIX
        rhs = % DEFINE THE RIGHT-HAND SIDE
        delta_xyz=J\rhs;
        x = x + delta_x yz(1);y = y + delta_{xyz}(2);z = z + delta xyz(3);error=max(abs(delta_xyz));
    end
    xroot(nroot)=x; yroot(nroot)=y; zroot(nroot)=z;
end
roots=[xroot;yroot;zroot];
fprintf('steady-state solution:\n')
fprintf('(x, y, z) = (\$2.0f, \$2.0f, \$2.0f) \ \n\infty, roots(:,1));
fprintf('(x, y, z) = (\frac{7.5f, \frac{2}{167.5f, \frac{2}{167.0f}}{n'}, \frac{3.0f}{n'}, \frac{3.0f}{n'});
fprintf('(x, y, z) = (87.5f, 87.5f, 83.0f) \n', roots(:,3));
```
#### **Solutions to the [Problems for Lecture 35](#page-85-0)**

**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
r=28; sigma=10; beta=8/3;
x1=0; y1=0; z1=0;
x2 = sqrt(beta * (r-1)); y2 = sqrt(beta * (r-1)); z2 = r-1;x3 = -sqrt(heta*(r-1)); y3 = -sqrt(heta*(r-1)); z3 = r-1;nx=500; ny=500;
xmin=-30; xmax=30; ymin=-30; ymax=30;
```

```
x_grid=linspace(xmin,xmax,nx); y_grid=linspace(ymin,ymax,ny);
[X,Y]=meshgrid(x_grid,y_grid);
% Write Newton's method using every gridpoint as the initial condition
% Perform enough iterations at each gridpoint to converge to the correct root
% Save the x-values of the converged roots in the matrix X
% To pass the assessment, every pixel in the figure must be correctly colored
%!!!!!!!!! Set initial value z=-10 for all values (x,y) on the grid !!!!!!!!!!
eps=0.001;
X1 = abs(X-x1) < eps; X2 = abs(X-x2) < eps; X3 = abs(X-x3) < eps;X4 = - (X1+X2+X3);
figure;
map = [1\ 0\ 0; \ 0\ 1\ 0; \ 0\ 0\ 1; \ 0\ 0\ 0]; colormap(map); \S[red;green;blue;black]X = (X1 + 2 \times X2 + 3 \times X3 + 4 \times X4);
image([xmin xmax], [ymin ymax], X); set(gca,'YDir','normal');
xlabel('$x$', 'Interpreter', 'latex', 'FontSize',14);
ylabel('$y$', 'Interpreter', 'latex', 'FontSize',14);
title('Fractal from the Lorenz Equations', 'Interpreter', 'latex','FontSize', 16)
```
**1.** The midpoint rule approximates the area under the curve of  $y = f(x)$  from zero to *h* by the area of a rectangle with base *h* and height *f*(*h*/2), as illustrated in the graph below.



## **2.** We have

$$
\int_0^h f(x) dx = \int_0^h (a + bx + cx^2) dx
$$
  
=  $(ax + \frac{bx^2}{2} + \frac{cx^3}{3}) \Big|_0^h$   
=  $ah + \frac{bh^2}{2} + \frac{ch^3}{3}$   
=  $h \left( a + \frac{bh}{2} + \frac{ch^2}{3} \right)$ .

Now,

$$
f(h/2) = a + \frac{bh}{2} + \frac{ch^2}{4}, \qquad f''(h/2) = 2c.
$$

Therefore,

$$
\int_0^h f(x) dx = h \left( a + \frac{bh}{2} + \frac{ch^2}{3} \right)
$$
  
=  $h \left( a + \frac{bh}{2} + \frac{ch^2}{4} \right) + h \left( \frac{ch^2}{3} - \frac{ch^2}{4} \right)$   
=  $h \left( a + \frac{bh}{2} + \frac{ch^2}{4} \right) + h \left( \frac{ch^2}{12} \right)$   
=  $hf(h/2) + \frac{h^3}{24} f''(h/2).$ 

# **Solutions to the [Problems for Lecture 37](#page-90-0)**

**1.** We approximate

$$
f(x) \approx f(0) + \frac{f(h) - f(0)}{h}x.
$$

Then

$$
\int_0^h f(x) dx = \int_0^h \left( f(0) + \frac{f(h) - f(0)}{h} x \right) dx
$$
  
=  $\left( f(0)x + \frac{f(h) - f(0)}{2h} x^2 \right) \Big|_0^h$   
=  $f(0)h + \frac{1}{2} (f(h) - f(0)) h$   
=  $\frac{h}{2} (f(0) + f(h)).$ 

**1.**

*a*) Let  $g(x) = a + bx + cx^2$ . The equations that we want to solve for *a*, *b* and *c* are given by

$$
g(0) = a = f(0),
$$
  
\n
$$
g(h) = a + hb + h2c = f(h),
$$
  
\n
$$
g(2h) = a + 2hb + 4h2c = f(2h),
$$

or in matrix form

$$
\begin{pmatrix} 1 & 0 & 0 \ 1 & h & h^2 \ 1 & 2h & 4h^2 \end{pmatrix} \begin{pmatrix} a \\ b \\ c \end{pmatrix} = \begin{pmatrix} f(0) \\ f(h) \\ f(2h) \end{pmatrix}.
$$

We form the augmented matrix and solve by Gaussian elimination:

$$
\begin{pmatrix}\n1 & 0 & 0 & f(0) \\
1 & h & h^2 & f(h) \\
1 & 2h & 4h^2 & f(2h)\n\end{pmatrix}\n\rightarrow\n\begin{pmatrix}\n1 & 0 & 0 & f(0) \\
0 & h & h^2 & f(h) - f(0) \\
0 & 2h & 4h^2 & f(2h) - f(0)\n\end{pmatrix}\n\rightarrow\n\begin{pmatrix}\n1 & 0 & 0 & f(0) \\
0 & h & h^2 & f(h) - f(0) \\
0 & 0 & 2h^2 & f(2h) - 2f(h) + f(0)\n\end{pmatrix}.
$$

Backsubstitution gives

$$
c = \frac{1}{2h^2} (f(0) - 2f(h) + f(2h)),
$$
  
\n
$$
b = \frac{1}{h} \left( f(h) - f(0) - \frac{1}{2} (f(2h) - 2f(h) + f(0)) \right) = \frac{1}{2h} (-3f(0) + 4f(h) - f(2h)),
$$
  
\n
$$
a = f(0).
$$

*b*) We approximate  $f(x)$  by the quadratic polynomial  $g(x)$ , that is,

$$
f(x) \approx f(0) + \frac{1}{2h} \left( -3f(0) + 4f(h) - f(2h) \right) x + \frac{1}{2h^2} \left( f(0) - 2f(h) + f(2h) \right) x^2.
$$

Then

$$
\int_0^{2h} f(x) dx \approx \int_0^{2h} \left( f(0) + \frac{1}{2h} \left( -3f(0) + 4f(h) - f(2h) \right) x \right. \\
\left. + \frac{1}{2h^2} \left( f(0) - 2f(h) + f(2h) \right) x^2 \right) dx
$$
\n
$$
= \left( f(0)x + \frac{1}{4h} \left( -3f(0) + 4f(h) - f(2h) \right) x^2 + \frac{1}{6h^2} \left( f(0) - 2f(h) + f(2h) \right) x^3 \right) \Big|_0^{2h}
$$
\n
$$
= 2hf(0) + h\left( -3f(0) + 4f(h) - f(2h) \right) + \frac{4h}{3} \left( f(0) - 2f(h) + f(2h) \right)
$$
\n
$$
= \frac{h}{3} \left( f(0) + 4f(h) + f(2h) \right).
$$

## **Solutions to the [Problems for Lecture 39](#page-94-0)**

**1.** We have

$$
\int_{a}^{b} f(x) dx = \frac{3h}{8} (f_0 + 3f_1 + 3f_2 + f_3) + \frac{3h}{8} (f_3 + 3f_4 + 3f_5 + f_6) + \dots
$$
  
+ 
$$
\frac{3h}{8} (f_{n-3} + 3f_{n-2} + 3f_{n-1} + f_n)
$$
  
= 
$$
\frac{3h}{8} (f_0 + 3f_1 + 3f_2 + 2f_3 + 3f_4 + 3f_5 + 2f_6 + \dots + 3f_{n-2} + 3f_{n-1} + f_n),
$$

where the first and last terms have a multiple of one, the terms with indices that have a remainder of one or two when divided by three have a multiple of three, the terms with indices that are a multiple of three have an index of two, and the entire sum is multiplied by 3*h*/8.

#### **Solutions to the [Problems for Lecture 40](#page-96-0)**

**1.** The three-point Legendre-Gauss quadrature rule is written as

$$
\int_{-1}^{1} f(x) dx = w_1 f(x_1) + w_2 f(x_2) + w_3 f(x_3).
$$

We will assume that  $x_1 = -x_3$ ,  $w_1 = w_3$  and  $x_2 = 0$ . Substituting for the function the basis polynomials  $f(x) = 1$ , x,  $x^2$ ,  $x^3$ ,  $x^4$  and  $x^5$ , we obtain the six equations given by

$$
2 = w_1 + w_2 + w_3, \quad 0 = w_1 x_1 + w_2 x_2 + w_3 x_3, \quad 2/3 = w_1 x_1^2 + w_2 x_2^2 + w_3 x_3^2,
$$
  
\n
$$
0 = w_1 x_1^3 + w_2 x_2^3 + w_3 x_3^3, \quad 2/5 = w_1 x_1^4 + w_2 x_2^4 + w_3 x_3^4, \quad 0 = w_1 x_1^5 + w_2 x_2^5 + w_3 x_3^5.
$$

Using the symmetry conditions, the third and fifth equations result in

$$
2w_1x_1^2 = 2/3, \qquad 2w_1x_1^4 = 2/5.
$$

Division of the second of these two equations by the first yields

$$
x_1^2 = 3/5
$$
, or  $x_1 = -\sqrt{3/5}$ ,

and the third of the original equations then yields  $w_1 = 5/9$ . The first equation yields  $w_2 = 8/9$ . To summarize, we have found

$$
w_1 = 5/9
$$
,  $w_2 = 8/9$ ,  $w_3 = 5/9$ ;  $x_1 = -\sqrt{3/5}$ ,  $x_2 = 0$ ,  $x_3 = \sqrt{3/5}$ .

#### **Solutions to the [Problems for Lecture 41](#page-98-0)**

**1.** For  $\int_0^h f(x) dx$  with  $f(x) = x^3$ , we have

$$
S_1 = \frac{h}{2} \left( f(0) + f(h) \right) = \frac{h^4}{2}, \quad S_2 = \frac{h}{4} \left( f(0) + 2f(h/2) + f(h) \right) = \frac{5h^4}{16}.
$$

Since  $I = h^4/4$ , we have

$$
E_1 = I - S_1 = -\frac{h^4}{4}, \quad E_2 = I - S_2 = -\frac{h^4}{16},
$$

and  $E_1 = 4E_2$ .

## **Solutions to the [Problems for Lecture 42](#page-100-0)**

**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
c = \theta(x) ... % assign the integrand for C(t)s = \theta(x) ... % assign the integrand for S(t)tmin=-8; tmax=8; nt=2000;
t=linspace(tmin,tmax,nt);
C=zeros(nt,1); S=zeros(nt,1);
for i=1:nt
    C(i)=integral(...); % compute C(i) using integral.m and the integrand c(x) defined on
    S(i)=intereral(...); % compute S(i) using integral.m and the integrand S(x) defined on
end
plot(S,C)
xlabel('$S(t)$', 'Interpreter', 'latex', 'FontSize',14);
ylabel('$C(t)$', 'Interpreter', 'latex', 'FontSize',14);
title('Cornu spiral', 'Interpreter', 'latex', 'FontSize',16);
```
Your plot should look like this:



#### **1.**

*a*) We interpolate the points  $(0, 0)$ ,  $(1, 1)$  and  $(2, 1)$ . Let  $y = ax^2 + bx + c$ . Our equations for *a*, *b* and *c* are

$$
c = 0
$$
,  $a+b+c=1$ ,  $4a+2b+c=1$ .

Solving, we find  $a = -1/2$ ,  $b = 3/2$  and  $c = 0$ . Our interpolating quadratic polynomial is

$$
g(x) = -\frac{1}{2}x^2 + \frac{3}{2}x,
$$

and we find  $g(1/2) = 5/8$  and  $g(3/2) = 9/8$ .

*b*) The piecewise linear polynomials are

$$
g_0(x) = x
$$
,  $g_1(x) = 1$ ,

and we find  $g(1/2) = 1/2$  and  $g(3/2) = 1$ .

*c*) The points and the two interpolating functions are plotted below.



#### **1.** We use

$$
g_i(x) = a_i(x - x_i)^3 + b_i(x - x_i)^2 + c_i(x - x_i) + d_i
$$
, for  $i = 0, ..., n - 1$  and  $x_i \le x \le x_{i+1}$ ,

and the additional constraints given by  $g'_0(x_0) = y'_0$  and  $g'_{n-1}(x_n) = y'_n$ . The two constraints yield the two additional equations given by

$$
y'_{0} = c_{0},
$$
  

$$
y'_{n} = 3a_{n-1}h_{n-1}^{2} + 2b_{n-1}h_{n-1} + c_{n-1}.
$$

## **Solutions to the [Problems for Lecture 45](#page-107-0)**

**1.** The points to interpolate are  $(0, 0)$ ,  $(1, 1)$ ,  $(2, 1)$  and  $(3, 2)$ . With  $h_i = x_{i+1} - x_i$  and  $\eta_i = y_{i+1} - y_i$ , we have

$$
h_i = 1,
$$
  $\eta_0 = 1,$   $\eta_1 = 0,$   $\eta_2 = 1.$ 

The not-a-knot condition yields the two equations

$$
b_0 - 2b_1 + b_2 = 0, \qquad b_1 - 2b_2 + b_3 = 0.
$$

The full matrix equation for the *b*-coefficients is then calculated to be

$$
\begin{pmatrix} 1 & -2 & 1 & 0 \ 1/3 & 4/3 & 1/3 & 0 \ 0 & 1/3 & 4/3 & 1/3 \ 0 & 1 & -2 & 1 \ \end{pmatrix} \begin{pmatrix} b_0 \ b_1 \ b_2 \ b_3 \end{pmatrix} = \begin{pmatrix} 0 \ -1 \ 1 \ 0 \end{pmatrix}.
$$

We can use Matlab as a calculator to solve for *b*, and then find *a*, *c* and *d*:

$$
b = \begin{pmatrix} -3/2 \\ -1/2 \\ 1/2 \\ 3/2 \end{pmatrix}, \quad a = \frac{1}{3} \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}, \quad c = \begin{pmatrix} 13/6 \\ 1/6 \\ 1/6 \end{pmatrix}, \quad d = \begin{pmatrix} 0 \\ 1 \\ 1 \end{pmatrix}.
$$

The three interpolating cubic polynomials are given by

$$
g_0(x) = \frac{1}{3}x^3 - \frac{3}{2}x^2 + \frac{13}{6}x,
$$
  
\n
$$
g_1(x) = \frac{1}{3}(x-1)^3 - \frac{1}{2}(x-1)^2 + \frac{1}{6}(x-1) + 1,
$$
  
\n
$$
g_2(x) = \frac{1}{3}(x-2)^3 + \frac{1}{2}(x-2)^2 + \frac{1}{6}(x-2) + 1.
$$

If fact, because of the not-a-knot condition,  $g_0 = g_1 = g_2$ . The points and the cubic spline interpolant are plotted below.



# **Solutions to the [Problems for Lecture 46](#page-109-0)**

**1.** Complete your solution on Coursera using MATLAB. Here is the Learner Template:

```
load data1; figure(1); plot(x1,y1);
load data2; figure(2); plot(x2,y2);
xx=linspace(0,2*pi,1000);
yy1=interp1(...); %interpolate y1 to the xx grid using 'spline'
yy2=interp1(...); %interpolate y2 to the xx grid using 'spline'
yyadd=yy1+yy2;
figure(3); plot(xx, yyadd);
```
#### Here are the resulting plots.





**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
num_roots=5; num_functions=6;
%initial guess for roots (from Wolfram MathWorld)
zeros_guess=[2.4,3.8,5.1,6,7.5,8.7;...
  5.5,7,8.4,9.7,11,12;...
  8.6 10,11.6,13,14,16;...
  11.8,13,15,16,18,19;...
   15,16.4,18,19.4,21,22];
%Compute first five roots of first six Bessel functions
Put in variable bzeros with size(bzeros) = [5, 6]%print table
fprintf('k J0(x) J1(x) J2(x) J3(x) J4(x) J5(x)\n')
for k=1:num_roots
   fprintf('%i',k)
   for n=0:num_functions-1
       fprintf('%10.4f',bzeros(k,n+1));
   end
   fprintf('n');
end
```
#### **Solutions to the [Problems for Lecture 48](#page-114-0)**

**1.** Let  $\dot{x} = b$  and  $x(0) = x_0$ . The Euler method applied to the first time step yields

$$
x_1 = x_0 + b\Delta t,
$$

where  $x_1 = x(\Delta t)$ . The next time step yields

$$
x_2 = x_1 + b\Delta t = x_0 + b(2\Delta t),
$$

where  $x_2 = x(2\Delta t)$ . Continuing to march the solution forward, we obtain

$$
x_n = x_0 + b(n\Delta t),
$$

which is equivalent to

$$
x(t)=x_0+bt,
$$

where  $t = n\Delta t$ .

## **Solutions to the [Problems for Lecture 49](#page-116-0)**

**1.** Let  $\dot{x} = bt$  and  $x(0) = x_0$ . The Modified Euler method applied to the first time step yields

$$
k_1 = 0
$$
,  $k_2 = b(\Delta t)^2$ ,  $x_1 = x_0 + \frac{1}{2}b(\Delta t)^2$ ,

where  $x_1 = x(\Delta t)$ . The next time step yields

$$
k_1 = b(\Delta t)^2
$$
,  $k_2 = 2b(\Delta t)^2$ ,  $x_2 = x_1 + \frac{3}{2}b(\Delta t)^2 = x_0 + 2b(\Delta t)^2 = x_0 + \frac{1}{2}b(2\Delta t)^2$ .

Continuing to march the solution forward, we obtain

$$
x_n = x_0 + \frac{1}{2}b(n\Delta t)^2,
$$

which is equivalent to

$$
x(t) = x_0 + \frac{1}{2}bt^2,
$$

where  $t = n\Delta t$ .

## **Solutions to the [Problems for Lecture 51](#page-119-0)**

**1.** Ralston's method is given by

$$
k_1 = \Delta t f(t_n, x_n), \qquad k_2 = \Delta t f\left(t_n + \frac{3}{4}\Delta t, x_n + \frac{3}{4}k_1\right),
$$

$$
x_{n+1} = x_n + \frac{1}{3}k_1 + \frac{2}{3}k_2.
$$

**2.** Consider the ode given by  $dy/dx = f(x)$ , with  $y(0)$  as the initial value. Separating variables, and integrating from  $x = 0$  to  $h$  yields

$$
\int_{y(0)}^{y(h)} dy = \int_0^h f(x) dx, \quad \text{or} \quad \int_0^h f(x) dx = y(h) - y(0).
$$

We now use our second-order Runge Kutta methods to construct elementary quadrature formulas.

(a) Midpoint rule

$$
k_1 = hf(0)
$$
,  $k_2 = hf(h/2)$ ,  $y(h) = y(0) + hf(h/2)$ .

The elementary quadrature formula (midpoint rule) is therefore given by

$$
\int_0^h f(x) \, dx = h f(h/2).
$$

(b) Modified Euler method

$$
k_1 = hf(0)
$$
,  $k_2 = hf(h)$ ,  $y(h) = y(0) + \frac{1}{2}h(f(0) + f(h))$ .

The elementary quadrature formula (trapezoidal rule) is therefore given by

$$
\int_0^h f(x) \, dx = \frac{1}{2} h \left( f(0) + f(h) \right).
$$

## **Solutions to the [Problems for Lecture 52](#page-121-0)**

**1.** Consider the ode given by  $dy/dx = f(x)$ , with  $y(0)$  as the initial value. Separating variables, and integrating from  $x = 0$  to 2*h* yields

$$
\int_{y(0)}^{y(2h)} dy = \int_0^{2h} f(x) dx, \quad \text{or} \quad \int_0^{2h} f(x) dx = y(2h) - y(0).
$$

We now use the fourth-order Runge Kutta methods to construct an elementary quadrature formula. We have

$$
k_1 = 2hf(0)
$$
,  $k_2 = 2hf(h)$ ,  $k_3 = 2hf(h)$ ,  $k_4 = 2hf(2h)$ ,

so that

$$
y(2h) = y(0) + \frac{h}{3} (f(0) + 4f(h) + f(2h)).
$$

The quadrature rule is therefore

$$
\int_0^{2h} f(x) dx = \frac{h}{3} (f(0) + 4f(h) + f(2h)),
$$

which is Simpson's rule.

## **Solutions to the [Problems for Lecture 53](#page-123-0)**

**1.**

$$
k_1 = \Delta t f(t_n, x_n, y_n, z_n), \qquad l_1 = \Delta t g(t_n, x_n, y_n, z_n), \qquad m_1 = \Delta t h(t_n, x_n, y_n, z_n);
$$

$$
k_2 = \Delta t f(t_n + \Delta t, x_n + k_1, y_n + l_1, z_n + m_1), \quad l_2 = \Delta t g(t_n + \Delta t, x_n + k_1, y_n + l_1, z_n + m_1),
$$

$$
m_2 = \Delta t h(t_n + \Delta t, x_n + k_1, y_n + l_1, z_n + m_1);
$$

$$
x_{n+1} = x_n + \frac{1}{2} (k_1 + k_2), \qquad y_{n+1} = y_n + \frac{1}{2} (l_1 + l_2), \qquad z_{n+1} = z_n + \frac{1}{2} (m_1 + m_2).
$$

**1.** Since the requested error is less than the estimated error, the current time step is rejected and redone using a smaller time step. The new time step is predicted to be

$$
\Delta \tau = \Delta t \left(\frac{\varepsilon}{e}\right)^{1/5} = 0.01 \times (1/1.1)^{1/5} = 0.0098112.
$$

For a margin of safety, one multiplies by the safety factor of 0.9 to obtain  $\Delta \tau = 0.0088301$ .

## **Solutions to the [Problems for Lecture 56](#page-129-0)**

**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
sigma=10; beta=8/3; r=28;
x0=1; y0=1; z0=1; tspan=[0 100];
ntrans=20;
options = odeset('RelTol',1.e-6);
[t,xyz]=ode45(@(t, xyz) lorenz_eqs(PROVIDE
                    FOUR ARGUMENTS HERE), PROVIDE TWO ARGUMENTS HERE, options);
x=xyz(ntrans:end,1); y=xyz(ntrans:end,2); z=xyz(ntrans:end,3);
plot3(x,y,z);
xlabel('$x$','Interpreter','latex','FontSize',14 );
ylabel('$y$','Interpreter','latex','FontSize',14 );
zlabel('$z$','Interpreter','latex','FontSize',14 );
title('Lorenz Equations','Interpreter','latex','FontSize',16);
function dxyzdt = lorenz_eqs(xyz,sigma,beta,r)x=xyz(1); y=xyz(2); z=xyz(3);
dxyzdt=
 [WRITE THE DIFFERENTIAL EQUATION HERE AS A COLUMN VECTOR. USE VARIABLES x, y AND z.];
end
```
Your plot should look like this:

# $\mbox{Lorenz}$  Equations



When viewed in the MATLAB environment, you can use your mouse to rotate the attractor for additional clarity.

**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
theta0=0; u0=2; %initial ode conditions. u0 is initial guess for root.
inf=8*pi; %inf is a large number. Takes a long time to get to top.
tspan=[0 inf];
options = odeset('RelTol',1.e-6);
%rootfind u0 such that theta(inf)=pi
alpha_i=linspace(0, 2, 100);
u0_i = zeros(100,1);for i=1:length(alpha_i)
    alpha=alpha_i(i);
    u0_i(i) = fzero(@(u0) F(tspan,theta0,u0,alpha,options), u0);
end
plot(alpha_i, u0_i);
xlabel('$\alpha$','Interpreter','latex','FontSize',14);
ylabel('$d \theta/dt$','Interpreter','latex','FontSize',14);
title('Shooting to the Pendulum Top','Interpreter','latex','FontSize',16);
function y=F(tspan,theta0,u0,alpha,options)
% use ode45 to define the root-finding problem
end
function d_theta_u_dt = pendulum(theta_u,alpha)
% define the differential equation here
end
```
#### Your plot should look like this:



#### **1.**

*a*) Consider  $m_1 = m_2$ . If  $\mathbf{r} = \cos(\omega t)\mathbf{i} + \sin(\omega t)\mathbf{j}$ , then

$$
\mathbf{r}_1 = \frac{1}{2} (\cos (\omega t) \mathbf{i} + \sin (\omega t) \mathbf{j}), \qquad \mathbf{r}_2 = -\frac{1}{2} (\cos (\omega t) \mathbf{i} + \sin (\omega t) \mathbf{j}).
$$

We can sketch the orbits as



*b*) Consider  $m_1 = 3m_2$ . If  $\mathbf{r} = \cos(\omega t)\mathbf{i} + \sin(\omega t)\mathbf{j}$ , then

$$
\mathbf{r}_1 = \frac{1}{4} (\cos (\omega t) \mathbf{i} + \sin (\omega t) \mathbf{j}), \qquad \mathbf{r}_2 = -\frac{3}{4} (\cos (\omega t) \mathbf{i} + \sin (\omega t) \mathbf{j}).
$$

We can sketch the orbits as



**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
e=0.7; m1=1; m2=4;
T=2*pi./(1-e).^1.5; tspan=linspace(0,T,1000);options=odeset('RelTol',1.e-6);
%%%%% Solve differential equations for x and y using ode45
%%%%% with arguments tspan and options.
%%%%% Determine x1, y1 and x2, y2
\mathbf{Q}\overline{\mathsf{S}}%
%%%%% graphics %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
k=0.1;R1=k*(m1)^(1/3); R2=k*(m2)^(1/3); %radius of masses
theta = linspace(0,2*pi);
figure; axis equal; hold on; set(gcf, 'color', 'w');
axis off;
xlim([-2,5]); ylim([-2.5,2.5]);
planet=fill(R1*cos(theta)+x1(1), R1*sin(theta)+y1(1),'b');
sun=fill(R2*cos(theta)+x2(1), R2*sin(theta)+y2(1),'r');
pause(1);
nperiods=5; %number of periods to plot
for j=1:nperiods
   for i=1:length(t)
        planet.XData=R1*cos(theta)+x1(i); planet.YData=R1*sin(theta)+y1(i);
        sun.XData=R2*cos(theta)+x2(i); sun.YData=R2*sin(theta)+y2(i);
        drawnow;
    end
end
%%%%% Write local function for differential equations %%%%%%%%%%%%%%
```
## **Solutions to the [Practice quiz: Classify partial differential equations](#page-138-0)**

**1.** a. The flow field is determined solely by the boundary conditions.

**2.** b. The turbulent motion is evolving in time.

**3.** a. The salt concentration is determined solely by the boundary conditions.

**4.** b. The salt concentration is evolving in time.

**5.** a. The potential is determined solely by the boundary conditions.

**6.** b. The potential is evolving in time.

#### **1.** We use

$$
y(x+2h) = y(x) + 2hy'(x) + 2h^2y''(x) + \frac{4}{3}h^3y'''(x) + \frac{2}{3}h^4y''''(x) + O(h^5),
$$
  
\n
$$
y(x+h) = y(x) + hy'(x) + \frac{1}{2}h^2y''(x) + \frac{1}{6}h^3y'''(x) + \frac{1}{24}h^4y''''(x) + O(h^5),
$$
  
\n
$$
y(x-h) = y(x) - hy'(x) + \frac{1}{2}h^2y''(x) - \frac{1}{6}h^3y'''(x) + \frac{1}{24}h^4y''''(x) + O(h^5),
$$
  
\n
$$
y(x-2h) = y(x) - 2hy'(x) + 2h^2y''(x) - \frac{4}{3}h^3y'''(x) + \frac{2}{3}h^4y''''(x) + O(h^5).
$$

Notice that

$$
y(x+h) - y(x-h) = 2hy'(x) + \frac{1}{3}h^3y'''(x) + O(h^5),
$$
  

$$
y(x+2h) - y(x-2h) = 4hy'(x) + \frac{8}{3}h^3y'''(x) + O(h^5).
$$

To eliminate the  $h^3$  term, we can multiply the first expression by eight and subtract the second expression to obtain

$$
8(y(x+h)-y(x-h)) - (y(x+2h)-y(x-2h)) = 12hy'(x) + O(h5).
$$

Solving for the first derivative and rearranging terms, we find

$$
y'(x) = \frac{-y(x+2h) + 8y(x+h) - 8y(x-h) + y(x-2h)}{12h} + O(h^4).
$$

# **Solutions to the [Problems for Lecture 62](#page-142-0)**

**1.** The discrete Laplace equation is given by

$$
4\Phi_{i,j} - \Phi_{i+1,j} - \Phi_{i-1,j} - \Phi_{i,j+1} - \Phi_{i,j-1} = 0.
$$

Solving for Φ*i*,*<sup>j</sup>* , we obtain

$$
\Phi_{i,j} = \frac{1}{4} \left( \Phi_{i+1,j} + \Phi_{i-1,j} + \Phi_{i,j+1} + \Phi_{i,j-1} \right).
$$

The right-hand side is just the average value of Φ at the neighboring four grid points.

**1.** The  $(i, j)$  coordinates for the four corners starting from the bottom left and moving clockwise are

$$
(i,j) = (1,1), (1,n_y), (n_x,n_y), (n_x,1).
$$

Using the mapping  $(i, j) \rightarrow k = i + (j - 1)n_x$ , the corresponding *k* coordinates are given by

$$
k=1, n_x n_y - n_x + 1, n_x n_y, n_x.
$$

# **Solutions to the [Problems for Lecture 64](#page-146-0)**

**1.** The grid and its *k* indexing is shown below.



#### The matrix equation is given by

$\left(1\right)$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi_1$												
$\overline{0}$	$\mathbf 1$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi_2$											
$\mathbf{0}$	$\boldsymbol{0}$	$\mathbf{1}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi_3$										
$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\,1$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi_4$									
$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\,1\,$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\mathbf{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi_5$
$\boldsymbol{0}$	$-1$	$\boldsymbol{0}$	$\boldsymbol{0}$	$-1$	$\overline{4}$	$-1$	$\boldsymbol{0}$	$\boldsymbol{0}$	$-1$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi_6$
$\boldsymbol{0}$	$\boldsymbol{0}$	$-1$	$\boldsymbol{0}$	$\mathbf{0}$	$-1$	$\overline{4}$	$-1$	$\boldsymbol{0}$	$\boldsymbol{0}$	$-1$	$\boldsymbol{0}$	$\mathbf{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi$ <sub>7</sub>
$\boldsymbol{0}$	$\mathbf{1}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi_8$						
$\boldsymbol{0}$	$\mathbf{1}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi_9$							
$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$-1$	$\boldsymbol{0}$	$\boldsymbol{0}$	$-1$	$\overline{4}$	$-1$	$\boldsymbol{0}$	$\boldsymbol{0}$	$-1$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi_{10}$
$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$-1$	$\boldsymbol{0}$	$\boldsymbol{0}$	$-1$	$\overline{4}$	$-1$	$\boldsymbol{0}$	$\overline{0}$	$-1$	$\boldsymbol{0}$	$\Phi_{11}$
$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\mathbf{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\mathbf{1}$	$\boldsymbol{0}$	$\mathbf{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi_{12}$
$\boldsymbol{0}$	$\mathbf{1}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi_{13}$											
$\boldsymbol{0}$	$\mathbf 1$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\Phi_{14}$												
$\boldsymbol{0}$	$\mathbf{0}$	$\mathbf{1}$	$\boldsymbol{0}$	$\Phi_{15}$												
$\overline{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\mathbf{0}$	$\boldsymbol{0}$	$\mathbf{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\boldsymbol{0}$	$\mathbf{0}$	$\mathbf{0}$	$\boldsymbol{0}$	$\frac{1}{2}$	$\Phi_{16}$
													$\begin{array}{c} b_1 \\ b_2 \\ b_3 \end{array}$ $b_4$ $b_5$ $\overline{0}$ $\overline{0}$ $b_8$ $=$ $b_9$ $\overline{0}$ $\overline{0}$ $b_{12}$ $b_{13}$ $b_{14}$ $b_{15}$ $b_{16}$			

**2.** The number of interior points is the number of points in a rectangle with sides  $n<sub>x</sub>$  − 2 and  $n_y - 2$  and is given by  $(n_x - 2)(n_y - 2)$ . The number of boundary points is given by  $2(n_x + n_y) - 4$ . When  $n_x = n_y = 100$ , there are 9,604 interior points and 396 boundary points. When  $n_x = n_y = 1000$ , there are 996,004 interior points and 3,996 boundary points. The percentage of boundary points to total number of grid points are, respectively, 3.96% and 0.3996%, or about 4% and 0.4%. Boundary points make up a small fraction of a large grid, which is why we can include them in our matrix without much additional cost.

**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
%%%%% Define the rectangle and grid parameters %%%%%%%%%%%%%%%%%%%%%%%%%%%%
Lx=1; Ly=1; %rectangle dimensions
Nx=100; Ny=100; %# of intervals
nx=Nx+1; ny=Ny+1; %# of gridpoints in x,y directions including boundaries
dx=Lx/Nx; dy=Ly/Ny; %grid size in x,y directions
x=(0:Nx)*dx; y=(0:Ny)*dy; %x, y values on the grid%%%%% Define the indices associated with the boundaries %%%%%%%%%%%%%%%%%%%
% boundary_index = [bottom, left, top, right]
boundary_index=[ 1:nx, 1:nx:1+(ny-1)*nx, ...
            1+(ny-1)*nx:nx*ny, nx:nx:nx*ny ];
%%%%% Set up matrix %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
diagonals = [4 \times ones(nx \times ny, 1), -ones(nx \times ny, 4)];
A=spdiags(diagonals, [0 -1 1 -nx nx], nx*ny, nx*ny); %use sparse matrices
I=speye(nx*ny);
A(boundary_index,:)=I(boundary_index,:);
%%%%% SET-UP RIGHT HAND SIDE %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
b=zeros(nx,ny);
b(:,1)=\ldots; & bottom
b(1,:)=...; % left
b(:,ny) = . . . ; %top
b(nx,:)=\ldots; %right
b=reshape(b,nx*ny,1); %make column vector
%%%%% Solve the Laplace equation using Gaussian elimination %%%%%%%%%%%%%%%
Phi=A\b; %solution step (all the computational time is here)
Phi=reshape(Phi,nx,ny); %make matrix
%%%%% Graphics %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
[X, Y]=meshgrid(x,y);
v=[DEFINE APPROPRIATE COUNTOUR LEVELS];
contour(X,Y,Phi',v,'ShowText','on');%requires transpose (read the notes)
axis equal;
set(gca, 'YTick', [0 0.2 0.4 0.6 0.8 1]);
set(gca, 'XTick', [0 0.2 0.4 0.6 0.8 1]);
xlabel('$x$','Interpreter','latex','FontSize',14 );
ylabel('$y$','Interpreter','latex','FontSize',14);
title('Solution of the Laplace equation','Interpreter','latex','FontSize',16);
```
## **Solutions to the [Problems for Lecture 66](#page-151-0)**

**1.**

$$
\begin{split} x_1^{(n+1)}&=\frac{1}{a_{11}}\left(b_1-a_{12}x_2^{(n)}-a_{13}x_3^{(n)}\right),\\ x_2^{(n+1)}&=\frac{1}{a_{22}}\left(b_2-a_{21}x_1^{(n)}-a_{23}x_3^{(n)}\right),\\ x_3^{(n+1)}&=\frac{1}{a_{33}}\left(b_3-a_{31}x_1^{(n)}-a_{32}x_2^{(n)}\right). \end{split}
$$

**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
%%%%% Define the rectangle and grid parameters %%%%%%%%%%%%%%%%%%%%%%%%%%%%
Lx=1; Ly=1; %rectangle dimensions
Nx=100; Ny=100; %# of intervals
nx=Nx+1; ny=Ny+1; %# of gridpoints in x,y directions including boundaries
dx=Lx/Nx; dy=Ly/Ny; %grid size in x,y directions
x=(0:Nx)*dx; y=(0:Ny)*dy; %x, y values on the grid
%%%%% Define the iteration parameters and initial condition %%%%%%%%%%%%%%%
eps=1.e-6; %convergence criteria for each value of Phi
index x=2:nx-1; index y=2:ny-1; %internal grid points
Phi=zeros(nx, ny); % matrix with solution and boundary conditions
%%set the boundary conditions
Phi(:,1)=\ldots; %bottom
Phi(1,:)=\ldots; % left
Phi(:, ny) = . . . ; %top
Phi(nx, :)=...; %right
\texttt{\$88888}Phi_old=Phi;
error=2*eps; ncount=0;
while (error > eps)
  ncount=ncount+1;
  Phi(index x,index y)=0.25*(Phi(index x+1,index y) ...
   +Phi(index_x-1,index_y)+Phi(index_x,index_y+1)+Phi(index_x,index_y-1));
  error=max(abs(Phi(:)-Phi_old(:)));
   if any (\text{isnan}(Phi(\cdot))) || any (\text{isinf}(Phi(\cdot)))fprintf('iterations diverge\n');
       return;
   end
   Phi_old=Phi;
    %fprintf('%g %e\n',ncount, error);
end
fprintf('%g\n',ncount);
%%%%% graphics %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
[X, Y]=meshgrid(x,y);
v=[]; %SET THE CONTOUR LEVELS
contour(X,Y,Phi',v,'ShowText','on');%requires transpose (read the notes)
axis equal;
set(gca, 'YTick', [0 0.2 0.4 0.6 0.8 1]);
set(gca, 'XTick', [0 0.2 0.4 0.6 0.8 1]);
xlabel('$x$','Interpreter','latex','FontSize',14 );
ylabel('$y$','Interpreter','latex','FontSize',14);
title('Solution of the Laplace equation','Interpreter','latex','FontSize',16);
```
#### **Solutions to the [Problems for Lecture 69](#page-156-0)**

**1.** The two-step solution of the one-dimensional diffusion equation using the modified

$$
\tilde{u}_{j}^{l+1} = u_{j}^{l} + \frac{\Delta t D}{(\Delta x)^{2}} \left( u_{j+1}^{l} - 2u_{j}^{l} + u_{j-1}^{l} \right),
$$
  

$$
u_{j}^{l+1} = u_{j}^{l} + \frac{\Delta t D}{2(\Delta x)^{2}} \left[ \left( u_{j+1}^{l} - 2u_{j}^{l} + u_{j-1}^{l} \right) + \left( \tilde{u}_{j+1}^{l+1} - 2\tilde{u}_{j}^{l+1} + \tilde{u}_{j-1}^{l+1} \right) \right].
$$

**2.** The FTCS scheme for the advection equation is given by

$$
u_j^{l+1} = u_j^l - \frac{c\Delta t}{2\Delta x} \left( u_{j+1}^l - u_{j-1}^l \right).
$$

#### **Solutions to the [Problems for Lecture 70](#page-158-0)**

**1.** We look for solutions of the form

$$
u_j^l = \xi^l e^{ikj\Delta x}.
$$

Substitution into the discrete advection equation and dividing by *ξ l e ikj*∆*x* results in

$$
\xi = 1 - \frac{c\Delta t}{2\Delta x} \left( e^{ik\Delta x} - e^{-ik\Delta x} \right) = 1 - i \frac{c\Delta t}{\Delta x} \sin (k\Delta x).
$$

The modulus of *ξ* is given by

$$
|\xi| = \sqrt{1 + \left(\frac{c\Delta t}{\Delta x}\right)^2 \sin^2(k\Delta x)}.
$$

The value of  $|\xi|$  is largest when  $\sin^2(k\Delta x) = 1$ , and since for this value  $|\xi| > 1$  no matter the value of ∆*t*, the method is always unstable.

## **Solutions to the [Problems for Lecture 71](#page-160-0)**

**1.**

*a*) The implicit scheme for the advection equation is given by

$$
u_j^{l+1} = u_j^l - \frac{c\Delta t}{2\Delta x} \left( u_{j+1}^{l+1} - u_{j-1}^{l+1} \right).
$$

*b*) We look for solutions of the form

$$
u_j^l = \xi^l e^{ikj\Delta x}.
$$

Substitution into the implicit discrete advection equation and dividing by *ξ l e ikj*∆*x* results in

$$
\xi = 1 - \frac{c\Delta t \xi}{2\Delta x} \left( e^{ik\Delta x} - e^{-ik\Delta x} \right) = 1 - i \frac{c\Delta t \xi}{\Delta x} \sin(k\Delta x).
$$

Solving for *ξ*, we obtain

$$
\xi = \frac{1}{1 + i \frac{c\Delta t}{\Delta x} \sin (k\Delta x)}.
$$

The modulus of *ξ* is given by

$$
|\xi| = \frac{1}{\sqrt{1 + \left(\frac{c\Delta t}{\Delta x}\right)^2 \sin^2(k\Delta x)}}.
$$

The value of  $|\xi|$  is largest when  $\sin^2(k\Delta x) = 0$ , and for this value  $|\xi| = 1$ . The method is unconditionally stable.

## **Solutions to the [Problems for Lecture 72](#page-162-0)**

**1.** The Lax scheme for the advection equation is given by

$$
u_j^{l+1} = \frac{1}{2}(u_{j+1}^l + u_{j-1}^l) - \frac{c\Delta t}{2\Delta x}(u_{j+1}^l - u_{j-1}^l).
$$

We look for solutions of the form

$$
u_j^l = \xi^l e^{ikj\Delta x}.
$$

Substitution into the Lax scheme and dividing by  $\tilde{\zeta}^l e^{ikj\Delta x}$  results in

$$
\zeta = \frac{1}{2} \left( e^{ik\Delta x} + e^{-ik\Delta x} \right) - \frac{c\Delta t}{2\Delta x} \left( e^{ik\Delta x} - e^{-ik\Delta x} \right)
$$

$$
= \cos (k\Delta x) - i\frac{c\Delta t}{\Delta x} \sin (k\Delta x).
$$

The modulus of *ξ* is given by

$$
|\xi| = \sqrt{\cos^2(k\Delta x) + \left(\frac{c\Delta t}{\Delta x}\right)^2 \sin^2(k\Delta x)}.
$$

Clearly, the Lax scheme is stable provide that

$$
\left(\frac{c\Delta t}{\Delta x}\right)^2 \le 1.
$$

which is called the Courant-Friedrichs-Lewy (CFL) stability criterion.

# **Solutions to the [Problems for Lecture 73](#page-165-0)**

**1.**

*a*) For *x* a boundary point on the left, we start with

$$
y(x+h) = y(x) + hy'(x) + \frac{1}{2}h^2y''(x) + O(h^3),
$$
  

$$
y(x+2h) = y(x) + 2hy'(x) + 2h^2y''(x) + O(h^3).
$$

To eliminate the term proportional to  $h^2$ , we multiply the first equation by four and subtract the second equation to obtain

$$
4y(x+h) - y(x+2h) = 3y(x) + 2hy'(x) + O(h3).
$$

Solving for  $y'(x)$ , we find

$$
y'(x) = \frac{-3y(x) + 4y(x+h) - y(x+2h)}{2h} + O(h^2).
$$

For *x* a boundary point on the right, we start with

$$
y(x - h) = y(x) - hy'(x) + \frac{1}{2}h^2y''(x) + O(h^3),
$$
  

$$
y(x - 2h) = y(x) - 2hy'(x) + 2h^2y''(x) + O(h^3).
$$

We multiply the first equation by four and subtract the second equation to obtain

$$
4y(x - h) - y(x - 2h) = 3y(x) - 2hy'(x) + O(h3).
$$

Solving for  $y'(x)$ , we find

$$
y'(x) = \frac{3y(x) - 4y(x - h) + y(x - 2h)}{2h} + O(h^{2}).
$$

*b*) Written in discretized form, the second-order method for the *x*- derivative at boundary points become

$$
y_1'=\frac{-3y_1+4y_2-y_3}{2\Delta x}, \qquad y_{n_x}'=\frac{3y_{n_x}-4y_{n_x-1}+y_{n_x-2}}{2\Delta x}.
$$

We apply these equations to  $u_j^l$  and set the partial derivatives with respect to *x* to be zero at the boundary points. We can then solve for *u* at the boundaries to obtain

$$
u_1^l = \frac{1}{3}(4u_2^l - u_3^l), \qquad u_{n_x}^l = \frac{1}{3}(4u_{n_x-1}^l - u_{n_x-2}^l).
$$

#### **2.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
D=1; %diffusion coefficient
%%%%% Define the x-domain and x-grid %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
Lx=1; %domain: -Lx < x < LxNx=500; %# of intervals
nx=Nx+1;%# of gridpoints in x
dx=2*Lx/Nx; %grid length in x
x=-Lx + (0:Nx)*dx; %x values on the grid
%%%%% Time step parameters %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
nsteps=10000; %number of time steps
nout=500; %plot every nout time steps
```

```
dt = (dx)^2/(2*D); %borderline stability of FTCS scheme
alpha=dt*D/dx^2; %equation parameter
%%%%% Construct the matrix %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
diagonals = [2*(1+a1pha)*ones(nx,1), -alpha*ones(nx,2)];A=spdiags(diagonals,[0 -1 1], nx, nx);
I=speye(nx);
A([1 nx],:)=I([1 nx],:); %boundaries
%%%%% Define initial conditions and plot %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
sigma=Lx/16;
u=1/(sigma*sqrt(2*pi))*exp(-0.5*(x/sigma).^2); u=u';
plot(x,u,'r'); hold on;
xlabel('$x$','Interpreter','latex','FontSize',14);
ylabel('$u(x, t)$','Interpreter','latex','FontSize',14);
title('Solution of the diffusion equation','Interpreter','latex','FontSize',16);
%%%%% Advance solution and plot %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
for m=1:nsteps
    b=[NO-FLUX BOUNDARY CONDITION; ...
        [alpha*u(1:nx-2) + 2*(1-alpha)*u(2:nx-1) + alpha*u(3:nx)]; ...NO-FLUX BOUNDARY CONDITION];
    u=A\b;
    if mod(m, nout) == 0, plot(x, u, 'b'), end
end
```
#### Your plot should look like this:


## **Solutions to the [Problems for Lecture 74](#page-168-0)**

**1.** Complete your solution on Coursera using Matlab. Here is the Learner Template:

```
%%%%% Define the square and grid parameters %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
L=1; %square is 2L x 2L
N=100; %# of intervals in x and y directions
n=N+1; %# of gridpoints in x,y directions including boundaries
h=2*L/N; %grid size in x,y directions
x=-L + (0:N) *h; %x values on the grid
y=-L + (0:N) *h; %y values on the grid
[X, Y] =meshqrid(x, y);
%%%%% Define the indices associated with the boundaries %%%%%%%%%%%%%%%%%%%
% boundary_index = [bottom, left, top, right]
boundary_index=[ 1:n, 1:n, 1:n:1+(n-1)*n, ...
                 1+(n-1)*n:n*n, n:n:n*n ];
%%%%% Diffusion constant and time-step parameters
D=1:
dt=h^2/(2*D); %borderline stability of FTCS scheme
alpha=dt*D/h^2; %equation parameter
nsteps=1000; %number of time steps
%%%%% CONSTRUCT THE MATRIX AND COMPUTE LU DECOMPOSITION %%%%%%%%%%%%%%%%%%%%
%
\frac{1}{2}\frac{1}{2}\frac{6}{10}\frac{1}{2}%%%%% Define initial conditions %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
u=zeros(n,n,nsteps);
sigma=L/4;
u(:,:,1) = 1/(2*pi*sigma^2)*exp(-0.5*(X.^2+Y.^2)/sigma^2);u(1,:,1)=0; u(n,:,1)=0; u(:,1,1)=0; u(:,n,1)=0; %b.c.
%%%%% ADVANCE SOLUTION u %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
for m=2:nsteps
\approx\frac{1}{2}\frac{1}{2}\frac{6}{10}\frac{1}{2}\frac{1}{2}end
%%%% Plot with animation %%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%%
figure('units','normalized','outerposition',[0 0 1 1])
s=surf(X, Y, u(:, :, 1)); zlim([0, 2.6]);
xlabel('$x$','Interpreter','latex','FontSize',14);
ylabel('$y$','Interpreter','latex','FontSize',14);
zlabel('$u(x,y,t)$','Interpreter','latex','FontSize',14);
title('Solution of the 2D diffusion equation','Interpreter','latex','FontSize',16);
pause(1)
for j=2:nsteps
    s.ZData=u(:,:,j); pause(0.01);
end
```